```
Why are statically-typed OOL's so inflexible?
```

```
Java programs require lots of type casts (as do C++, Object Pascal, etc.).
Why?
```

The Object class in Java illustrates most of the problems.

```
public class Object{
   protected Object clone(){...}
   public boolean equal(Object other){...}
}
```

Recall not allowed to change signature of methods in extensions.

Also all classes automatically inherit from Object.

```
public class A implements Cloneable{
  protected B b = ...;
  •••
  public Object clone(){
     A other = super.clone();
     other.b = b.clone();
                                           // type error
     return other;
  }
}
A = new A();
A = a1.clone();
                                      // type error
Both errors would disappear if could change
  return types of methods. (OK in C++)
Suppose also want to override equals in A:
  public boolean equal(Object other){
     if (other instanceof A) {
```

```
A aother = (A)other;
return (b = aother.b);
else
???
```

Problems:

- Inconvenient (and slow) to have run-time test No static check possible.
- What to do in else clause?
 If raise exception must declare in method header.

Get similar problems if try to define doubly-linked node as subclass of linked node.

```
public void setNext(Node newNext){...}
```

Want parameter type to be doubly-linked in subclass.

Methods where parameters should be of same type (class) as receiver called "binary methods".

Source of many typing problems in OOL's.

Still other problems:

```
public class Circle{
    protected Point center;
    ...
}
```

If define ColorCircle as extension, might want center to be ColorPoint.

Can't make any of these changes in signature (types) in Java, Object Pascal, and C++ (aside from return type in C++)

Why are there these restrictions?

Is there any way to overcome them?

Look at the following example:

```
class A{
   D m(C c) {...}
   void n() {... self.m(someC) ...}
}
class B{
   D' m(C' c') {...}
}
```

For which C', D' will B end up being type safe if A is?

Homework asks similar question for instance variables.

GJ: Adding F-bounded polymorphism to Java

Odersky, Wadler, et al (follow up to Pizza)

GJ adds parametric polymorphism w/ syntax like C++'s templates:

```
public class Stack<Elt> extends Vector<Elt>{
    public Elt push(Elt item){...}
    public Elt pop(){...}
    public Elt peek(){...}
    public boolean empty(){...}
    public int search(Elt o){...}
}
Stack<Point> myStack = new Stack<Point>();
Point aPoint = new Point(2,3);
myStack.push(aPoint);
```

Can also add constraints to type parameters:

```
public interface Comparing {
  public boolean equal(Comparing other);
  public boolean greaterThan(Comparing other);
 public boolean lessThan(Comparing other);
}
public class OrderedList<Elt implements Comparing>
                                 extends ... {
  protected Elt[] elts = new Elt[0];
  public void insert(Elt item){...
     while (elts[current].greaterThan(item))
       current ++;
          •••
     }
  }
  public Elt removeFirst(){...}
  public boolean empty(){...}
 public int searchFor(Elt o){...}
}
How to define ordered objects?
public class KeyedObj implements Comparing{
  protected int key ;
  . . . ;
  public int getKey(){...}
  public lessThan(Comparing other){
     return this.key < other.getKey();}</pre>
}
```

```
Won't work: other.getKey() not well-typed!
```

```
public lessThan(Comparing other){
```

```
if (other instanceof Comparing)
  return this.key < other.getKey();}
else
  ?????</pre>
```

```
Same problem as earlier!
```

```
F-bounded polymorphism (1989) can help:
public interface Comparing <Elt>{
  boolean lessThan(Elt other);
  boolean greaterThan(Elt other);
}
public class OrderedList
                < Elt implements Comparing<Elt> >{
  protected Elt [] elts;
  public void insert(Elt newVal){
    while (elts[current].greaterThan(newVal))
        current ++;
        •••
  }
}
public class KeyedObj
                 implements Comparing <KeyedObj>{
  protected int key ;
  . . . ;
  public int getKey(){...}
  public lessThan(KeyedObj other){
    return this.key < other.getKey();}</pre>
}
```

```
Now OrderedList<KeyedObj> is fine!
```

```
Generally works well (though confusing at first).
Still one problem -- F-bounded not preserved under subclass:
```

```
public class NuKeyedObj extends KeyedObj {
  protected String nuField;
   ...;
  public lessThan(KeyedObj other){
    return this.key < other.getKey() &&
        other.getNuField() ...;}
}</pre>
```

```
Unfortunately, NuKeyedObj does not implement Comparing <NuKeyedObj>.
```

Can't be used with OrderedList!

Other info on GJ:

- Works with existing JVM -- essentially translates to original code w/Object and casts.
- Authors designed so that existing library classes can be used as though they were polymorphic.
- Because of translation, cannot get accurate info using Java's reflection facilities or debugger.
- See GJ web page available through hmwk page.

Eiffel

Designed by Bertrand Meyer in mid-80's

Class-based OOL w/multiple inheritance

Assertions: pre- and post-conditions, loop invariants and variants built into language.

Supports bounded polymorphism.

Reference semantics like Java, garbage collection, etc.

Information hiding: private, public, or could list classes visible to (like C++'s friends)

No interfaces or modules.

In subclasses, can redefine or even rename methods.

Can also change type of instance variables, parameters and return types covariantly. Seen this can cause type-safety problems!

Introduced "anchor" types: Can declare type to be "like" another feature:

x: A; y: like x;

Current is Eiffel's name for self.

Example:

class LINKABLE [G]

feature

```
item: G; -- value held
right: like Current; -- Right neighbor
putRight (other: like Current) is
    -- Put `other' to right of current cell.
    do
        right := other
    ensure
        chained: right = other
end;
```

end -- class LINKABLE

```
class BILINKABLE [G] inherit
        LINKABLE [G]
           redefine
              putRight
           end
feature -- Access
  left: like Current; -- Left neighbor
  putRight (other: like Current) is
         -- Put `other' to right of current cell.
      do
         right := other;
         if (other /= Void) then
            other.simplePutLeft (Current)
         end
       end;
    putLeft (other: like Current) is
           -- Put `other' to left of current cell.
        do
           left := other;
           if (other /= Void) then
              other.simplePutRight (Current)
           end
        ensure
           chained: left = other
        end;
feature {BILINKABLE}
   simplePutRight (other: like Current) is
         -- set `right' to `other'
```

```
do
          right := other
       end;
   simplePutLeft (other: like Current) is
          -- set `left' to `other'
      do
          left := other
       end;
invariant
   rightSymmetry:
       (right /= Void) implies (right.left = Current);
   leftSymmetry:
       (left /= Void) implies (left.right = Current)
end -- class BILINKABLE
Notice BILINKABLE is subclass of LINKABLE.
Can't do this with Java, C++, etc.
Secret is use of "like Current" as type of instance variables and in types
  of methods.
Can define:
   class LINKEDLIST[NODE -> LINKABLE] ...
Can be instantiated with either

    LINKABLE (and get singly-linked list) or

• BILINKABLE (and get doubly-linked list).
Very expressive w/out F-bounded polymorphism:
deferred class Comparing
```

feature

lessThan(other: like Current): boolean

is deferred

end

```
greaterThan(other: like Current): boolean;
end
```

Unfortunately, use of like Current gives rise to implict covariant change to types of instance variables and method parameter and return types.

Thus BILINKABLE is not a subtype of LINKABLE. Though BILINKABLE is internally consistent.