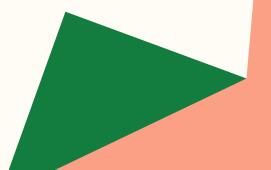
MPI and Shared Memory

Lecture 7 March 4, 2025



Reading for next time

Program 3 submission on Thursday

Questions about program 3?

To Dos

Review: MPI Communication

- Communicator
- Point-to-Point
 - Send
 - Receive
- Collective communication
 - MPI_Reduce, MPI_Allreduce
 - MPI_Bcast
 - MPI_Scatter
 - MPI_Gather, MPI_Allgather
 - MPI_Barrier

Communication vs. Computation

- Communication way more expensive than local computation
- Every message has a high fixed overhead cost for sending with some additional time dependent on payload size
 - latency = fixed_overhead_1_byte + (payload-size / bandwidth)
 - 2 messages will take longer than 1 message with same data

Data Consolidation in Messages

- count parameter to communication functions
 - For sending of arrays
- Derived datatypes
 - Like creating a struct for the data values you need to send
 - Sequence of basic MPI datatypes with a displacement for each
 - e.g., {(MPI_DOUBLE, 0), (MPI_DOUBLE, 16), (MPI_INT, 24)}
- MPI_Pack/Unpack

Derived Datatypes

- count number of elements in datatype
- array_of_block_lengths number of elements in each datatype
- array_of_displacements displacement in bytes from start of first data
 - Use MPI_GetAddress (void *location_p, MPI_AInt* address p) to get address of data reference by location p
- array_of_types stores datatypes of elements
- Need to call MPI_Type_commit(MPI_Datatype* new_type_p) before using in communication function (just like any other MPI type)
- Need to MPI_Type_free (MPI_Datatype *new_type_p) when done

Parallel Performance

Collecting Timing Info

- Surround code we care about with timing collection function calls and subtract to find total time code ran
- double MPI_Wtime(void)
 - Returns wall clock time
- To ensure all parallel processes start at same time, have every process call MPI_Barrier (MPI_Comm comm) before calling MPI_Wtime so they start in synch
- Need to then find the largest time for the parallel processors as that is the time required by the entire set of parallel processes

Limitations to Performance Improvements

• Amdahl's Law

 $T_{enhanced} = (1 - fraction_{enhanced}) \times T_{unenhanced} + (\frac{fraction_{enhanced}}{Speedup_{ehnhancement}} \times T_{unenhanced})$

- Inherently sequential parts of code
- Overhead in parallel parts of code
 - Communication of data between processes
 - Load balancing
 - Synchronization

Speedup and Efficiency

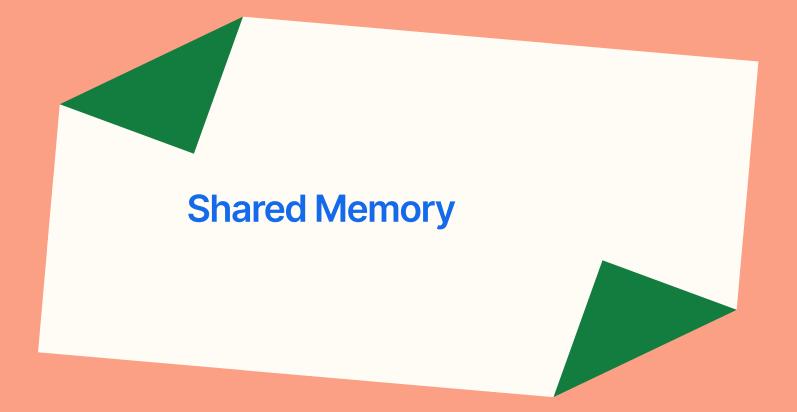
• Speedup(n, p) = $T_{\text{serial(n)}} / T_{\text{parallel(n,p)}}$

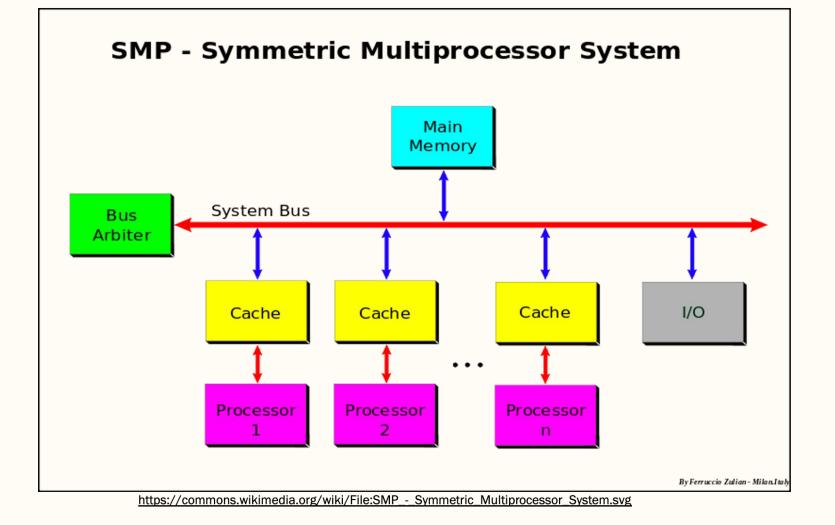
• Linear speedup = p

• Efficiency (n, p) = S(n,p) / p = $T_{\text{serial}(n)}$ / (p * $T_{\text{parallel}(n,p)}$)

Linear speedup corresponds to Efficiency = p/p = 1

- Strongly scalable maintain constant efficiency w/out increasing problem size
- Weakly scalable maintain constant efficiency if problem size increases at same rate as number of processes





Shared Memory Programming (aka Multithreaded Programming)

Processes

- Instance of running program One process per instance
 - Memory is private to process
- Associated resources (e.g. files)
- Access permissions
- Includes state for at least one thread of execution
 - HW registers, 0 runnable/blocked

Threads

- Contained within a process 1+ threads per process All memory shared among threads Each thread has own stack
- memory All resources and permissions defined by process Each thread has own state of
- execution
 - HW registers, runnable/blocked 0

POSIX Threads (aka Pthreads)

- Standard for Unix-like OSes
- Specifies API for multithreaded programming
- Like MPI, just a library linked with C programs

Basic Pthreads program setup

- #include <pthread.h>
- When compiling, link in pthreads library

gcc -g -o hello hello.c -lpthread

• When running, just run as normal

./hello

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
int num threads;
void *Hello(void *rank)
 printf("Hello from thread %ld out of %d\n", (long)rank, num threads);
 return NULL;
int main(int argc, char *argv[])
{
 long pthread;
 num threads= 4;
 pthread t ids[num threads];
 for(long i = 0; i < num threads; i++) {</pre>
    pthread create(&ids[i], NULL, Hello, (void*)i);
  }
 printf("Hello from main!\n");
 for(int i = 0; i < num threads; i++) {</pre>
    pthread join(ids[i], NULL);
  }
 return 0;
```

Some key features

- All non-stack variables shared by all threads
 - Where the challenges of multithreaded programming comes from!
- Variables allocated on a thread's stack are considered private
 - Technically, they can be accessed by other threads, but that's not good practice
- Running the program multiple times may result in different outcomes

Creating a Thread

int pthread_create(pthread_t* thread_p, const
 pthread_attr_t* attr_p, void* (*start_routine)(void*),
 void* arg_p)

- pthread_t* :
 - Stores thread-specific info that uniquely identifies thread
 - User program cannot access the contents
- pthread_attr_p* :
 - Specifies new threads attributes
 - Set stack, set stack size, set scheduling policy, set affinity, etc.
 - NULL indicates use of default values

Creating a Thread

int pthread_create(pthread_t* thread_p, const
 pthread_attr_t* attr_p, void* (*start_routine)(void*),
 void* arg_p)

- void* (*start_routine)(void*) :
 - Function pointer for code thread should execute
 - Function must take single void* argument and return void *
- void* arg_p :
 - Arguments to be passed to function
 - Can be single value cast to void*
 - Often is array or other complex data structure that contains 1+ values

Waiting for a Thread

int pthread_join(pthread_t thread, void **ret_val_p)

- Waits for thread specified to finish execution
 - If multiple threads call join on same thread, undefined action
 - Unjoined threads are considered to be zombie threads
- void** ret_val_p:
 - Pointer to location of the item returned by thread's return statement
 - Can be NULL