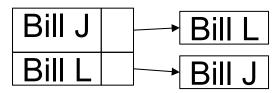
CSCI 136 Data Structures & Advanced Programming

Lecture 30

Fall 2017

Instructors:



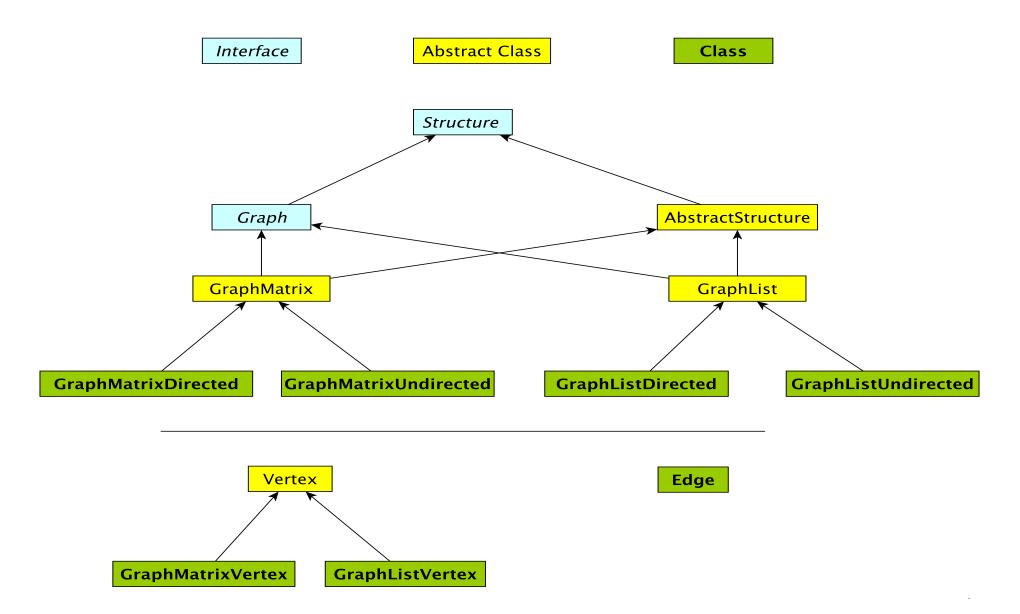
Last Time

- Graph Interface
 - Adjacency Array Implementation Basic Concepts
 - Adjacency List Implementation Basic Concepts
- Structure5 Graph classes + hierarchy

Today's Outline

- Graph Data Structures: Implementation
 - Adjacency Array Implementation Details
- Greedy Algorithms for Optimization
- Lab II: Exam Scheduling
 - Defining the problem
 - Sketching a design

Graph Classes in structure5



Graph Classes in structure5

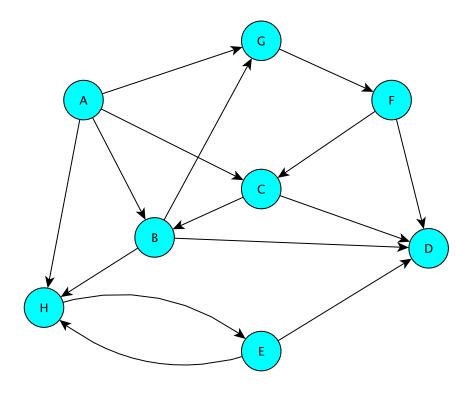
Why so many?!

- There are two types of graphs: undirected & directed
- There are two implementations: arrays and lists
- Strategy: implement as much code as can be written without assuming directedness
 - (Un)Directed Subclasses implement the rest

We'll tackle array-based graphs first....

Adjacency Array: Directed Graph

	Α	В	С	D	Ε	F	G	Н
Α	0	I	I	0	0	0	ı	I
В	0	0	0	I	0	0	I	I
С	0	I	0	I	0	0	0	0
D	0	0	0	0	0	0	0	0
Ε	0	0	0	I	0	0	0	I
F	0	0	I	ı	0	0	0	0
G	0	0	0	0	0	I	0	0
Н	0	0	0	0	I	0	0	0



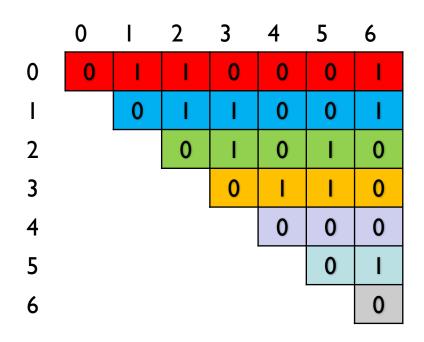
Challenges to having our rows/columns be "vertices"

- Can't use Objects as array indices
- •How does adding/deleting a vertex work?!

Adjacency Array: Undirected Graph

Halving the Space (not in structure5)

	0	I	2	3	4	5	6
0	0	1	1	0	0	0	1
ı	ı	0	I	1	0	0	1
2	ı	I	0	Τ	0	_	0
3	0	I	I	0	_	_	0
4	0	0	0	I	0	0	0
5	0	0	I	I	0	0	I
6	I	I	0	0	0	I	0



in general case: (i,j) maps to i* |V|+7

Vertex and GraphMatrixVertex

- We need to define a Vertex class
 - Unlike the Edge class, Vertex class is not public
 - Useful Vertex methods:

```
V label(), boolean visit(),
boolean isVisited(), void reset()
```

- GraphMatrixVertex class adds one more useful attribute to Vertex class
 - Index of node (int) in adjacency matrix int index()
 - Why do we only need one int to represent index?

Choosing a Dictionary Structure

- We need a structure that will let us retrieve the index of a vertex given the vertex label (a dictionary)
- Many choices
 - Vector of Associations:
 - Vector<Association<V, GraphMatrixVertex<V>>>
 - OrderedVector of Associations
 - BinarySearchTree of Associations
- Problem: We don't want to allow multiple vertices with same label.... [Why?]
- We'll use the Map Interface [Chapter 15]
 - Maps require a unique key for each entry

Digression: Map Interface

- Maps unique keys to values (V is value not vertex!!!)
- Methods for Map<K, V>
 - int size() returns number of entries in map
 - boolean isEmpty() true iff there are no entries
 - boolean containsKey(K key) true iff key exists in map
 - boolean containsValue(V val) true iff val exists at least once in map
 - V get(K key) get value associated with key
 - V put(K key, V val) insert mapping from key to val, returns value replaced (old value) or null
 - V remove(K key) remove mapping from key to val
 - void clear() remove all entries from map
- We'll study this more in a week or so....

Implementing the Matrix Model

Abstract class – partially implements Graph

```
public abstract class GraphMatrix<V,E> implements Graph<V,E>
```

- This class will implement features common to directed and undirected graphs
- Instance variables

```
protected int size; //max size of matrix
protected Object data[][]; //matrix of edges
protected Map<V, GMV<V>> dict; //labels -> vertices
// This is structure5.Map, NOT java.util.Map!
protected List<Integer> freeList; //avail indices
protected boolean directed;
```

GraphMatrix Constructor

(Yes, abstract classes can have constructors!)

```
protected GraphMatrix(int size, boolean dir) {
   this.size = size; // set maximum size
   directed = dir; // fix direction of edges
   // the following constructs a size x size matrix
   // (the "Objects" will be "Edges")
   // (can't use generics with arrays!)
   data = new Object[size][size];
   // label → index translation table
   dict = new Hashtable<V,GraphMatrixVertex<V>>(size);
   // put all indices in the free list
   freeList = new SinglyLinkedList<Integer>();
   for (int row = size-1; row \geq 0; row--)
       freeList.add(new Integer(row));
}
```

GraphMatrix add()

```
public void add(V label) {
    // if there already, do nothing
    if (dict.containsKey(label)) return;

Assert.pre(!freeList.isEmpty(), "Matrix not full");
    // allocate a free row and column
    int row = freeList.removeFirst().intValue();
    // add vertex to dictionary
    dict.put(label, new GraphMatrixVertex<V>(label, row));
```

GraphMatrix remove()

```
public V remove(V label) {
       // find and extract vertex
       GraphMatrixVertex<V> vert;
       vert = dict.remove(label);
       if (vert == null) return null;
       // remove vertex from matrix
       int index = vert.index();
       // clear row and column entries
       for (int row=0; row<size; row++) {</pre>
           data[row][index] = null;
           data[index][row] = null;
       // add node index to free list
       freeList.add(new Integer(index));
       return vert.label();
}
```

Neighbors Iterator: GraphMatrix

neighbors Iterator

```
public Iterator<V> neighbors(V label) {
      GraphMatrixVertex<V> vert = dict.get(label);
      List<V> list = new SinglyLinkedList<V>();
       for (int row=size-1; row>=0; row--) {
             Edge<V,E> e = (Edge<V,E>)data[vert.index()][row];
             if (e != null)
                    if (e.here().equals(vert.label()))
                           list.add(e.there());
                           else list.add(e.here());
      return list.iterator();
   }
```

- Completes the implementation of GraphMatrix to ensure graph is directed
- GraphMatrixUndirected is very similar...
- How do we implement GraphMatrixDirected?
 - We'll discuss some methods
 - Read Ch 16 for complete details...

Constructor

```
public GraphMatrixDirected(int size) {
    // pre: size > 0
    // post: constructs an empty graph that may be
    // expanded to at most size vertices. Graph
    // is directed if dir true and undirected
    // otherwise

// call GraphMatrix constructor
    super(size,true);
}
```

addEdge

removeEdge

```
// pre: vLabel1 and vLabel2 are labels of existing vertices
public E removeEdge(V vLabel1, Vlabel2) {
    // get indices
    int row = dict.get(vLabel1).index();
    int col = dict.get(vLabel2).index();
    // cache old value
    Edge<V,E> e = (Edge<V,E>)data[row][col];
    // update matrix
    data[row][col] = null;
    if (e == null) return null;
    else return e.label(); // return old value
}
```

GraphMatrix Efficiency

- Assume Map operations are O(I) (for now)
 - |E| = number of edges
 - |V| = number of vertices
- Runtime of add, addEdge, getEdge, removeEdge, remove?
- Space usage?
- Conclusions
 - Matrix is good for dense graphs
 - Have to commit to maximum # of vertices in advance

Efficiency: Assuming Fast Map

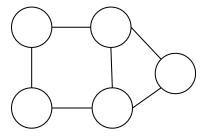
	GraphMatrix		
add	O(I)		
addEdge	O(I)		
getEdge	O(I)		
removeEdge	O(I)		
remove	O(V)		
space	O(V ²)		

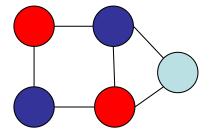
Lab II Overview: Graph Algorithms using structure5

Greedy Algorithms

- A greedy algorithm attempts to find a globally optimum solution to a problem by making locally optimum (greedy) choices
- Example: Walking in Manhattan
- Example: Graph Coloring
 - A (proper) coloring of a graph G=(V,E) is an assignment of a value (color) to each vertex so that adjacent vertices get different values (colors)
 - Typically one strives to minimize the number of colors used

Graph Coloring Example





Greedy Coloring: Math

```
Here's a greedy coloring algorithm
```

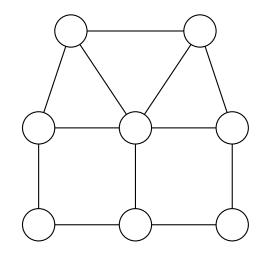
```
Build a collection C = \{C_1, ..., C_k\} of sets of vertices
i = 0; C_i = \{\} //empty set
while G is has more vertices
       for each vertex u in G
               if u is not adjacent to any vertex of C_i
                       remove u from G and add u to C_i
       add C_i to C
       i++:
```

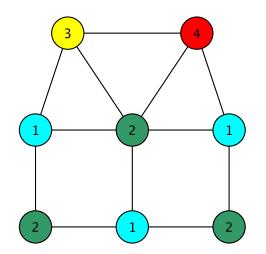
Return C as the coloring

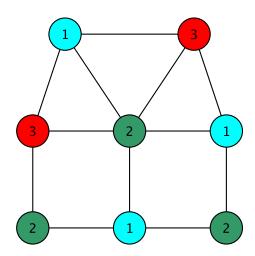
Greedy Coloring: CS

```
Here's a greedy coloring algorithm
Create a structure C to hold a collection of lists
while G is not empty
      pick a vertex v in G; create an empty list L; add v to L
      for each vertex u \neq v in G
              if u is not adjacent to any vertex of L
                     add u to L
       remove all vertices of L from G
       add L to C
Return C as the coloring
```

Greedy Coloring







Greedy Coloring

Some observations

- Each list (color class) L is a set of vertices no two of which are adjacent (an independent set)
- Each color class is maximal: cannot be made any larger
 - The hope is that this results in fewer colors being needed
 - But the solution is not always optimum!
 - This is a very hard problem
- The coloring problem is the same as finding a partition of the vertex set into independent sets
 - Partition means union of disjoint sets

Lab II: Exam Scheduling

Find a schedule (set of time slots) for exams so that

- No student has two exams in the same slot
- Every course is in a slot
- The number of slots is as small as possible

This is just the graph coloring problem in disguise!

- Each course is a vertex
- Two vertices are adjacent if the courses share students
- A slot must be an independent set of vertices (that is, a color class)

Lab 11 Notes: Using Graphs

- Create a new graph in structure5
 - GraphListDirected, GraphListUndirected,
 - GraphMatrixDirected, GraphMatrixUndirected
- Graph<V,E> conflictGraph = new GraphListUndirected<V,E>();

Lab II: Useful Graph Methods

- void add(V label)
 - add vertex to graph
- void addEdge(V vtx1, V vtx2, E label)
 - add edge between vtx1 and vtx2
- Iterator<V> neighbors(V vtx1)
 - Get iterator for all neighbors to vtx l
- boolean isEmpty()
 - Returns true iff graph is empty
- Iterator<V> iterator()
 - Get vertex iterator
- V remove(V label)
 - Remove a vertex from the graph
- E removeEdge(V vLabel1, V vLabel2)
 - Remove an edge from graph