URGENT

Dear Computer Science Majors,

My musical In Transit is opening on Broadway tonight!

Just this morning I learned that my 90 minute show needs to be divided exactly into two 45 minute halves. (Couldn't they just *reschedule* that fire alarm test?!)

I have contacted everyone that I know and they have not been able to divide the shows by hand. Someone suggested that this could be a Computer Science problem — something about algorithms and complexity? — so I am turning to you.

As a Williams College alum I know that your education has been outstanding, so I know that you can do this.

The list of performances and their times is attached. The performances can be reordered however you like, but they must divide into two 45 minute halves.

Thank you so much! Go Ephs!

Sincerely, Kristen Anderson-Lopez



"Deep Beneath the City"	193
"Not There Yet"	131
"Do What I Do"	180
"Four Days Home"	277
"Broke"	345
"Saturday Night Obsession"	234
"Wingman"	454
"But, Ya Know"	264
"Not There Yet (Reprise)"	377
"Keep It Goin'"	178
"A Little Friendly Advice"	372
"Choosing Not to Know"	330
"The Moving Song"	299
"We Are Home"	434
"Getting There"	520
"Finale"	812
	5400 sec
	(90 min)

Objects Binary strings of length n

ObjectsBinary strings of length n



ObjectsBinary strings of length n

Operation[Complement one bit

Objects Binary strings of length n

Operation Complement one bit (Hamming distance one)

00010011

Objects Binary strings of length n

Operation Complement one bit (Hamming distance one)





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Operation Complement one bit (Hamming distance one)

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Objects Binary strings of length n

Operation Complement one bit (Hamming distance one) OOODODD before: 0000000

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Objects Binary strings of length n

Operation[Complement one bit

Lexicographic order 0 0 0 0 n = 4

Objects Binary strings of length n

Operation[Complement one bit

Gray code



Objects Binary strings of length n

Operation **Complement one bit**

Lexicographic order

0

0

NOT a Gray code

Objects Binary strings of length n

Operation[Complement one bit



n = 4



Objects Binary strings of length n

Operation[Complement one bit





Objects Binary strings of length n

Operation[Complement one bit



Lexicographic order



n = 4

Objects Binary strings of length n

Operation[Complement one bit



<u>;;;</u>

n = 4



ObjectsBinary strings of length n

Operation[Complement one bit

The Gray Code is <u>cyclic</u> if the last and first strings differ in one bit.



<u>;;;</u>

n = 4





A Gray Code has the following properties:

- Every n-bit binary string is listed exactly once.
- Consecutive strings differ in exactly one bit.



A Gray code for the 2-bit binary strings.



A Gray code for the 3-bit binary strings.

Activity

Arrange the 4-bit binary strings as a Gray code.

Questions

- Is your Gray code cyclic?
- How could you build a Gray code for any n?

The BRGC for n-bits uses two copies of the BRGC for (n-1)-bits:

Reflect the second copy.

Prefix 0 to each string in the first copy.

Prefix 1 to each string in the second copy.

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Ruler Sequence

The sequence of complemented bits in the BRGC for n-bits uses two copies of the sequence for (n-1)-bits:
The sequence of complemented bits in the BRGC for n-bits uses two copies of the sequence for (n-1)-bits:

Insert n between the two copies.

The sequence of complemented bits in the BRGC for n-bits uses two copies of the sequence for (n-1)-bits:

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n = 2

9

n = 3

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n = 2

9

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3

2

Applications

The binary reflected Gray codes has thousands of applications.

NP-Complete decision problem:

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There are n objects

NP-Complete decision problem:
— There are n objects
Object i has weight(i)

1. Bowie	3:15
2. Robots	3:43
3. Business Time	4:05
4. Hiphopopotamus	2:09
5. Hurt Feelings	2:38
6. Fashion is Danger	2:20
7. Carol Brown	3:26

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NP-Complete decision problem: Exhaustive computation:
— [There are n objects
— [Object i has weight(i)
Can the objects be partitioned
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Exhaustive computation: — There are 2ⁿ partitions

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0.60



Exhaustive computation: ——[There are 2ⁿ partitions —[O(1)-time with Gray code

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Exhaustive computation: — [There are 2ⁿ partitions — [O(1)-time with Gray code Total is O(2ⁿ)-time instead of O(n·2ⁿ)-time for naive algorithm.



ADATE .	YES ONO		т <u>е</u>	. 0	YES ONO
1. Bowie	3:15				
2. Robots	3:43				
3. Business Time	4:05				
4. Hiphopopotamus	3 2:09				
5. Hurt Feelings	2:38				
6. Fashion is Dang	er 2:20				4.4
7. Carol Brown	3:26				
8. Sugalumps	2:11				
Total	23:47	Total			0:00
0000	0	0	0	0	0
7. 6	. 5.	4.	3.	2.	1.

ADATE .	YES ONO		TE I.	. 0'	YES ONO
		1. Bo	wie		3:15
2. Robots	3:43				
3. Business Time	4:05				
4. Hiphopopotamus	2:09				
5. Hurt Feelings	2:38				
6. Fashion is Dange	er 2:20				
7. Carol Brown	3:26				
8. Sugalumps	2:11				
Total	20:32	Total			3:15
0000	0	0	0	0	
7. 6.	5.	4.	3.	2.	1.

ADATE .	YES ONO		TE 1.	. 01	YES ONO
		1. Bo	owie		3:15
		2. Ro	bots		3:43
3. Business Time	4:05				
4. Hiphopopotamus	2:09				
5. Hurt Feelings	2:38				
6. Fashion is Dange	r 2:20				
7. Carol Brown	3:26				
8. Sugalumps	2:11				
Total	13:49	Total			6:58
0000	0	0	0		
7. 6.	5.	4.	3.	2.	1.

ADATE .	YES ONO	B DAT	E	. 01	YES ONO
1. Bowie	3:15				
		2. Ro	bots		3:43
3. Business Time	4:05				
4. Hiphopopotamus	2:09				
5. Hurt Feelings	2:38		Route of the		
6. Fashion is Dange	r 2:20				
7. Carol Brown	3:26				
8. Sugalumps	2:11				
Total	17:04	Total		and area	3:43
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ADATE .	YES ONO	BATE	S ONO
1. Bowie	3:15		
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7. Carol Brown	3:26		
8. Sugalumps	2:11		
Total	12:59	Total	7:48
	1		
	$\mathbf{}$		\frown
	U		
7. 6.	5.	4. 3. 2.	1.



Volume



Volume











Which volume is selected?Encode each position with a binary string





Lexicographic order

Which volume is selected?

• Encode each position with a binary string





Which volume is selected?

• Encode each position with a binary string

• A unique n-bit string per volume level





- Encode each position with a binary string
 - A unique n-bit string per volume level
 - Sensor behind the lever





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What happens between positions?





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• Non-issue if there is one ambiguous bit




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Problem is solved by using a Gray code





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Problem is solved by using a Gray code

Gray code

0 0

Λ

0

Volume



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- Problem if there are multiple ambiguous bits

Problem is solved by using a Gray code • Requires ranking/unranking

U.S. Patent on Pulse Communication

The term "Gray code" comes from this type of application.



Frank Gray at Bell Labs circa 1950

Rotary encoder

From Wikipedia, the free encyclopedia

A rotary encoder, also called a **shaft encoder**, is an electro-mechanical device that converts the angular position or motion of a shaft or axle to an analog or digital code.

There are two main types: absolute and incremental (relative). The output of absolute encoders indicates the current position of the shaft, making them angle transducers. The output of incremental encoders provides information about the *motion* of the shaft, which is typically further processed elsewhere into information such as speed, distance, RPM and position.

Rotary encoders are used in many applications that require precise shaft unlimited rotation including industrial controls, robotics, special purpose photographic lenses,^[1] computer input devices (such as optomechanical mice and trackballs), controlled stress rheometers, and rotating radar platforms.

Contents [hide]

1 Absolute rotary encoder

- 1.1 Construction
- 1.2 Mechanical absolute encoders
- 1.3 Optical absolute encoders
- 1.4 Standard binary encoding
- 1.5 Gray encoding
- 1.6 Single-track Gray encoding
- 1.7 Absolute encoder output formats
- 2 Incremental rotary encoder
- 3 Absolute versus incremental encoder terminology
 - 3.1 Traditional absolute encoders
 - 3.2 Traditional incremental encoders
- 4 Sine wave encoder

5 Use in industry





Gray code pattern n = 13

graycodes.com

Home

Permutations

Multiperms

Combinations

Bitstrings

Middle levels

Partitions

Dyck paths

Plane trees

Spanning trees

. .

De Bruijn

Semigroups

References

Generate permutations of a multiset

Generate all permutations of a multiset. Provide the multiset by specifying its frequency sequence, e.g. "1 1 3 2" for the multiset $\{0, 1, 2, 2, 2, 3, 3\}$. Output limit is 10.000 objects. Frequency 1132 sequence (max. 20) Algorithm Prefix rotation (Williams' cool-lex)

The first three algorithms running on this page are part of Jörg Arndt's FXT library. The iterative algorithm for lexicographic ordering is explained in Jörg Arndt's FXT library [Section 13.2.2, Arn10]. The minimal-change algorithm follows a Steinhaus-Johnson-Trotter adjacent transposition

Run

Results

T	•		L		0		3		3	2		2		2		T			
2	:		[3		0	1.1	3	2		2		2		1]	
3	:		[3		3	()	2		2		2		1]	
4	:		[2		3	1.1	3	0		2		2		1]	
5	•		[3		2	1.1	3	0		2		2		1]	
6	:		[0		3	2	2	3		2		2		1]	
7	:		[3		0	2	2	3		2		2		1]	
8	:		[2		3	()	3		2		2		1]	
9	:		[0		2	1.1	3	3		2		2		1]	
1(0	:		[2	(С	3)	3		2		2		1]
11	1	:		[3		2	0)	3		2		2		1]
12	2	:		[3		3	2		0		2		2		1]
1:	3	:		[2		3	3)	2		0		2		1]
1.	4	:		[3		2	3)	2		0		2		1]
1	5	:		[2		3	2		3		0		2		1]
1	6	:		[2	2	2	3		3		0		2		1]
1'	7	:		[3		2	2		3		0		2		1]
18	8	:		[0		3	2)	2		3		2		1]
1	9	:		[3	(С	2		2		3		2		1]
20	0	:		[2		3	С)	2		3		2		1]
2:	1	:		[0	4	2	3	}	2		3		2		1]
22	2	:		[2	(С	3	}	2		3		2		1]
23	3	:		[3	4	2	0)	2		3		2		1]
24	4	:		[2		3	2		0		3		2		1]
2!	5	:		[2	4	2	3	}	0		3		2		1]
2	6	:		[0	4	2	2		3		3		2		1]
2'	7	:		[2	(С	2		3		3		2		1]
28	8	•		[2	2	2	0)	3		3		2		1]
2	\circ			Г		2	,	2	0		\circ		2		0		1		٦







































THERE ARE 10 TYPES OF PEOPLE.

THOSE UHO KNOU BINARY AND THOSE UHO DON⁷T. THERE ARE 11 TYPES OF PEOPLE.

Those who know Gray Code And Those who don't.