CSCI 136
Data Structures & Advanced Programming

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Lecture 9

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Administrative Details

- Lab 3 is due next Monday
  - Run-time (big-O notation) in comments above methods
  - Don't forget to define "n" if you say O(n)
- I posted Warmup solutions with and without check for empty set on Handouts page

Last Time

- Discussed runtime analysis techniques
- Reviewed and discussed recursion
  - Looked at factorial, digit sum, subset sum

Today's Outline

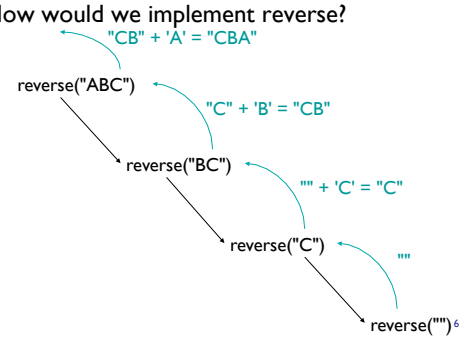
- Wrapup recursion review
- Begin reviewing mathematical induction
- Begin learning about searching and sorting
  - Two of the most important classes of algorithms
- Two searches:
  - Linear search
  - Binary search

Palindromes

- Implement boolean isPalindrome(String str)
- What are our base cases?
  - Strings of length 0
  - Strings of length 1
- What is our recursive case?
  - If str.length() > 1
    - Pull off first and last char. If they match, recurse on remaining string.

Reversing Strings

- How would we implement reverse?



Towers of Hanoi

- Demo
- Base case:
  - One disk: Move from start to finish
- Recursive case (n disks):
  - Move smallest n-1 disks from start to temp
  - Move bottom disk from start to finish
  - Move smallest n-1 disks from temp to finish

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Hanoi.java

```
class Hanoi {

  public static void moveDisk(String start, String finish) {
        System.out.println(start + " -> " + finish);
  }

  public static void moveTower(int n, String start, String finish, String temp) {
        if (n == 1) {
            moveDisk(start, finish);
        } else {
            moveTower(n-1, start, temp, finish);
            moveDisk(start, finish);
            moveTower(n-1, temp, finish, start);
        }
  }

  public static void main(String args[]) {
        moveTower(Integer.parseInt(args[0]), "A", "B", "C");
  }

}
```

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Recursion Tradeoffs

- Advantages
  - Often easier to construct recursive solution
  - Code is usually cleaner
  - Some problems do not have obvious non-recursive solutions
- Disadvantages
  - Overhead of recursive calls
  - Can use lots of memory (need to store state for each recursive call until base case is reached)

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Mathematical Induction

- The mathematical equivalent of recursion is induction
- Induction is a proof technique
- Comes from how natural numbers are defined:
  - A is the set of natural numbers such that
    1. 0 is an element of A
    2. For each n, if 0, 1, 2, …, n-1 are in A, than n is in A.

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Mathematical Induction

- Examples

$$P = \sum_{i=0}^{n} i = 0 + 1 + \ldots + n = \frac{n(n+1)}{2}$$

- Proof by induction:
  - Base case: P is true for 0
  - Inductive hypothesis: If P is true for all k<n, then P is true for n.

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Mathematical Induction

- Prove: $\sum_{i=0}^{n} 2^i = 2^0 + 2^1 + 2^2 + \ldots + 2^n = 2^{n+1} - 1$

- Prove: $0^3 + 1^3 + \ldots + n^3 = (0 + 1 + \ldots + n)^2$

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Proof: $0^3 + 1^3 + ... + n^3 = (0 + 1 + ... + n)^2$

- Base case: $n = 0, 0^3 = (0)^2$.
- Ind. Hyp.: For $k < n$ assume true.
  - $(0^3 + 1^3 + ... + k^3) = (0 + 1 + ... + k)^2$
- Ind Step: Show true for n.

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Proof: $0^3 + 1^3 + ... + n^3 = (0 + 1 + ... + n)^2$

$$\begin{aligned}
(0^3 + 1^3 + ...n^3) &= (0^3 + 1^3 + ... + (n-1)^3) + n^3 \\
&= (0 + 1 + ... + (n-1))^2 + n^3 \\
&= \left(\frac{n(n-1)}{2}\right)^2 + n^3 \\
&= n^2\left(\frac{(n-1)^2 + 4n}{4}\right) \\
&= n^2\left(\frac{n^2 + 2n + 1}{4}\right) \\
&= n^2\left(\frac{(n+1)^2}{4}\right) \\
&= \left(\frac{n(n+1)}{2}\right)^2 \\
&= (0 + 1 + ... + n)^2
\end{aligned}$$

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What about Recursion?

- What does induction have to do with recursion?
  - Same form!
    - Base case
    - Inductive case that uses simpler form of problem
- Example: factorial
  - Prove that fact(n) requires n multiplications
    - Base case: n = 0 returns 1, 0 multiplications (note that technically 1! = 1*0!)
    - Assume true for all k<n, so fact(k) requires k multiplications.
    - fact(n) performs one multiplication (n*fact(n-1)). We know that fact(n-1) requires n-1 multiplications by inductive hypothesis.
      1+n-1 = n, therefore fact(n) requires n multiplications.

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Moving on…

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Searching

- What is searching?
  - Locate element in collection
  - Examples: Number in list, grade in gradebook, etc
  - Complexity analysis, induction, recursion
- Today's algorithms
  - Linear search - move down line
  - Binary search - divide elements in half
- Next up
  - Sorting
  - Designing data structures to support other searching/ sorting algorithms

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Linear Search

- Where have we seen a linear search?
  - Dictionary.java!

```
// post: returns the definition of word, or "" if not found.
   public String lookup(String word) {
      for (int i = 0; i < words.length; i++) {
           Association a = words[i];
           if (a.getKey().equals(word)) {
               return (String)a.getValue();
           }
       }
       return null;
   }
```

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Linear Search

```
public class LinearSearchComp {
    // post: returns index of value in a, or -1 if not found
    public static <E> int linearSearch(E a[], E value) {
        for (int i = 0; i < a.length; i++) {
            if (a[i].equals(value)) {
                return i;
            }
        }
        return -1;
    }

    public static void main(String args[]) {
        // search a String array
        System.out.println(linearSearch(args, "cow"));

        // search a Linear array
        Integer odds[] = new Integer[] { 1,3,5,7,9 };
        System.out.println(linearSearch(odds, 7));
    }
}
```

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Linear Search

- Complexity analysis of linear search:
  - Best case:
    - O(1)
  - Worst case:
    - O(n)
  - Average case
    - O(n)

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