

## Computer Science 136

### Data Structures

#### Lecture #21 (November 10, 2021)

1. Questions?
2. Skew-heap implementation (**SkewHeap**).
  - (a) Notion: a **merge** of two heaps `h1` and `h2`.
    - i. If one heap is empty, use the other as the result.
    - ii. Otherwise, assume the `h1` root is smallest:
    - iii. Case 1: If `h1` has no left child: make `h2` its left.
    - iv. Case 2: Otherwise, swap the children of `h1`, and merge `h2` with `h1`'s new left (former right).
  - (b) Notice how the leftmost branch appears to be the target of all merges. But: at each stage, children are swapped/twisted. Result:
  - (c) Has amortized logarithmic cost even though the tree is not necessarily very balanced. Very cool analysis based on some clever bookkeeping/accounting tricks.
  - (d) **getFirst**: return root.
  - (e) **remove**: return root after merging children.
  - (f) **add**: merge new value with existing heap.

3. Binary Search Trees.

- (a) An implementation of an **OrderedStructure**: **add**, **remove**, **get**, **contains**, **iterator**.
- (b) Comparable values are kept in an (internal) binary tree.
- (c) All values to the left of the root are smaller *or equal*.
- (d) All values to the right of the root are larger.
- (e) We write a method **locate** that determines the correct location for the value in the tree. **Locate** can be used to determine if the tree contains a value, or to find the best location to insert it.
- (f) We have an important notion of the predecessor and successor of a node in a tree. The predecessor is the rightmost descendent of the left child. The successor is the opposite. Adding a right child to the predecessor installs a new predecessor, and *vice versa* for the left child of the successor.
- (g) **removeTop** is an important method that allows you to remove the top node of a (sub)tree. It has to be done with care. Several cases: study these.

- (h) The iterator is an **inorderIterator** on the root of the underlying binary tree.
- (i) Rotations.
  - i. Tree is balanced if, at each node, children have heights within 1.
  - ii. Left- and right- rotations fix problems of balance (See Figure 4.4). Rotations can be seen as bringing a node higher in the tree.

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#### Notes.