

CSCI 136 Data Structures & Advanced Programming

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Lecture 32
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Administrative Details

- Lab today: Exam Scheduling
 - Focuses on using and manipulating graphs
- Darwin tournament on Friday
 - I'll need your final creatures by tomorrow at noon!

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Lab 10 Notes: Graphs in structure5

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Lab 10 Notes: Using Graphs

- Create a new graph in structure5
 - GraphList, GraphListDirected, GraphListUndirected,
 - GraphMatrix, GraphMatrixDirected, GraphMatrixUndirected
- `Graph<V,E> dep = new GraphListUndirected<V,E>();`

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Lab 10 Notes: Greedy Algorithm

- Pick a starting vertex – add it to list
- Walk through other vertices in graph
 - If vertex is not connected to all other vertices in list, add it to list
- Remove all vertices in list from graph
- Repeat until all vertices are gone

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Lab 10 Notes: Useful Graph Methods

- `void add(V label)`
 - add vertex to graph
- `void addEdge(V vtx1, V vtx2, E label)`
 - add edge between vtx1 and vtx2
- `Iterator<V> neighbors(V vtx1)`
 - Get iterator for all neighbors to vtx1
- `boolean isEmpty()`
 - Returns true iff graph is empty
- `Iterator<V> iterator()`
 - Get vertex iterator
- `V remove(V label)`
 - Remove a vertex from the graph
- `E removeEdge(V vLabel1, V vLabel2)`
 - Remove an edge from graph

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Last Time

- Looked at GraphMatrix implementation
 - Define graph using an adjacency matrix
 - Matrix keeps track of edge weights

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Recap: GraphMatrix

- Abstract class – partially implements Graph

```
abstract public class GraphMatrix<V,E> implements Graph<V,E>
```

- Instance variables

```
protected int size; //max size of matrix
protected Object data[][]; //matrix of edges
protected Map<V, GMV<V>> dict; //labels -> vertices
protected List<Integer> freeList; //avail indices
protected boolean directed;
```

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Recap: GraphMatrix Constructor

```
protected GraphMatrix(int size, boolean dir) {
    this.size = size; // set maximum size
    directed = dir; // fix direction of edges

    // the following constructs a size x size matrix
    data = new Object[size][size];

    // label to index translation table
    dict = new Hashtable<V,GraphMatrixVertex<V>>(size);

    // put all indices in the free list
    freeList = new SinglyLinkedList<Integer>();
    for (int row = size-1; row >= 0; row--)
        freeList.add(new Integer(row));
}
```

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Recap: GraphMatrixDirected

- Constructor

```
public GraphMatrixDirected(int size) {
    // pre: size > 0
    // post: constructs an empty graph that may be
    //        expanded to at most size vertices. Graph
    //        is directed if dir true and undirected
    //        otherwise
    super(size, true);
}
```

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Recap: GraphMatrixDirected

- add and addEdge

```
public void add(V label) {
    (dict.containsKey(label)) return;
    int row = freeList.removeFirst().intValue();
    dict.put(label, new GraphMatrixVertex<V>(label, row));
}

public void addEdge(V vLabel1, V vLabel2, E label) {
    GraphMatrixVertex<V> vtx1,vtx2;
    vtx1 = dict.get(vLabel1);
    vtx2 = dict.get(vLabel2);
    Edge<V,E> e = new Edge<V,E>(vtx1.label(), vtx2.label(),
        label, true);
    data[vtx1.index()][vtx2.index()] = e;
}
```

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Today's Outline

- Continue talking about graphs
 - Finish up GraphMatrix
 - Discuss GraphList

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Removing

- Now let's look at removing vertices and edges in GraphMatrixDirected

```

• public V remove(V label)

• public E removeEdge(V vlabel1, vlabel2)

```

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Removing Edges

```

public E removeEdge(V vlabel1, vlabel2) {
    // get indices
    int row = dict.get(vLabel1).index();
    int col = dict.get(vLabel2).index();
    // save old value
    Edge<V,E> e = (Edge<V,E>)data[row][col];
    // update matrix
    data[row][col] = null;
    if (e == null) return null;
    else return e.label(); // return old value
}

```

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Removing Vertices (same for GMD and GMU)

```

public V remove(V label) {
    // find and extract vertex
    GraphMatrixVertex<V> vert;
    vert = dict.remove(label);
    if (vert == null) return null;
    // remove vertex from matrix
    int index = vert.index();
    // clear row and column entries
    for (int row=0; row<size; row++) {
        data[row][index] = null;
        data[index][row] = null;
    }
    // add node index to free list
    freeList.add(new Integer(index));
    return vert.label();
}

```

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Efficiency

- Assume map operations are $O(1)$ (for now)
 - $|E|$ = number of edges
 - $|V|$ = number of vertices
- Runtime of add, addEdge, getEdge, removeEdge?
- Runtime of remove?
- Space usage?
 - Can we do better?
 - YES! Don't store empty slots in array.

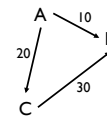
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GraphList

- Rather than keep an adjacency matrix, maintain an *adjacency list of edges* at each vertex (only keep outgoing edges for directed graphs)
- Support both directed and undirected graphs (GraphListDirected, GraphListUndirected)

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Example



Adjacency Matrix:

	A	B	C
A	-	10	20
B	-	-	-
C	-	30	-

Adjacency List:

A: [A->B,10] [A->C,20]
 B:
 C: [C->B,30]

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GraphListDirected

- To implement GLD, we use the GraphListVertex (GLV) class
- GraphListVertex class
 - Maintain linked list of edges at each vertex
 - Instance vars: label, visited flag, linked list of edges
- GraphList abstract class
 - Instance vars:
 - Map<V, GraphListVertex<V, E>> dict; // label -> vertex
 - boolean directed; // is graph directed?
- How do we implement key GLD methods?
 - add, addEdge, removeEdge, remove

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```
// add is same for GraphListDirected and GraphListUndirected,
// so implement in GraphList.java
public void add(V label) {
    // check to see if vertex already exists
    if (dict.containsKey(label)) return;
    GraphListVertex<V,E> v = new GraphListVertex<V,E>(label);
    dict.put(label,v);
}

// addEdge in GraphListDirected.java
// first vertex is source, second is destination
public void addEdge(V vLabel1, V vLabel2, E label) {
    // first get the vertices
    GraphListVertex<V,E> v1 = dict.get(vLabel1);
    GraphListVertex<V,E> v2 = dict.get(vLabel2);
    // create the new edge
    Edge<V,E> e = new Edge<V,E>(v1.label(), v2.label(), label, true);
    // add edge only to source vertex linked list (aka adjacency list)
    v1.addEdge(e);
}
```

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```
public V remove(V label) {
    //Get vertex out of map/dictionary
    GraphListVertex<V,E> v = dict.get(label);

    //Iterate over all vertex labels (called the map "keyset")
    Iterator<V> vi = iterator();
    while (vi.hasNext()) {
        //Get next vertex label in iterator
        V v2 = vi.next();

        //Skip over the vertex label we're removing
        //(Nodes don't have edges to themselves...)
        if (!label.equals(v2)) {
            //Remove all edges to "label"
            //If edge does not exist, removeEdge returns null
            removeEdge(v2,label);
        }
    }
    //Remove vertex from map
    dict.remove(label);
    return v.label();
}
```

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```
public E removeEdge(V vLabel1, V vLabel2) {
    //Get vertices out of map
    GraphListVertex<V,E> v1 = dict.get(vLabel1);
    GraphListVertex<V,E> v2 = dict.get(vLabel2);

    //Create a "temporary" edge connecting two vertices
    Edge<V,E> e = new Edge<V,E>(v1.label(), v2.label(), null, true);

    //Remove edge from source vertex linked list
    e = v1.removeEdge(e);
    if (e == null) return null;
    else return e.label();
}
```

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Efficiency Revisited

- Assume map operations are $O(1)$ (for now)
 - $|E|$ = number of edges
 - $|V|$ = number of vertices
- Runtime of add, addEdge, getEdge, removeEdge, remove?
- Space usage?
- Conclusions
 - Matrix is better for dense graphs
 - List is better for sparse graphs
 - For graphs "in the middle" there is no clear winner

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Efficiency

	Matrix	Sparse GraphList	Dense GraphList
add	$O(1)$	$O(1)$	$O(1)$
addEdge	$O(1)$	$O(1)$	$O(1)$
getEdge	$O(1)$	$O(1)$	$O(V)$
removeEdge	$O(1)$	$O(1)$	$O(V)$
remove	$O(V)$	$O(E) = O(V)$	$O(E) = O(V ^2)$
space	$O(V ^2)$	$O(E) = O(V)$	$O(E) = O(V ^2)$

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Moving On...

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Graph Traversal Algorithms

- General concepts
 - Reachability
 - Cycle detection
 - Shortest path
 - ...
 - (Read Ch 16!)

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Recall our campus map...

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Recall our campus map...separated by Rt. 2

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Suppose we applied a direction to our edges...

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Reachability

- There are two ways to measure reachability in our graph
 - Depth-first search and breadth-first search
- How did we do DFS and BFS in trees?
- DFS uses a stack
 - Stack records path from src to current node
 - Like pre-order tree traversal with visited flags to only visit nodes once
 - Runtime: $O(|E|)$
- BFS uses a queue
 - Queue records nodes whose out edges have not been explored
 - Like level-order tree traversal
 - Runtime: $O(|E|)$

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Depth-First Search

```
protected void reachableFrom(Graph<V,E> g, V src) {  
  
    if (!g.visited(src)) {  
        g.visit(src);  
  
        Iterator<V> neighbors = g.neighbors(src);  
        while (neighbors.hasNext()) {  
            V next = neighbors.next();  
            if (!g.visited(next))  
                reachableFrom(g, next);  
        }  
    }  
}
```

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