

CSCI 136 Data Structures & Advanced Programming

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Lecture 36
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Administrative Details

- Final exam - self scheduled
 - You get 2.5 hours to complete it
 - Covers everything, w/ strong emphasis on Ch 14-16
 - BSTs, HashTables, Maps, Graphs
- NOTE: I will be out of town this weekend
 - Otherwise I'll be mostly around next week

Last Time

- Wrapped up graphs
- Started discussing Maps

Today's Outline

- Finish discussing HashMaps/HashTables

Recap: Hashing in a Nutshell

- Group objects into "bins"
- When searching for object, go directly to appropriate bin
- If there are multiple objects in bin, then search (linearly) for correct one
- Important Insight: This works best when objects are evenly distributed among bins
- Must deal with collisions

Recap: Linear Probing

- If a collision occurs at a given bin, just move forward (linearly) until an empty slot is available
- Need a "placeholder" for removed values...

```
//OVERSIMPLIFIED VERSION OF PUT
public void put (K key, V val) {
    int index = key.hashCode() % arraySize;
    while (index < array.length && array[index]!=null &&
        !array[index].getKey().equals(key)) {
        index++;
    }
    array[index] = new Association(key, val);
}
```

Recap: Linear Probing

- If a collision occurs at a given bin, just move forward (linearly) until an empty slot is available
- Need a “placeholder” for removed values...

```
//OVERSIMPLIFIED VERSION OF GET
public V get (K key) {
    int index = key.hashCode() % arraySize;
    while (index < array.length && array[index]!=null &&
        !array[index].getKey().equals(key)) {
        index++;
    }
    if (index==array.length) return null;
    return array[index];
}
```

Linear Probing

- Runtime
 - put, get, remove
 - $O(1)$
 - ...as long as the array isn't too full!

External Chaining

- Downsides of linear probing
 - What if array is almost full?
 - Linear probing is extremely difficult on almost-full arrays
- How can we avoid this problem?
 - Keep all values that hash to same bin in a “collection”
 - Usually a SLL
 - External chaining “chains” objects with the same hash value together

Example

- Recall our previous example
 - Add algorithm
 - Add data
 - Add queue
 - Now remove algorithm
 - No need for placeholders!

External Chaining

- Let's look at put(key, val)...

```
public V put(K key, V val) {
    //locate returns SLL in the appropriate bin
    SLL list = locate(key);
    ComparableAssociation newA =
        new ComparableAssociation(key, val);
    ComparableAssociation oldA = list.remove(newA);
    list.add(newA);
    return oldA.getValue();
}
```

External Chaining

- Let's look at put(key, val)...
- Runtime
 - put, get, and remove
 - $O(\ell)$ (ℓ is size of linked list at given bin)
 - As long as table isn't too full, this is almost $O(1)$

Load Factor

- Need to keep track of how full the table is
 - Why?
 - What happens when array fills completely?
- Load factor is a measure of how full the hash table is
 - $LF = \# \text{ elements} / \text{table size}$
- When LF reaches some threshold, need to double size of array (typically threshold = 0.6)
 - How?

Doubling Array

- Cannot just copy values
 - Why?
 - Hash values may change (i.e., element "list")
 - Example
 - $\text{key.hashCode()} \% 8 = 3$;
 - $\text{key.hashCode()} \% 16 = 11$;
- Have to recompute all hash codes
- This is expensive!

Good Hashing Functions

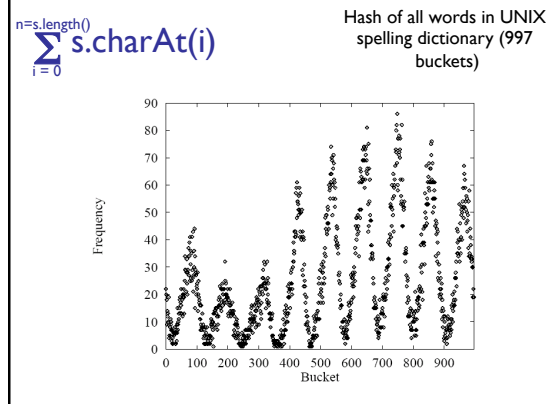
- Important point:
 - All of this hinges on using "good" hash functions that spread keys "evenly"
- Two properties of good hash functions:
 - Fast to compute
 - Uniformly distribute keys
- Almost always have to test "goodness" empirically

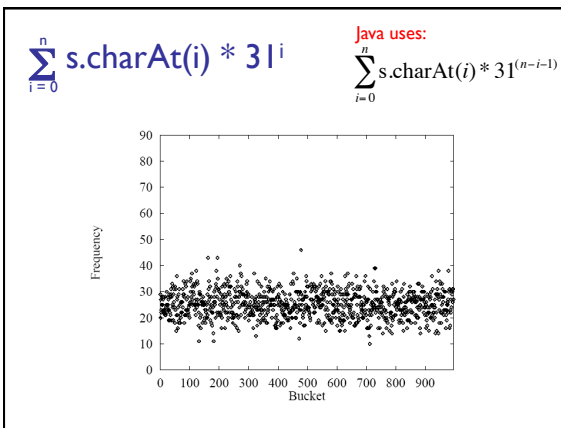
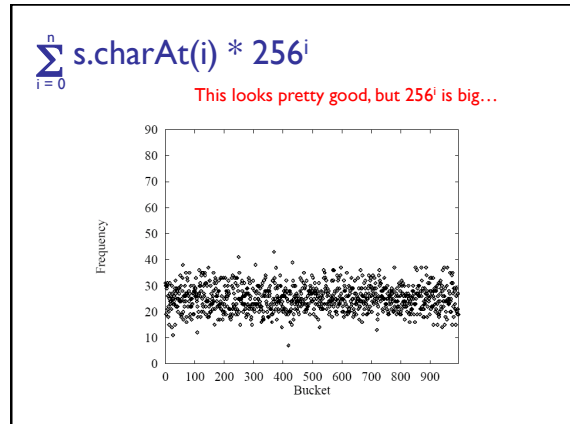
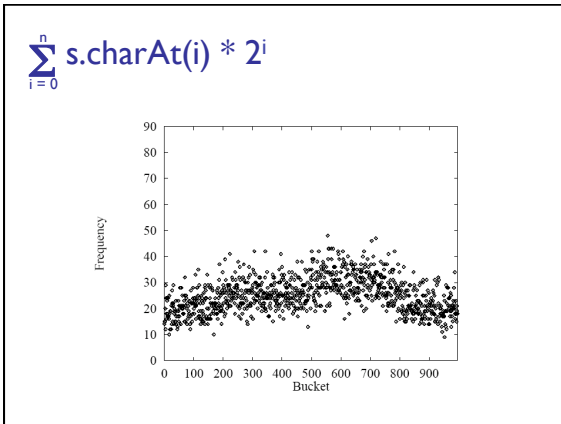
Example Hash Functions

- What are some feasible hash functions for Strings?
 - First char ASCII value mapping
 - 0-255 only
 - Not uniform (some letters more popular than others)
 - Sum of ASCII characters
 - Not uniform - lots of small words
 - smile, limes, miles, slime are all the same

Example Hash Functions

- String hash functions
 - Weighted sum
 - Small words get bigger codes
 - Distributes keys better than non-weighted sum
 - Let's look at different weights...





Dictionary Summary

	put	get	space
unsorted vector	$O(n)$	$O(n)$	$O(n)$
unsorted list	$O(n)$	$O(n)$	$O(n)$
ordered vector	$O(n)$	$O(\log n)$	$O(n)$
balanced BST	$O(\log n)$	$O(\log n)$	$O(n)$
array indexed by key	$O(1)$	$O(1)$	$O(\text{key range})$

- ### Course Evals – Wrap Up
- This semester we have explored structures to represent data
 - Started with simple arrays
 - Ended with complicated graphs and hash tables
 - We studied algorithms to manipulate and use these data structures
 - We studied ways to evaluate and visualize them
 - Emphasized importance of abstraction, modular design, correctness, and efficiency

- ### Topics Covered
- Vectors (and arrays)
 - Complexity (big O)
 - Recursion + Induction
 - Searching
 - Sorting
 - LinkedLists
 - Stacks
 - Queues
 - Iterators
 - Comparables/Comparators
 - OrderedStructures
 - Binary Trees
 - Priority Queues
 - Heaps
 - Binary Search Trees
 - Graphs
 - Maps/Hashtables

What's Next?

- Data structures provides you with a complete tool box to study the rest of Computer Science
 - Programming Languages
 - Artificial Intelligence
 - Compilers
 - Networks
 - Distributed Systems
 - Algorithms
 - Operating Systems
 - ...
- Also applicable to many areas outside of CS

Thanks!!

- Any questions?
- Thanks for a great semester!
- Have a great summer!
- Congrats seniors!