

Fire OS



Amazon's operating system for Kindle Fire

Kindle Fire History and Overview

- Kindle fire is Amazon's version of the iPad
- It provides access to movies, books, music, social media, and work-related content.

- 2004: Amazon started selling e-books
- 2009: Apple launched the iPad, which contained access to iTunes and iBooks stores
- 2011: Jeff Bezos disclosed the "Otter project" and announced the launch of Kindle Fire



fire

Brilliant and responsive entertainment

Kindle Fire Overview

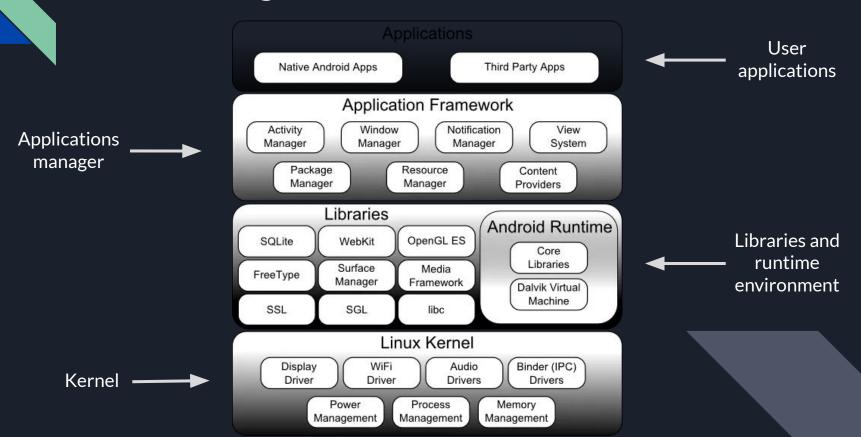
Hardware

- 7" multi-touch display with a 1024 x 600 pixel resolution capable of displaying 16 million colors at 169 pixels per inch
- 8GB of internal storage
- 2GB of this memory is reserved for the OS
- Texas Instruments OMAP 4430, a 1GHz dual-core processor

Software

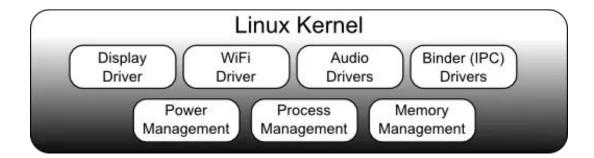
- Fire OS
- OS is based on Google's Android 2.3
 (Gingerbread) mobile operating system
- Highly customized to match low price sales (\$199) and reduced storage
- Customized in this case mostly means simplified

Diving into Fire OS Architecture



Kernel

- Based on Linux version 2.6
- Multitasking execution environment dual core
- Android applications do not run as processes directly on the Linux kernel, instead they
 run on within its own instance of the Dalvik VM
- Why? Applications are essentially sandboxed and cannot interfere with the OS
- Also enforces level of abstraction so that applications aren't tied to specific hardware
- Fun fact: the Dalvik executable (.dex) format has a 50% smaller memory footprint than standard Java bytecode



Resource Management

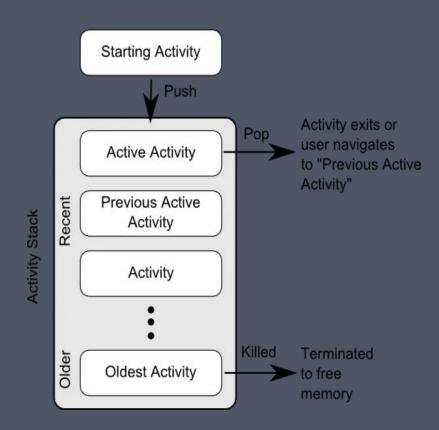
- Each running Android application is viewed by the OS as a separate process
- If resources on the device are reaching capacity, it will terminate processes to free up memory
- Importance hierarchy: considers both the priority and state of all running processes
- Processes are terminated from the lowest priority and working up the hierarchy until sufficient resources have been liberated

Process states



Activity Stack

- For each application that is running, the runtime system maintains an Activity Stack
- When an application is launched, the first of the application's activities to be started is placed onto the top of the stack
- When the active activity exits, it is popped off the stack and the activity located beneath it becomes the current active activity
- 4 activity states: Running, Paused, Stopped, Killed



thread_create()

```
static int kthread(void * create)
        /* Copy data: it's on kthread's stack */
        struct kthread_create_info *create = _create;
        int (*threadfn)(void *data) = create->threadfn;
        void *data = create->data:
        struct kthread self;
        int ret;
        self.should stop = 0:
        init completion(&self.exited);
        current->vfork_done = &self.exited;
        /* OK, tell user we're spawned, wait for stop or wakeup */
        __set_current_state(TASK_UNINTERRUPTIBLE);
        create->result = current;
        complete(&create->done);
        schedule();
        ret = -EINTR;
        if (!self.should stop)
                ret = threadfn(data);
        /* we can't just return, we must preserve "self" on stack */
        do_exit(ret);
```

```
void
create(void (*func)())
{
    //look through all threads
    for(int i = 0; i < NTHREAD; i++){
        //found free thread
        if(thread[i].state == FREE){
            thread[i].ra = (uint64)(func);
            thread[i].sp = (uint64)(thread[i].stack + TSSIZE);
            thread[i].state = RUNNABLE;
            break;
    }
}</pre>
```

```
struct task struct *kthread create(int (*threadfn)(void *data),
                                   void *data.
                                   const char namefmt[],
                                   ...)
       struct kthread create info create;
        create.threadfn = threadfn;
       create.data = data;
        init completion(&create.done);
        spin_lock(&kthread_create_lock);
        list_add_tail(&create.list, &kthread_create_list);
        spin unlock(&kthread create lock);
       wake_up_process(kthreadd_task);
       wait for completion(&create.done);
        if (!IS_ERR(create.result)) {
                struct sched_param param = { .sched_priority = 0 };
                va list args:
                va start(args, namefmt);
                vsnprintf(create.result->comm, sizeof(create.result->comm),
                          namefmt, args);
                va_end(args);
                * root may have changed our (kthreadd's) priority or CPU mask.
                * The kernel thread should not inherit these properties.
                sched setscheduler nocheck(create.result, SCHED NORMAL, &param);
                set cpus allowed ptr(create.result, cpu all mask);
        return create.result:
```

```
/**
* sys_sched_yield - yield the current processor to other threads.
 * This function yields the current CPU to other tasks. If there are no
* other threads running on this CPU then this function will return.
*/
SYSCALL_DEFINE0(sched_yield)
        struct rq *rq = this_rq_lock();
        schedstat_inc(rq, yld_count);
        current->sched_class->yield_task(rg);
        /*
         * Since we are going to call schedule() anyway, there's
         * no need to preempt or enable interrupts:
         release(rq->lock);
        spin_release(&rq->lock.dep_map, 1, _THIS_IP_);
        do_raw_spin_unlock(&rq->lock);
        preempt_enable_no_resched();
        schedule();
        return 0;
```

yield_thread()

```
void
yield(void)
{
  if(current->state != MONITOR) {
    current->state = RUNNABLE;
  }
  schedule();
}
```

Potentially Interesting Stuff

```
/*
* Schedules idle task to be the next runnable task on current CPU.
* It does so by boosting its priority to highest possible.
* Used by CPU offline code.
void sched_idle_next(void)
        int this_cpu = smp_processor_id();
        struct rq *rq = cpu_rq(this_cpu);
        struct task struct *p = rq->idle;
        unsigned long flags:
        /* cpu has to be offline */
        BUG ON(cpu online(this cpu));
        /*
        * Strictly not necessary since rest of the CPUs are stopped by now
        * and interrupts disabled on the current cpu.
         */
        raw_spin_lock_irgsave(&rg->lock, flags);
        __setscheduler(rq, p, SCHED_FIF0, MAX_RT_PRI0-1);
        activate_task(rq, p, 0);
        raw_spin_unlock_irqrestore(&rq->lock, flags);
```

Thank you!

Questions?

Some quality comments from the smart programmers at Amazon:

```
/*
 * Ahh, all good. It wasn't running, and it wasn't
 * runnable, which means that it will never become
 * running in the future either. We're all done!
 */
break;
```

```
/*
 * Was it really running after all now that we
 * checked with the proper locks actually held?
 *
 * Oops. Go back and try again..
 */
if (unlikely(running)) {
         cpu_relax();
         continue;
}
```