8.5 Sequencing Problems

Basic genres.

- Packing problems: SET-PACKING, INDEPENDENT SET.
- Covering problems: SET-COVER, VERTEX-COVER.
- Constraint satisfaction problems: SAT, 3-SAT.
- **Sequencing problems**: HAMILTONIAN-CYCLE, TSP.
- Partitioning problems: 3D-MATCHING, 3-COLOR.
- Numerical problems: SUBSET-SUM, KNAPSACK.
Hamiltonian Cycle

**HAM-CYCLE:** given an undirected graph $G = (V, E)$, does there exist a simple cycle $\Gamma$ that contains every node in $V$.

**YES:** vertices and faces of a dodecahedron.
**Hamiltonian Cycle**

**HAM-CYCLE:** given an undirected graph $G = (V, E)$, does there exist a simple cycle $\Gamma$ that contains every node in $V$.

![Graph example](image)

**NO:** bipartite graph with odd number of nodes.
**Directed Hamiltonian Cycle**

**DIR-HAM-CYCLE:** given a digraph $G = (V, E)$, does there exists a simple directed cycle $\Gamma$ that contains every node in $V$?

**Claim.** $\text{DIR-HAM-CYCLE} \leq_p \text{HAM-CYCLE}$. 

**Pf.** Given a directed graph $G = (V, E)$, construct an undirected graph $G'$ with $3n$ nodes.
Claim. $G$ has a Hamiltonian cycle iff $G'$ does.

Pf. $\Rightarrow$
- Suppose $G$ has a directed Hamiltonian cycle $\Gamma$.
- Then $G'$ has an undirected Hamiltonian cycle (same order).

Pf. $\Leftarrow$
- Suppose $G'$ has an undirected Hamiltonian cycle $\Gamma'$.
- $\Gamma'$ must visit nodes in $G'$ using one of following two orders:
  - ..., B, G, R, B, G, R, B, G, R, B, ...
- Blue nodes in $\Gamma'$ make up directed Hamiltonian cycle $\Gamma$ in $G$, or reverse of one. □
Claim. 3-SAT $\leq_p$ DIR-HAM-CYCLE.

Pf. Given an instance $\Phi$ of 3-SAT, we construct an instance of DIR-HAM-CYCLE that has a Hamiltonian cycle iff $\Phi$ is satisfiable.

Construction. First, create graph that has $2^n$ Hamiltonian cycles which correspond in a natural way to $2^n$ possible truth assignments.
Construction. Given 3-SAT instance $\Phi$ with $n$ variables $x_i$ and $k$ clauses.

- Construct $G$ to have $2^n$ Hamiltonian cycles.
- Intuition: traverse path $i$ from left to right $\Leftrightarrow$ set variable $x_i = 1$. 

![Diagram](image-url)
3-SAT Reduces to Directed Hamiltonian Cycle

Construction. Given 3-SAT instance $\Phi$ with $n$ variables $x_i$ and $k$ clauses.
- For each clause: add a node and 6 edges.

\[ C_1 = x_1 \lor \overline{x_2} \lor x_3 \]

\[ C_2 = \overline{x_1} \lor x_2 \lor \overline{x_3} \]
3-SAT Reduces to Directed Hamiltonian Cycle

Claim. \( \Phi \) is satisfiable iff \( G \) has a Hamiltonian cycle.

Pf. \( \Rightarrow \)

- Suppose 3-SAT instance has satisfying assignment \( x^* \).
- Then, define Hamiltonian cycle in \( G \) as follows:
  - if \( x^*_i = 1 \), traverse row \( i \) from left to right
  - if \( x^*_i = 0 \), traverse row \( i \) from right to left
  - for each clause \( C_j \), there will be at least one row \( i \) in which we are going in "correct" direction to splice node \( C_j \) into tour
3-SAT Reduces to Directed Hamiltonian Cycle

Claim. $\Phi$ is satisfiable iff $G$ has a Hamiltonian cycle.

Pf. $\Leftarrow$

- Suppose $G$ has a Hamiltonian cycle $\Gamma$.
- If $\Gamma$ enters clause node $C_j$, it must depart on mate edge.
  - thus, nodes immediately before and after $C_j$ are connected by an edge $e$ in $G$
  - removing $C_j$ from cycle, and replacing it with edge $e$ yields Hamiltonian cycle on $G - \{C_j\}$
- Continuing in this way, we are left with Hamiltonian cycle $\Gamma'$ in $G - \{C_1, C_2, \ldots, C_k\}$.
- Set $x^*_i = 1$ iff $\Gamma'$ traverses row $i$ left to right.
- Since $\Gamma$ visits each clause node $C_j$, at least one of the paths is traversed in "correct" direction, and each clause is satisfied.
Longest Path

**SHORTEST-PATH.** Given a digraph $G = (V, E)$, does there exists a simple path of length at most $k$ edges?

**LONGEST-PATH.** Given a digraph $G = (V, E)$, does there exists a simple path of length at least $k$ edges?

**Claim.** $3$-SAT $\leq_p$ $\text{LONGEST-PATH}$. 

**Pf 1.** Redo proof for $\text{DIR-HAM-CYCLE}$, ignoring back-edge from $t$ to $s$. 

**Pf 2.** Show $\text{HAM-CYCLE} \leq_p \text{LONGEST-PATH}$. 
The Longest Path †

Lyrics.  Copyright © 1988 by Daniel J. Barrett.

Music.  Sung to the tune of The Longest Time by Billy Joel.

Woh-oh-oh-oh, find the longest path!
Woh-oh-oh-oh, find the longest path!

If you said P is NP tonight,
There would still be papers left to write,
I have a weakness,
I'm addicted to completeness,
And I keep searching for the longest path.

The algorithm I would like to see
Is of polynomial degree,
But it's elusive:
Nobody has found conclusive
Evidence that we can find a longest path.

I have been hard working for so long.
I swear it's right, and he marks it wrong.
Some how I'll feel sorry when it's done:
GPA 2.1
Is more than I hope for.

Garey, Johnson, Karp and other men (and women)
Tried to make it order N log N.
Am I a mad fool
If I spend my life in grad school,
Forever following the longest path?

Woh-oh-oh-oh, find the longest path!
Woh-oh-oh-oh, find the longest path!
Woh-oh-oh-oh, find the longest path.

† Recorded by Dan Barrett while a grad student at Johns Hopkins during a difficult algorithms final.
Traveling Salesperson Problem

**TSP.** Given a set of \( n \) cities and a pairwise distance function \( d(u, v) \), is there a tour of length \( \leq D \)?

All 13,509 cities in US with a population of at least 500
Reference: http://www.tsp.gatech.edu
Traveling Salesperson Problem

**TSP.** Given a set of $n$ cities and a pairwise distance function $d(u, v)$, is there a tour of length $\leq D$?

Reference: [http://www.tsp.gatech.edu](http://www.tsp.gatech.edu)
Traveling Salesperson Problem

TSP. Given a set of $n$ cities and a pairwise distance function $d(u, v)$, is there a tour of length $\leq D$?

11,849 holes to drill in a programmed logic array
Reference: http://www.tsp.gatech.edu
Traveling Salesperson Problem

**TSP.** Given a set of $n$ cities and a pairwise distance function $d(u, v)$, is there a tour of length $\leq D$?

Optimal TSP tour
Reference: http://www.tsp.gatech.edu
Traveling Salesperson Problem

**TSP.** Given a set of \( n \) cities and a pairwise distance function \( d(u, v) \), is there a tour of length \( \leq D \)?

**HAM-CYCLE:** given a graph \( G = (V, E) \), does there exists a simple cycle that contains every node in \( V \)?

**Claim.** \( \text{HAM-CYCLE} \leq_p \text{TSP} \).

**Pf.**
- Given instance \( G = (V, E) \) of \( \text{HAM-CYCLE} \), create \( n \) cities with distance function
  \[
  d(u, v) = \begin{cases} 
  1 & \text{if } (u, v) \in E \\
  2 & \text{if } (u, v) \notin E
  \end{cases}
  \]
- TSP instance has tour of length \( \leq n \) iff \( G \) is Hamiltonian.

**Remark.** TSP instance in reduction satisfies \( \Delta \)-inequality.
8.6 Partitioning Problems

Basic genres.

- Packing problems: SET-PACKING, INDEPENDENT SET.
- Covering problems: SET-COVER, VERTEX-COVER.
- Constraint satisfaction problems: SAT, 3-SAT.
- Sequencing problems: HAMILTONIAN-CYCLE, TSP.
- **Partitioning problems:** 3D-MATCHING, 3-COLOR.
- Numerical problems: SUBSET-SUM, KNAPSACK.
3-Dimensional Matching

**3D-MATCHING.** Given \( n \) instructors, \( n \) courses, and \( n \) times, and a list of the possible courses and times each instructor is willing to teach, is it possible to make an assignment so that all courses are taught at different times?

<table>
<thead>
<tr>
<th>Instructor</th>
<th>Course</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>Wayne</td>
<td>COS 423</td>
<td>MW 11-12:20</td>
</tr>
<tr>
<td>Wayne</td>
<td>COS 423</td>
<td>TTh 11-12:20</td>
</tr>
<tr>
<td>Wayne</td>
<td>COS 226</td>
<td>TTh 11-12:20</td>
</tr>
<tr>
<td>Wayne</td>
<td>COS 126</td>
<td>TTh 11-12:20</td>
</tr>
<tr>
<td>Tardos</td>
<td>COS 523</td>
<td>TTh 3-4:20</td>
</tr>
<tr>
<td>Tardos</td>
<td>COS 423</td>
<td>TTh 11-12:20</td>
</tr>
<tr>
<td>Tardos</td>
<td>COS 423</td>
<td>TTh 3-4:20</td>
</tr>
<tr>
<td>Kleinberg</td>
<td>COS 226</td>
<td>TTh 3-4:20</td>
</tr>
<tr>
<td>Kleinberg</td>
<td>COS 226</td>
<td>MW 11-12:20</td>
</tr>
<tr>
<td>Kleinberg</td>
<td>COS 423</td>
<td>MW 11-12:20</td>
</tr>
</tbody>
</table>
3-Dimensional Matching

**3D-MATCHING.** Given disjoint sets $X$, $Y$, and $Z$, each of size $n$ and a set $T \subseteq X \times Y \times Z$ of triples, does there exist a set of $n$ triples in $T$ such that each element of $X \cup Y \cup Z$ is in exactly one of these triples?

**Claim.** $3\text{-SAT} \leq_p 3\text{D-MATCHING}$.  

**Pf.** Given an instance $\Phi$ of $3\text{-SAT}$, we construct an instance of $3\text{D-matching}$ that has a perfect matching iff $\Phi$ is satisfiable.
3-Dimensional Matching

Construction. (part 1)

- Create gadget for each variable $x_i$ with $2k$ core and tip elements.
- No other triples will use core elements.
- In gadget $i$, 3D-matching must use either both grey triples or both blue ones.

$k = 2$ clauses
$n = 3$ variables

true

false

clause 1 tips

core

set $x_i = \text{true}$

set $x_i = \text{false}$

Number of clauses
Construction. (part 2)

- For each clause $C_j$ create two elements and three triples.
- Exactly one of these triples will be used in any 3D-matching.
- Ensures any 3D-matching uses either (i) grey core of $x_1$ or (ii) blue core of $x_2$ or (iii) grey core of $x_3$.

\[ C_j = x_1 \lor \bar{x}_2 \lor x_3 \]
3-Dimensional Matching

Construction. (part 3)

- For each tip, add a cleanup gadget.

```
  x_1   x_2   x_3
```

```
clause 1 gadget
```

```
true  false
```

```
clause 1 tips
```

```
cleanup gadget
```

```
core
```

```
z
```
3-Dimensional Matching

**Claim.** Instance has a 3D-matching iff $\Phi$ is satisfiable.

**Detail.** What are $X$, $Y$, and $Z$? Does each triple contain one element from each of $X$, $Y$, $Z$?
3-Dimensional Matching

Claim. Instance has a 3D-matching iff $\Phi$ is satisfiable.

Detail. What are $X$, $Y$, and $Z$? Does each triple contain one element from each of $X$, $Y$, $Z$?
8.7 Graph Coloring

Basic genres.

- Packing problems: SET-PACKING, INDEPENDENT SET.
- Covering problems: SET-COVER, VERTEX-COVER.
- Constraint satisfaction problems: SAT, 3-SAT.
- Sequencing problems: HAMILTONIAN-CYCLE, TSP.
- Partitioning problems: 3D-MATCHING, 3-COLOR.
- Numerical problems: SUBSET-SUM, KNAPSACK.
3-Colorability

3-COLOR: Given an undirected graph $G$ does there exists a way to color the nodes red, green, and blue so that no adjacent nodes have the same color?
Register Allocation

Register allocation. Assign program variables to machine register so that no more than $k$ registers are used and no two program variables that are needed at the same time are assigned to the same register.

Interference graph. Nodes are program variables names, edge between $u$ and $v$ if there exists an operation where both $u$ and $v$ are "live" at the same time.

Observation. [Chaitin 1982] Can solve register allocation problem iff interference graph is $k$-colorable.

Fact. $3$-COLOR $\leq_p$ $k$-REGISTER-ALLOCATION for any constant $k \geq 3$. 
3-Colorability

Claim. \( 3\text{-}\text{SAT} \leq_p 3\text{-}\text{COLOR} \).

Pf. Given 3-SAT instance \( \Phi \), we construct an instance of 3-COLOR that is 3-colorable iff \( \Phi \) is satisfiable.

Construction.

i. For each literal, create a node.

ii. Create 3 new nodes T, F, B; connect them in a triangle, and connect each literal to B.

iii. Connect each literal to its negation.

iv. For each clause, add gadget of 6 nodes and 13 edges.

\[ \uparrow \]

\text{to be described next}
3-Colorability

**Claim.** Graph is 3-colorable iff $\Phi$ is satisfiable.

**Pf.** Suppose graph is 3-colorable.

- Consider assignment that sets all T literals to true.
- (ii) ensures each literal is T or F.
- (iii) ensures a literal and its negation are opposites.
3-Colorability

Claim. Graph is 3-colorable iff $\Phi$ is satisfiable.

Pf. $\Rightarrow$ Suppose graph is 3-colorable.
- Consider assignment that sets all T literals to true.
- (ii) ensures each literal is T or F.
- (iii) ensures a literal and its negation are opposites.
- (iv) ensures at least one literal in each clause is T.

![Diagram of 3-Colorability](image)

$$C_i = x_1 \lor \overline{x_2} \lor x_3$$
Claim. Graph is 3-colorable iff $\Phi$ is satisfiable.

Pf. $\Rightarrow$ Suppose graph is 3-colorable.

- Consider assignment that sets all $T$ literals to true.
- (ii) ensures each literal is $T$ or $F$.
- (iii) ensures a literal and its negation are opposites.
- (iv) ensures at least one literal in each clause is $T$.

\[ C_i = x_1 \lor \overline{x_2} \lor x_3 \]

not 3-colorable if all are red

contradiction
3-Colorability

Claim. Graph is 3-colorable iff \( \Phi \) is satisfiable.

Pf. \( \Leftarrow \) Suppose 3-SAT formula \( \Phi \) is satisfiable.
   - Color all true literals T.
   - Color node below green node F, and node below that B.
   - Color remaining middle row nodes B.
   - Color remaining bottom nodes T or F as forced. \( \blacksquare \)

\[ C_i = x_1 \lor \neg x_2 \lor x_3 \]
8.8 Numerical Problems

Basic genres.

- Packing problems: SET-PACKING, INDEPENDENT SET.
- Covering problems: SET-COVER, VERTEX-COVER.
- Constraint satisfaction problems: SAT, 3-SAT.
- Sequencing problems: HAMILTONIAN-CYCLE, TSP.
- Partitioning problems: 3-COLOR, 3D-MATCHING.
- Numerical problems: SUBSET-SUM, KNAPSACK.
Subset Sum

**SUBSET-SUM.** Given natural numbers $w_1, \ldots, w_n$ and an integer $W$, is there a subset that adds up to exactly $W$?

**Ex:** $\{1, 4, 16, 64, 256, 1040, 1041, 1093, 1284, 1344\}$, $W = 3754$.

**Yes.** $1 + 16 + 64 + 256 + 1040 + 1093 + 1284 = 3754$.

**Remark.** With arithmetic problems, input integers are encoded in binary. Polynomial reduction must be polynomial in binary encoding.

**Claim.** $3$-SAT $\leq_P$ SUBSET-SUM.

**Pf.** Given an instance $\Phi$ of $3$-SAT, we construct an instance of SUBSET-SUM that has solution iff $\Phi$ is satisfiable.
Construction. Given 3-SAT instance $\Phi$ with $n$ variables and $k$ clauses, form $2n + 2k$ decimal integers, each of $n+k$ digits, as illustrated below.

Claim. $\Phi$ is satisfiable iff there exists a subset that sums to $W$.

Pf. No carries possible.

\[ C_1 = \overline{x} \lor y \lor z \]
\[ C_2 = x \lor \overline{y} \lor z \]
\[ C_3 = \overline{x} \lor \overline{y} \lor \overline{z} \]

\[
\begin{array}{cccc|ccc}
  x & y & z & C_1 & C_2 & C_3 \\
  \hline
  x & 1 & 0 & 0 & 0 & 1 & 0 & 100,110 \\
  \neg x & 1 & 0 & 0 & 1 & 0 & 1 & 100,001 \\
  y & 0 & 1 & 0 & 1 & 0 & 0 & 10,000 \\
  \neg y & 0 & 1 & 0 & 0 & 1 & 1 & 10,111 \\
  z & 0 & 0 & 1 & 1 & 1 & 0 & 1,010 \\
  \neg z & 0 & 0 & 1 & 0 & 0 & 1 & 1,101 \\
  0 & 0 & 0 & 1 & 0 & 0 & 100 \\
  0 & 0 & 0 & 2 & 0 & 0 & 200 \\
  0 & 0 & 0 & 0 & 1 & 0 & 10 \\
  0 & 0 & 0 & 0 & 2 & 0 & 20 \\
  0 & 0 & 0 & 0 & 0 & 1 & 1 \\
  0 & 0 & 0 & 0 & 0 & 2 & 2 \\
  1 & 1 & 1 & 4 & 4 & 4 & 111,444 \\
\end{array}
\]
Scheduling With Release Times

\textbf{SCHEDULE-RELEASE-TIMES.} Given a set of \( n \) jobs with processing time \( t_i \), release time \( r_i \), and deadline \( d_i \), is it possible to schedule all jobs on a single machine such that job \( i \) is processed with a contiguous slot of \( t_i \) time units in the interval \([r_i, d_i] \)?

\textbf{Claim.} \( \text{SUBSET-SUM} \leq_p \text{SCHEDULE-RELEASE-TIMES} \).
\textbf{Pf.} Given an instance of \( \text{SUBSET-SUM} \ w_1, \ldots, w_n \), and target \( W \),
\begin{itemize}
  \item Create \( n \) jobs with processing time \( t_i = w_i \), release time \( r_i = 0 \), and no deadline \( (d_i = 1 + \sum_j w_j) \).
  \item Create job 0 with \( t_0 = 1 \), release time \( r_0 = W \), and deadline \( d_0 = W+1 \).
\end{itemize}

\textit{Can schedule jobs 1 to n anywhere but \([W, W+1]\)}

\begin{tikzpicture}
  \draw[fill=gray!50] (0,0) rectangle (10,1);
  \draw[fill=blue] (1,0) rectangle (2,1);
  \draw[very thick, ->] (1.5,0.5) -- (2.5,0.5);
  \draw[very thick, ->] (1.5,0.5) -- (0.5,0.5);
  \node at (0.5,0.5) {job 0};
  \node at (1,0) {0}; \node at (2,0) {W}; \node at (3,0) {W+1}; \node at (4,0) {S+1};
\end{tikzpicture}
8.10 A Partial Taxonomy of Hard Problems
Polynomial-Time Reductions

3-SAT

INDEPENDENT SET

VERTEX COVER

SET COVER

packing and covering

constraint satisfaction

3-SAT reduces to INDEPENDENT SET

DIR-HAM-CYCLE

HAM-CYCLE

TSP

sequencing

GRAPH 3-COLOR

PLANAR 3-COLOR

partitioning

SUBSET-SUM

SCHEDULING

numerical

Dick Karp (1972)
1985 Turing Award
Extra Slides
Construction. Let \( X \cup Y \cup Z \) be an instance of 3D-MATCHING with triplet set \( T \). Let \( n = |X| = |Y| = |Z| \) and \( m = |T| \).

- Let \( X = \{ x_1, x_2, x_3, x_4 \} \), \( Y = \{ y_1, y_2, y_3, y_4 \} \), \( Z = \{ z_1, z_2, z_3, z_4 \} \).
- For each triplet \( t = (x_i, y_j, z_k) \in T \), create an integer \( w_t \) with \( 3n \) digits that has a 1 in positions \( i \), \( n+j \), and \( 2n+k \).

\[ \uparrow \]

use base \( m+1 \)

Claim. 3D-matching iff some subset sums to \( W = 111, \ldots, 111 \).
SUBSET-SUM. Given natural numbers $w_1, ..., w_n$ and an integer $W$, is there a subset that adds up to exactly $W$?

PARTITION. Given natural numbers $v_1, ..., v_m$, can they be partitioned into two subsets that add up to the same value?

Claim. SUBSET-SUM $\leq_p$ PARTITION.

Pf. Let $W, w_1, ..., w_n$ be an instance of SUBSET-SUM.

- Create instance of PARTITION with $m = n+2$ elements.
  - $v_1 = w_1$, $v_2 = w_2$, ..., $v_n = w_n$, $v_{n+1} = 2 \sum_i w_i - W$, $v_{n+2} = \sum_i w_i + W$

- There exists a subset that sums to $W$ iff there exists a partition since two new elements cannot be in the same partition.

\[ v_{n+1} = 2 \sum_i w_i - W \quad W \quad \text{subset } A \]
\[ v_{n+2} = \sum_i w_i + W \quad \sum_i w_i - W \quad \text{subset } B \]
Extra Slides: 4 Color Theorem
Planar 3-Colorability

**PLANAR-3-COLOR.** Given a planar map, can it be colored using 3 colors so that no adjacent regions have the same color?

YES instance.
PLANAR-3-COLOR. Given a planar map, can it be colored using 3 colors so that no adjacent regions have the same color?

NO instance.
Def. A graph is **planar** if it can be embedded in the plane in such a way that no two edges cross.

**Applications:** VLSI circuit design, computer graphics.

**Kuratowski's Theorem.** An undirected graph $G$ is non-planar iff it contains a subgraph homeomorphic to $K_5$ or $K_{3,3}$. 

homeomorphic to $K_{3,3}$ →
Planarity testing. [Hopcroft-Tarjan 1974] \(O(n)\).

\[\uparrow\]

simple planar graph can have at most 3n edges

Remark. Many intractable graph problems can be solved in poly-time if the graph is planar; many tractable graph problems can be solved faster if the graph is planar.
Planar 3-Colorability

Claim. \(3\text{-COLOR} \leq_P \text{PLANAR-3-COLOR}\).

Proof sketch: Given instance of \(3\text{-COLOR}\), draw graph in plane, letting edges cross if necessary.

- Replace each edge crossing with the following planar gadget \(W\).
  - in any 3-coloring of \(W\), opposite corners have the same color
  - any assignment of colors to the corners in which opposite corners have the same color extends to a 3-coloring of \(W\)
Planar $k$-Colorability

**PLANAR-2-COLOR.** Solvable in linear time.

**PLANAR-3-COLOR.** NP-complete.

**PLANAR-4-COLOR.** Solvable in $O(1)$ time.

**Theorem.** [Appel-Haken, 1976] Every planar map is 4-colorable.

- Resolved century-old open problem.
- Used 50 days of computer time to deal with many special cases.
- First major theorem to be proved using computer.

**False intuition.** If PLANAR-3-COLOR is hard, then so is PLANAR-4-COLOR and PLANAR-5-COLOR.
Graph minor theorem. [Robertson-Seymour 1980s]

Corollary. There exist an $O(n^3)$ algorithm to determine if a graph can be embedded in the torus in such a way that no two edges cross.

Pf of theorem. Tour de force.
Polynomial-Time Detour

**Graph minor theorem.** [Robertson-Seymour 1980s]

**Corollary.** There exist an $O(n^3)$ algorithm to determine if a graph can be embedded in the torus in such a way that no two edges cross.

**Mind boggling fact 1.** The proof is highly non-constructive!
**Mind boggling fact 2.** The constant of proportionality is enormous!

Unfortunately, for any instance $G = (V, E)$ that one could fit into the known universe, one would easily prefer $n^{70}$ to even constant time, if that constant had to be one of Robertson and Seymour’s. - David Johnson

**Theorem.** There exists an explicit $O(n)$ algorithm.
**Practice.** LEDA implementation guarantees $O(n^3)$. 