

Administrative Details

- Lab 7
- Due Monday at noon
- Questions?


## Last Time

- Learned about ordered structures (Ch II)
- Talked about OrderedVector and OrderedList


## Today's Outline

- Begin learning about trees
- Very important data structure!
- Main advantage is that the data is always sorted
- Restrict add method so that objects are always added "in the right spot"
- Easy to find min, max, median
- Be careful not to use mutable keys!


## Data Structures so far...

- So far all data has been stored in a linear fashion
- Stacks, queues
- Even arrays, vectors, SLLs are visualized using linear structures
- By linear we mean that each element has only one successor and one predecessor...



## Introducing Trees

- A tree is a data structure where elements can have multiple successors (called children)
- But still only one predecessor (called parent)



## Other Trees

- Phylogenetic tree
- Directories of files
- Game tree
- Build tree
- Search for moves with high likelihood of winning
- Expression trees (we'll come back to these in a bit)



## Trees in CS 136

- A tree is a data structure where elements can have multiple successors (but generally only one predecessor)



## Tree Features

- Hierarchical relationship
- Root at the top
- Leaf at the bottom
- Interior nodes in middle
- Parents, children, ancestors, descendants, siblings
- Degree: max number of children per node
- Depth of node $n$ : number of edges from root to $n$
- Height: max depth (across all nodes)


## Binary Trees

- Degree of all nodes <= 2
- Recursive nature of tree
- Base case: Empty
- Rec. case: Root with left and right subtrees (also BTs)
- SLL: Recursive nature was captured by elements (SLLEs) that pointed to other elements (SLLEs)
- Binary Tree: No second element class; single BinaryTree class does it all!


## Full vs. Complete

- Full tree - A full binary tree of height $h$ has leaves only on level h, and each internal node has exactly 2 children.

- Complete tree - A complete binary tree of height $h$ is a full tree with 0 or more rightmost leaves of level h removed.


All full trees are complete, but not all complete trees are full!

## Expression Trees

## - General strategy

- Make a binary tree (BT) for each leaf node
- Move from bottom to top, creating BTs
- Eventually reach the root
- Call "evaluate" on final BT


## - Example

- How do we make a binary expression tree for (((4+3)*(10-5))/2)
- Postfix notation: $43+105-* 21$
public int evaluate(BinaryTree<String> expr) \{
if (expr.height() $==0$ ) \{
return Integer.parseInt(expr.value());
\} else \{
int left = evaluate(expr.left()); int right = evaluate(expr.right()); String op = expr.value(); if (op.equals("+")) return left + right; if (op.equals("-")) return left - right; if (op.equals("*")) return left * right; if (op.equals(")")) return left / right; Assert.fail("Bad op"); return -1; //Why do we need this?
\}
\}

BinaryExpressionTree.java


## Implementing BinaryTree

- Methods (on board - see Ch 12 for more info):
- (All "left" methods have equivalent "right" methods)
- public BinaryTree()
- // generates an empty node (EMPTY)
- I/ parent and value are null, left=right=this
- public BinaryTree(E value)
- // generates a tree with a non-null value and two empty (EMPTY) subtrees
- public BinaryTree(E value, BinaryTree<E> left, BinaryTree<E> right)
- // returns a tree with a non-null value and two subtrees
- public void setLeft(BinaryTree<E> newLeft)
- // sets left subtree to newLeft
- // re-parents newLeft (if not null) by calling newLeft.setParent(this)
- protected void setParent(BinaryTree<E> newParent)
- // sets parent subtree to newParent
- // called from setLeft and setRight to keep all "links" consistent

