

CSCI 136

Data Structures & Advanced Programming

Lecture 20
Spring 2018
Profs Bill & Jon

Last Time

- Iterators Recap
- Iterating over Iterators

Today

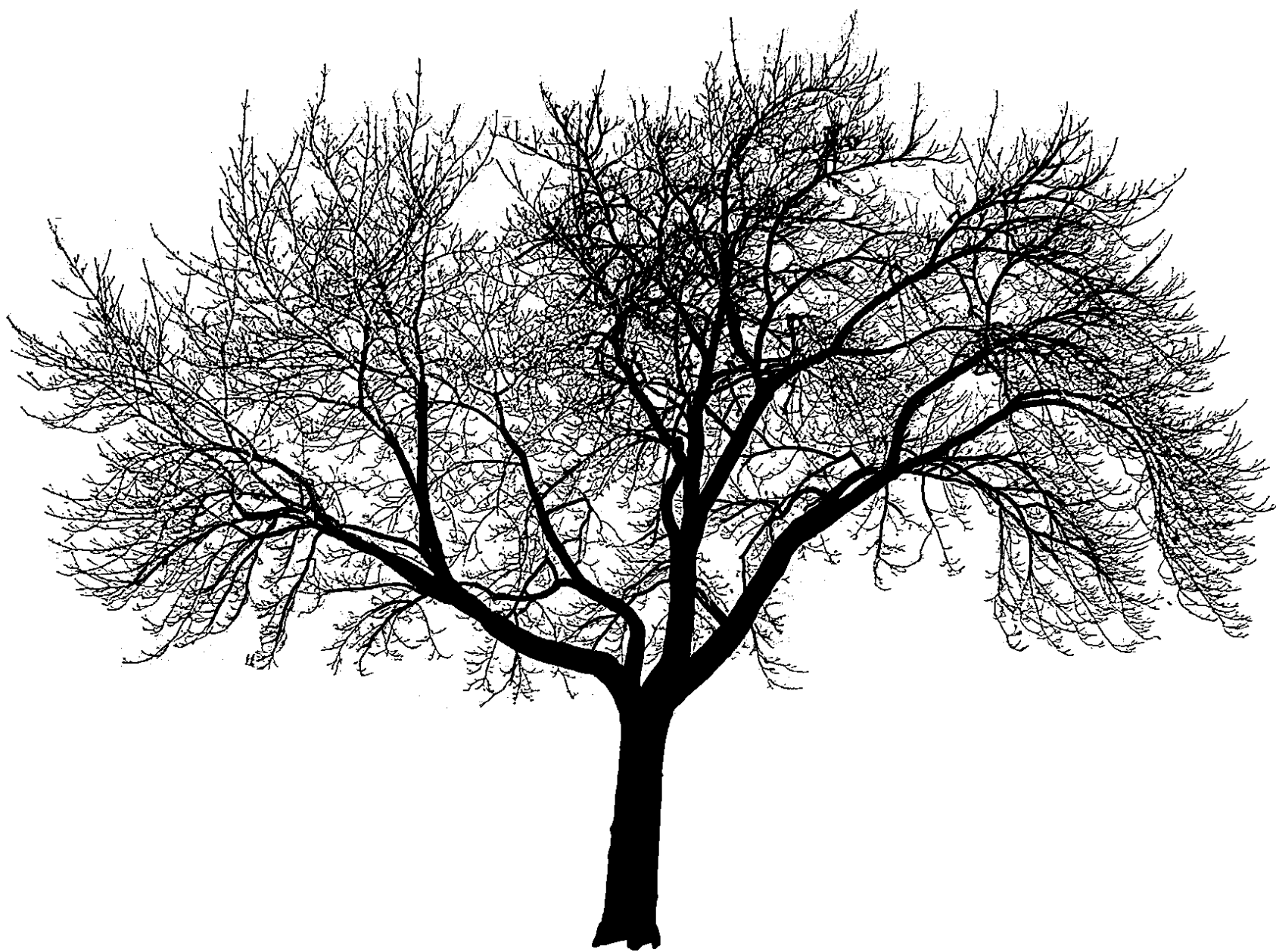
- Trees!
 - General Idea and Uses
 - Terminology
 - Some examples
 - Expression trees
 - Introduction to `structure5 BinaryTree` class
- `BinaryTree` class implementation details
- Proofs and theory
- Traversing trees

Introducing Trees

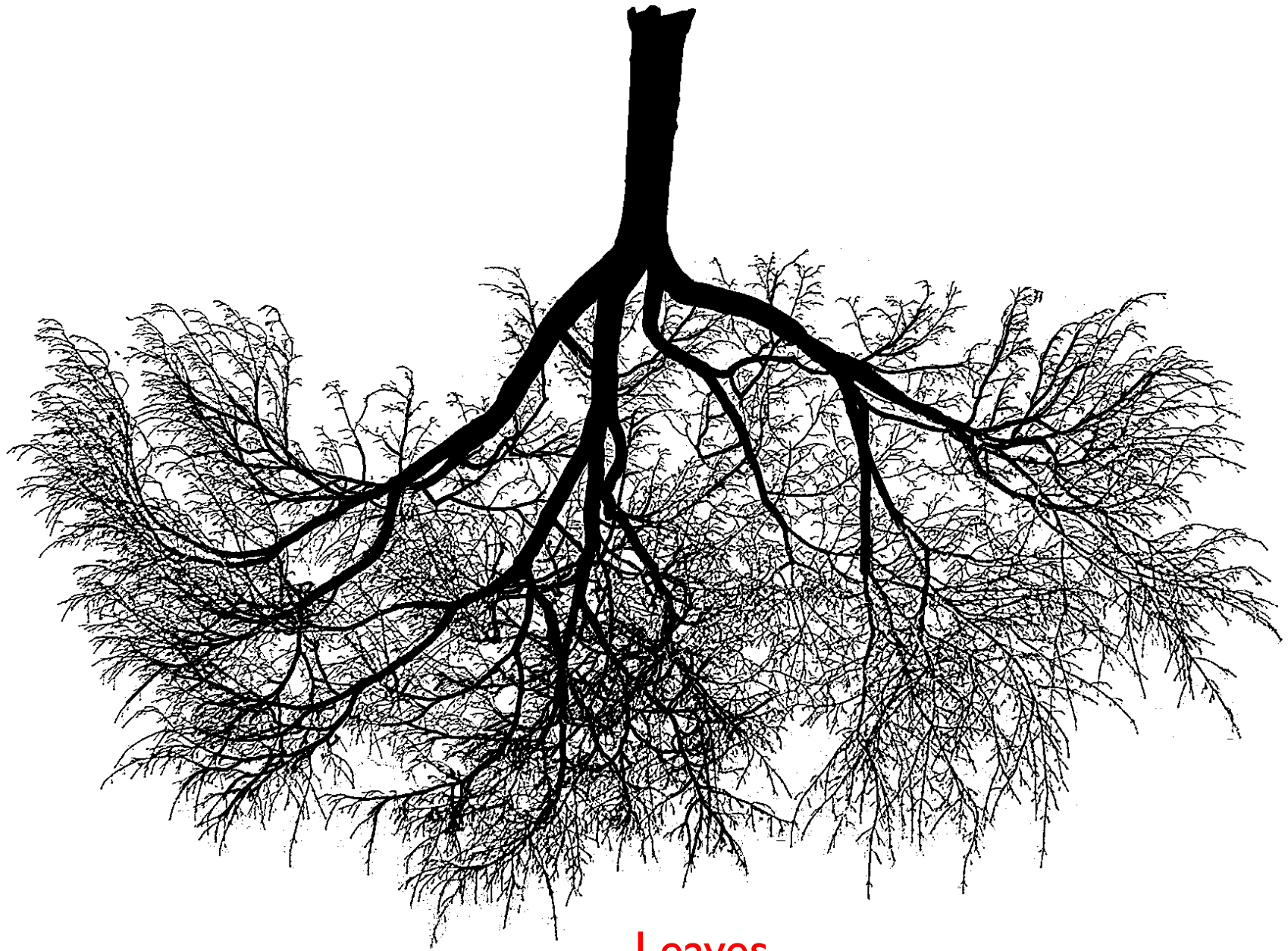
- Our structures have had a linear organization
 - Stacks, queues
 - Even ordered vectors, ordered lists, arrays, vectors, lists are visualized linearly
- By linear we essentially mean that each element has at most one successor and at most one predecessor...

Branching Out: Trees

- A tree is a data structure where elements can have multiple successors (called children)
- But still only one predecessor (called parent)



Root

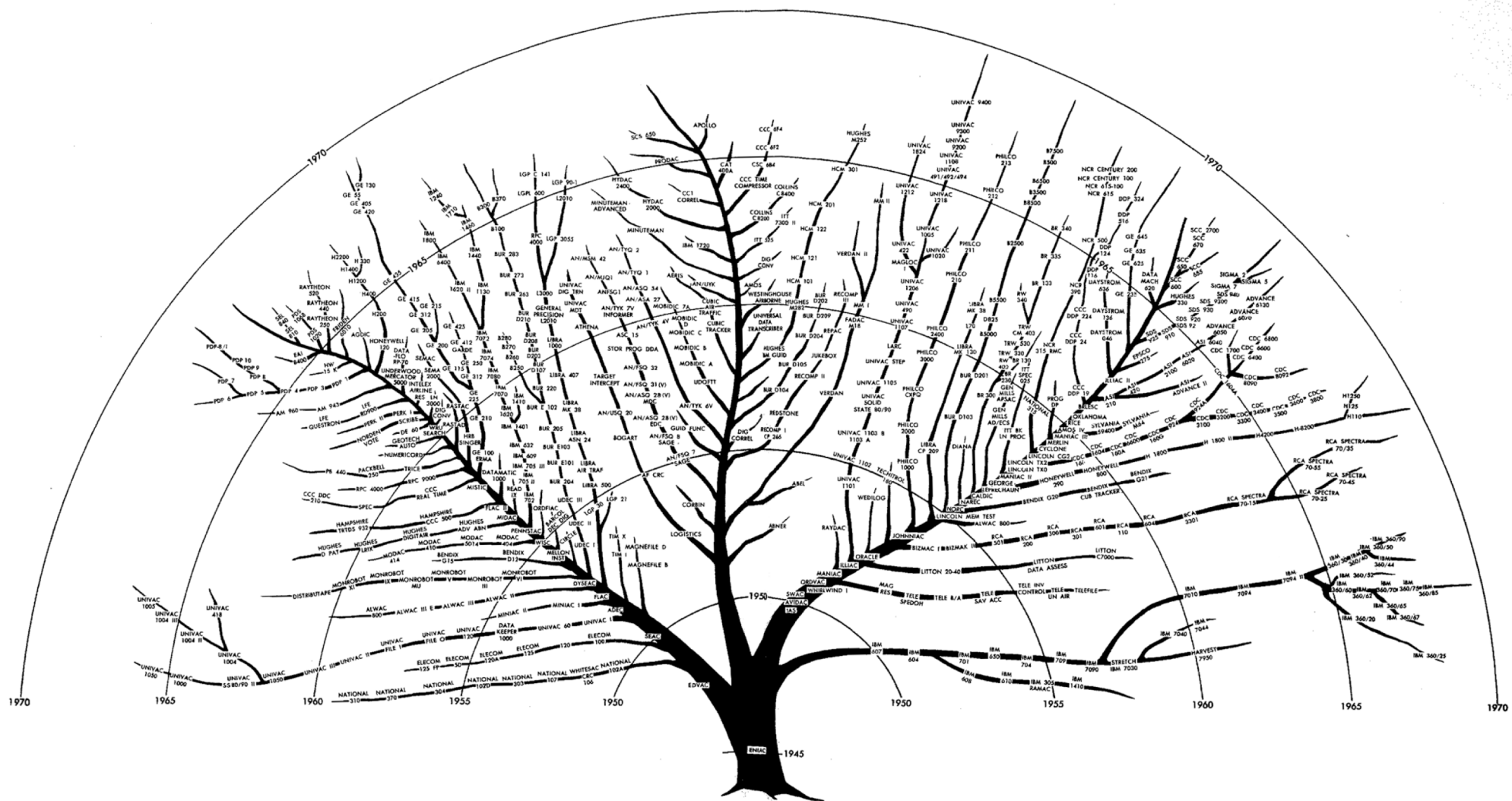


Leaves

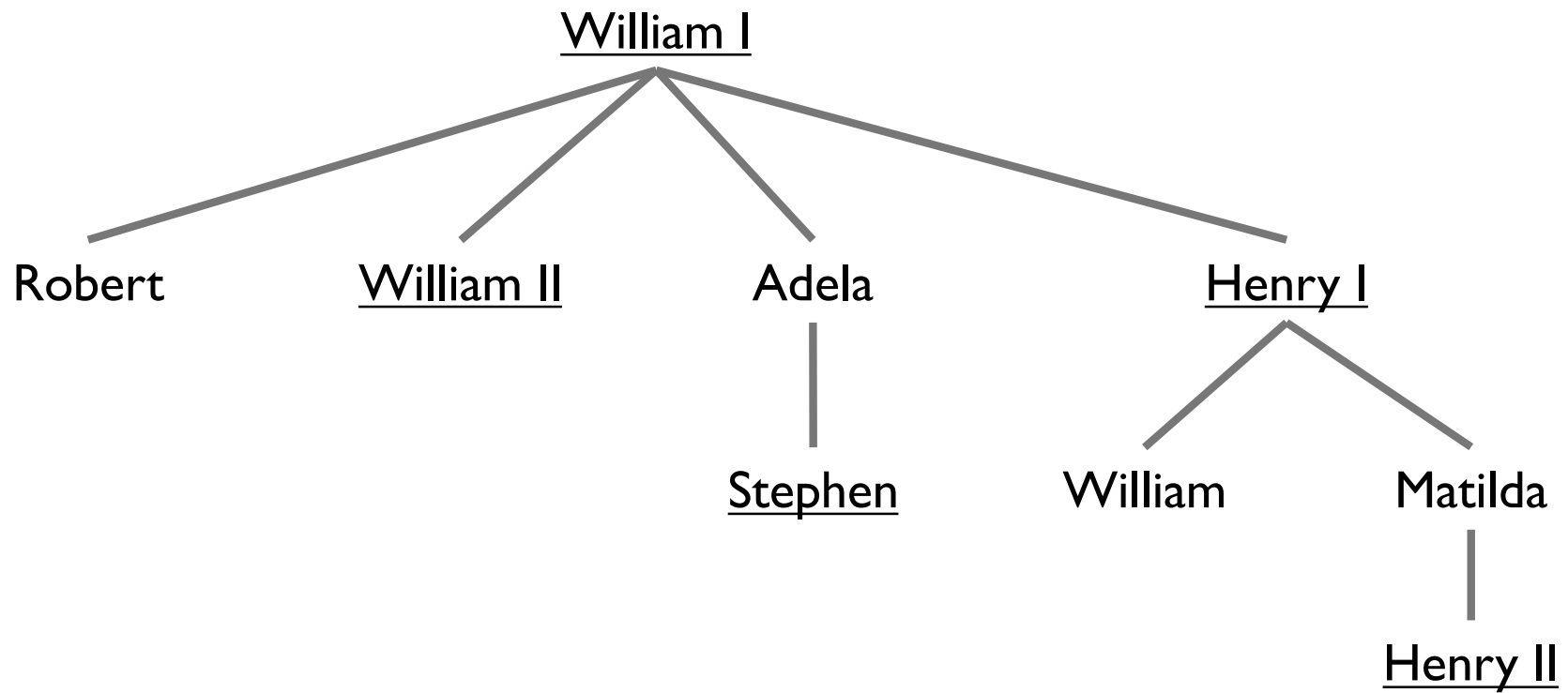


Tree Logic (Natalie Jeremijenko) at Mass MoCA

"Computer Tree"



House of Normandy, Battle of Hastings, 1066

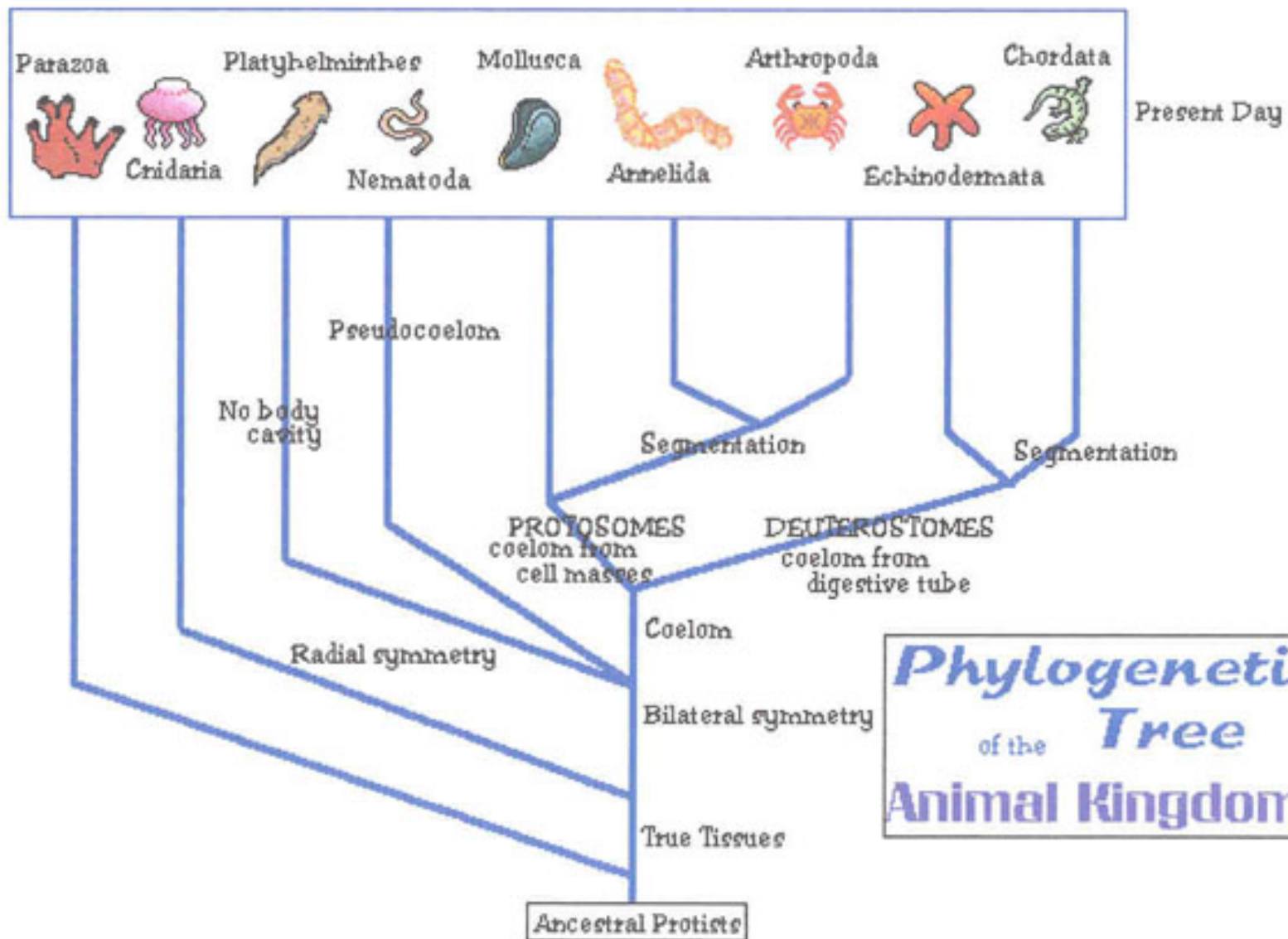


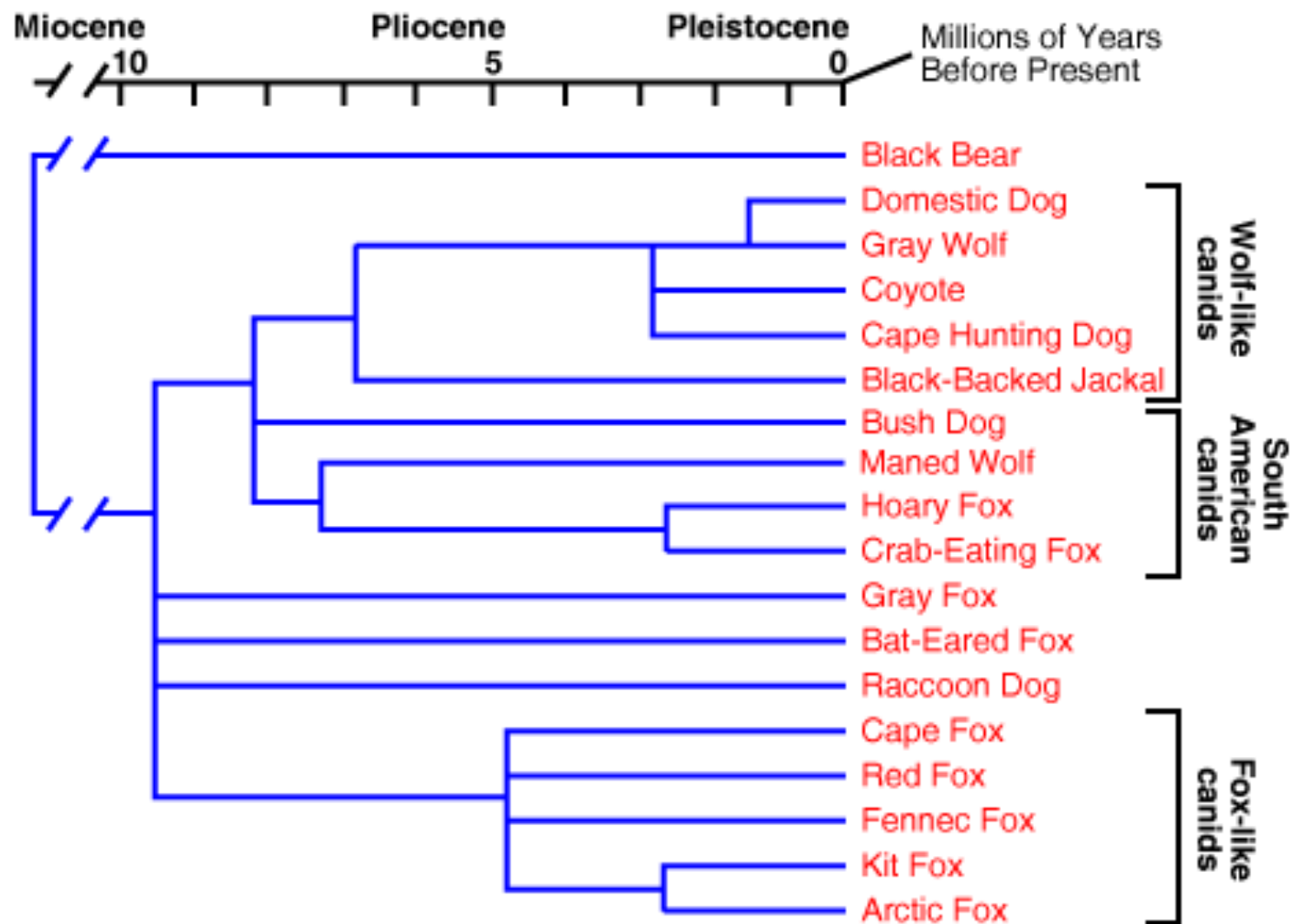
Tree Features

- Trees express hierarchical relationships
 - Directed: root to leaf
- **Root** at the top
- **Leaf** at the bottom
- **Interior nodes** in middle
- Parent, children, ancestors, descendants, siblings
- **Degree (of node)**: number of children of node
- **Degree (of tree)**: maximum degree (across all nodes)
- **Depth** of node: number of edges from root to node
- **Height**: maximum depth (across all nodes)

Other Trees

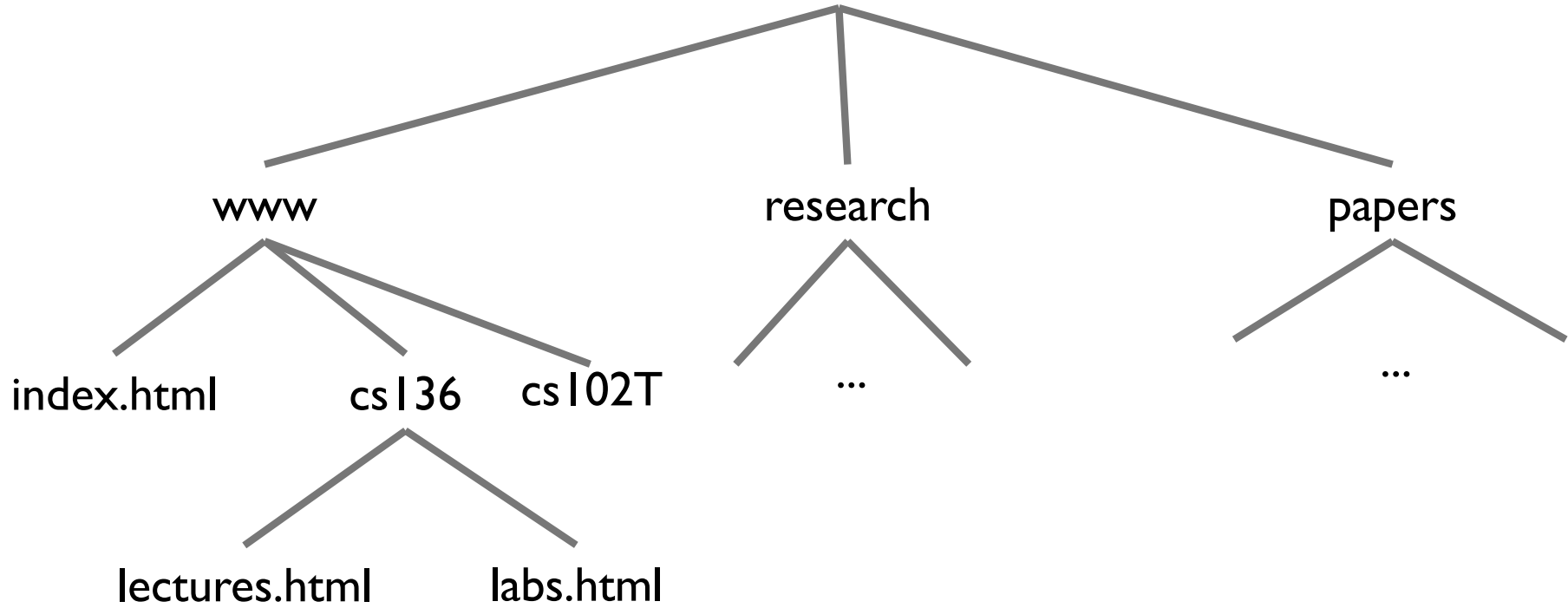
- Phylogenetic tree
- Directories of files
- Game trees
 - Build a tree
 - Search it for moves with high likelihood of winning
- Expression trees

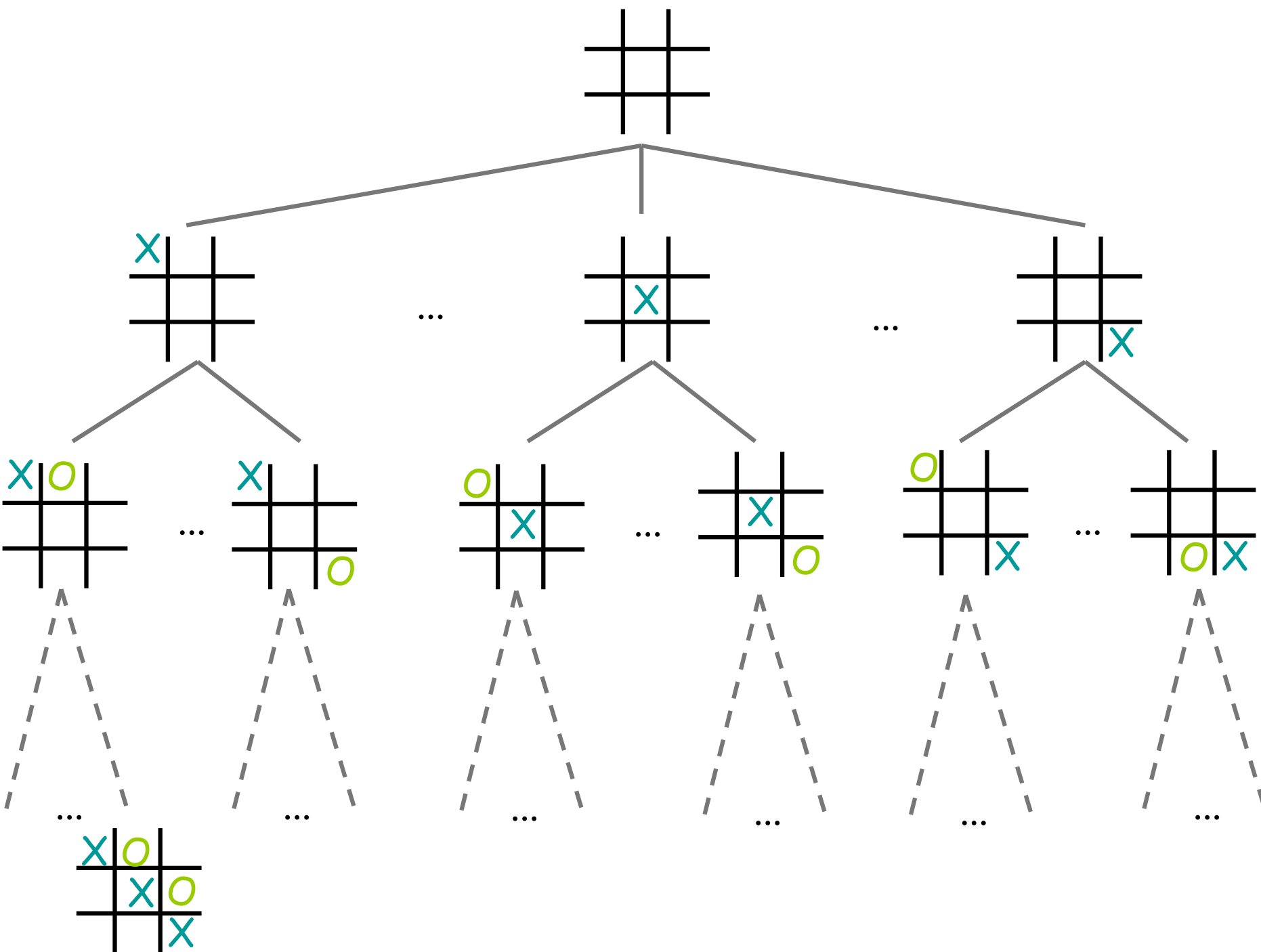






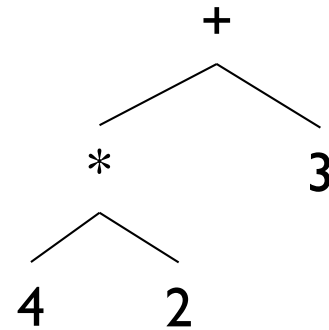
~jannen



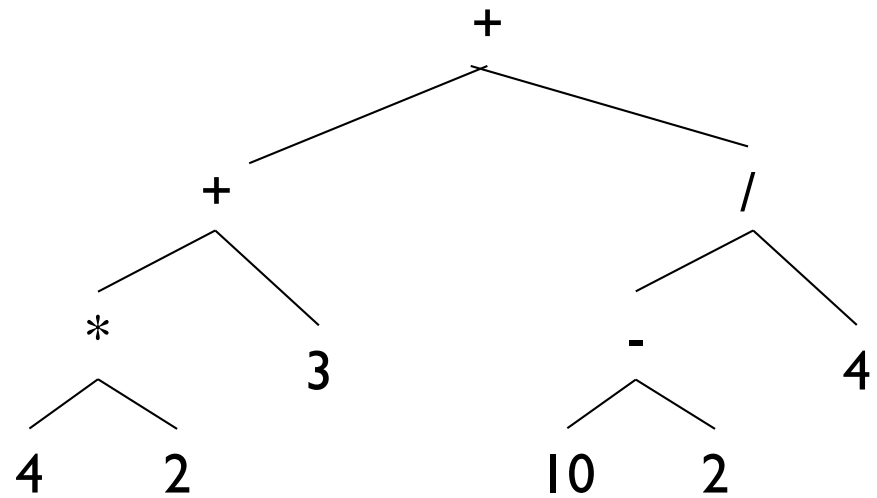


Expression Trees

$4 * 2 + 3$



$(4 * 2 + 3) + ((10 - 2) / 4)$



Introducing Binary Trees

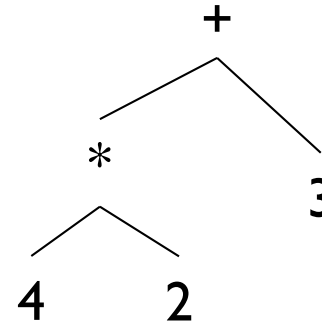
- **Degree** of each node ≤ 2
- Recursively defined. A tree can either be:
 - Empty
 - Root with left and right subtrees
- SLL: Recursive nature was captured by nodes (Node<E>) on inside
- Binary Tree: No “inner” node class; single BinaryTree class does it all
- (Not part of the structure hierarchy)

Binary Trees for (Math) Expressions

- General strategy
 - Make a binary tree (BT) for each leaf node
 - Move from bottom to top, creating BTs
 - Eventually reach the root
 - Call “evaluate” on final BT
- Example
 - How do we make a binary expression tree for: $(4*2)+3$
 - Leaves are numbers
 - Non-leaf nodes are operators
 - We will apply each operator to its children (ex: left + right)

Example: Expression Trees

$4 * 2 + 3$



Build using constructor

```
new BinaryTree<E>(value, leftSubTree, rightSubTree)
```

```
BinaryTree<String> fourTimesTwo =
```

```
    new BinaryTree<String>("*",  
        new BinaryTree<String>("4"),  
        new BinaryTree<String>("2"));
```

```
BinaryTree<String> fourTimesTwoPlusThree =
```

```
    new BinaryTree<String>("+",  
        fourTimesTwo,  
        new BinaryTree<String>("3"));
```

Evaluating Expression Trees

- Starting at the root,
 - Evaluate left subtree
 - Evaluate right subtree
 - Perform operation (+, -, *, /) with left and right

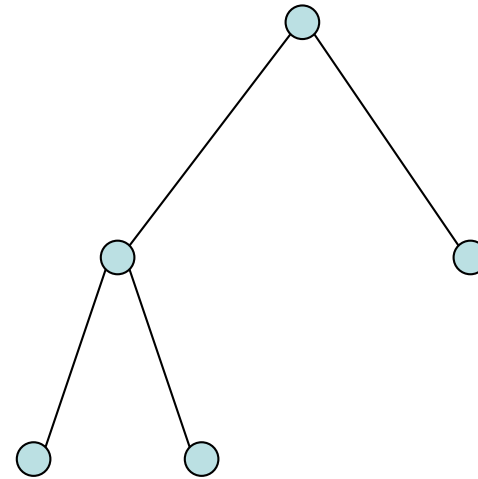
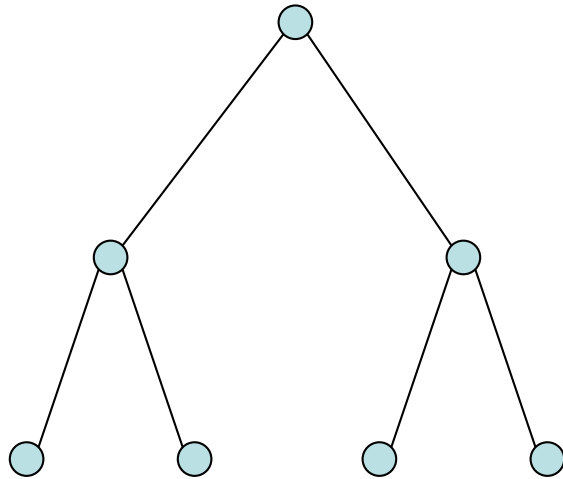
```
int evaluate(BinaryTree<String> expr) {
    if (expr.height() == 0) {
        return Integer.parseInt(expr.value());
    } else {
        int left = evaluate(expr.left());
        int right = evaluate(expr.right());
        String op = expr.value();
        switch (op) {
            case "+" : return left + right;
            case "-" : return left - right;
            case "*" : return left * right;
            case "/" : return left / right;
        }
        Assert.fail("Bad op");
        return -1;
    }
}
```

More Tree Terminology

- Some of the terminology is non-standard
- We will try to be consistent in this class, but...
 - We want to be able to communicate to our friends outside of Williams CS too!
- I *hate* jargon, but having a language for our data structures gives us the ability to express ideas and describe algorithms

Full vs. Complete (non-standard!)

- **Full** tree – A full binary tree of height h has *leaves only* on level h , and each internal node has exactly 2 children.
- **Complete** tree – A *complete* binary tree of height h is *full* to height $h-1$ and has all leaves at level h in leftmost locations.



All full trees are complete, but not all complete trees are full!