

[TAP:XLKTU] Balanced Trees

- Which of the following are not guaranteed to be “balanced”?
 - A. AVL Tree
 - B. Red-black Tree
 - > C. Heap
 - D. They are all balanced
 - E. Whatever

Administrative Details

- Lab 9 Today: Gardner's Hex-a-Pawn
 - Another partner lab!
 - Challenging to design & debug
 - Design doc is worth 2 points of your lab grade

Today's Outline

- Balanced Binary Search Trees
 - AVL Tree
 - Red-Black Tree
 - • (Splay Tree)
- Game Trees

Splay Trees

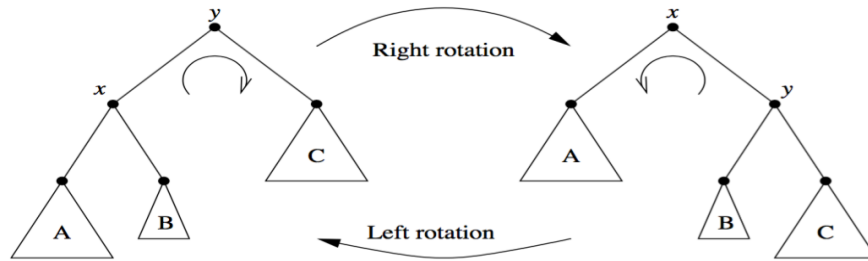
Splay trees are self-adjusting binary trees

- Each time a node is accessed, it is moved to root position via rotations *includes "contains()"*
- No guarantee of balance (or shallow height)
- But good *amortized* performance

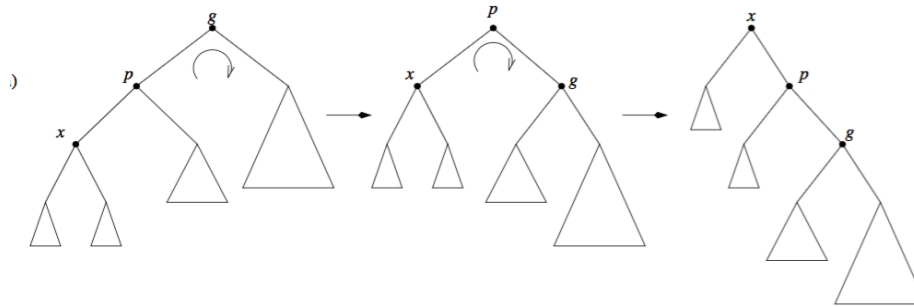
Theorem: Any set of m operations (add, remove, contains, get) on an n -node splay tree take at most $O(m \log n)$ time.

Splay Tree Rotations

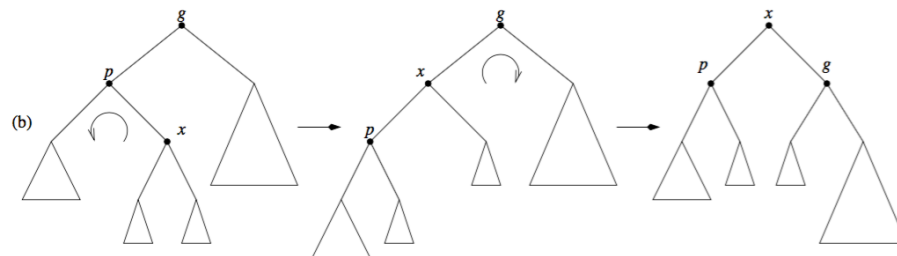
Right Zig Rotation (left version too)



Right Zig-Zig Rotation (left version too)



Right Zig-Zag Rotation (left version too)



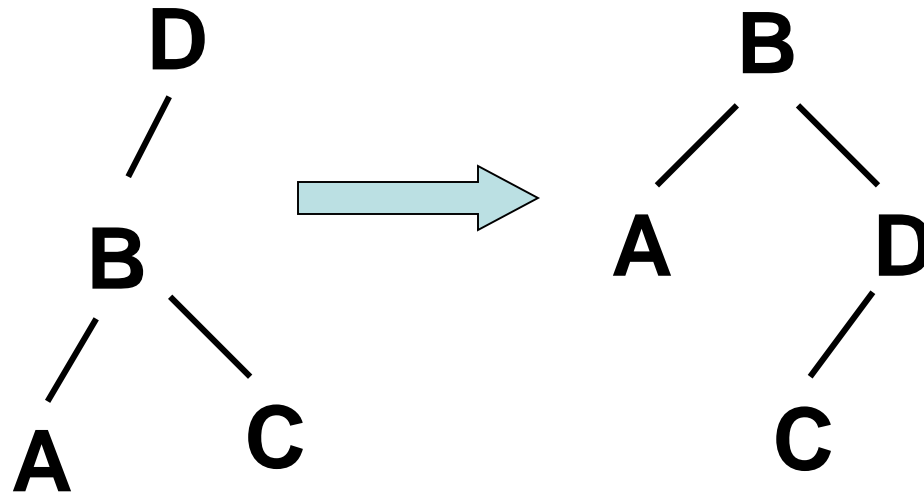
AVL Tree Facts

- A tree that is AVL except at root, where root balance factor equals ± 2 can be rebalanced with at most 2 rotations
- $\text{add}(v)$ requires at most $O(\log n)$ balance factor changes and **one** (single or double) rotation to restore AVL structure
- $\text{remove}(v)$ requires at most $O(\log n)$ balance factor changes and $O(\log n)$ (single or double) rotations to restore AVL structure
- An AVL tree on n nodes has height $O(\log n)$

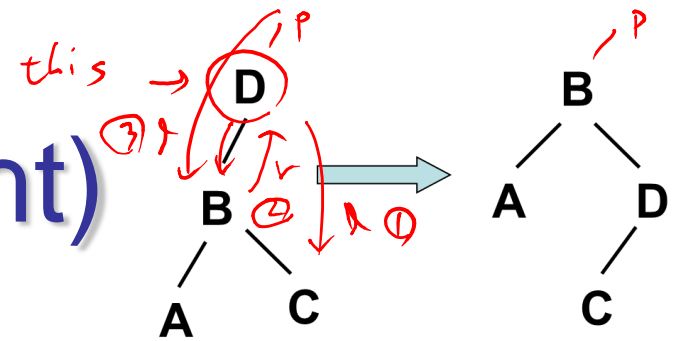
Red-Black Tree Facts

- The coloring rules lead to the following result
 - No leaf has depth more than twice that of any other leaf.
 - A Red-Black tree with n nodes has height $O(\log n)$

Single Rotation (Right)



Single Rotation (Right)



```
// pre: this has a left subtree
```

```
// post: rotates local portion of tree so left child is root
```

```
protected void rotateRight() {
```

```
    BinaryTree<E> p = parent;  
    boolean wasLeftChild = isLeftChild();
```

```
    BinaryTree<E> newRoot = left;
```

```
    setLeft(newRoot.right); // ①
```

```
    newRoot.setRight(this); // ②
```

```
    if (parent != null) {  
        if (isLeftChild() wasLeftChild)
```

```
            parent.setLeft(newRoot); // ③
```

```
        } else Pparent.setRight(newRoot);
```

```
        newRoot.setParent(parent);  
        setParent(newRoot);      left.setParent(this);
```

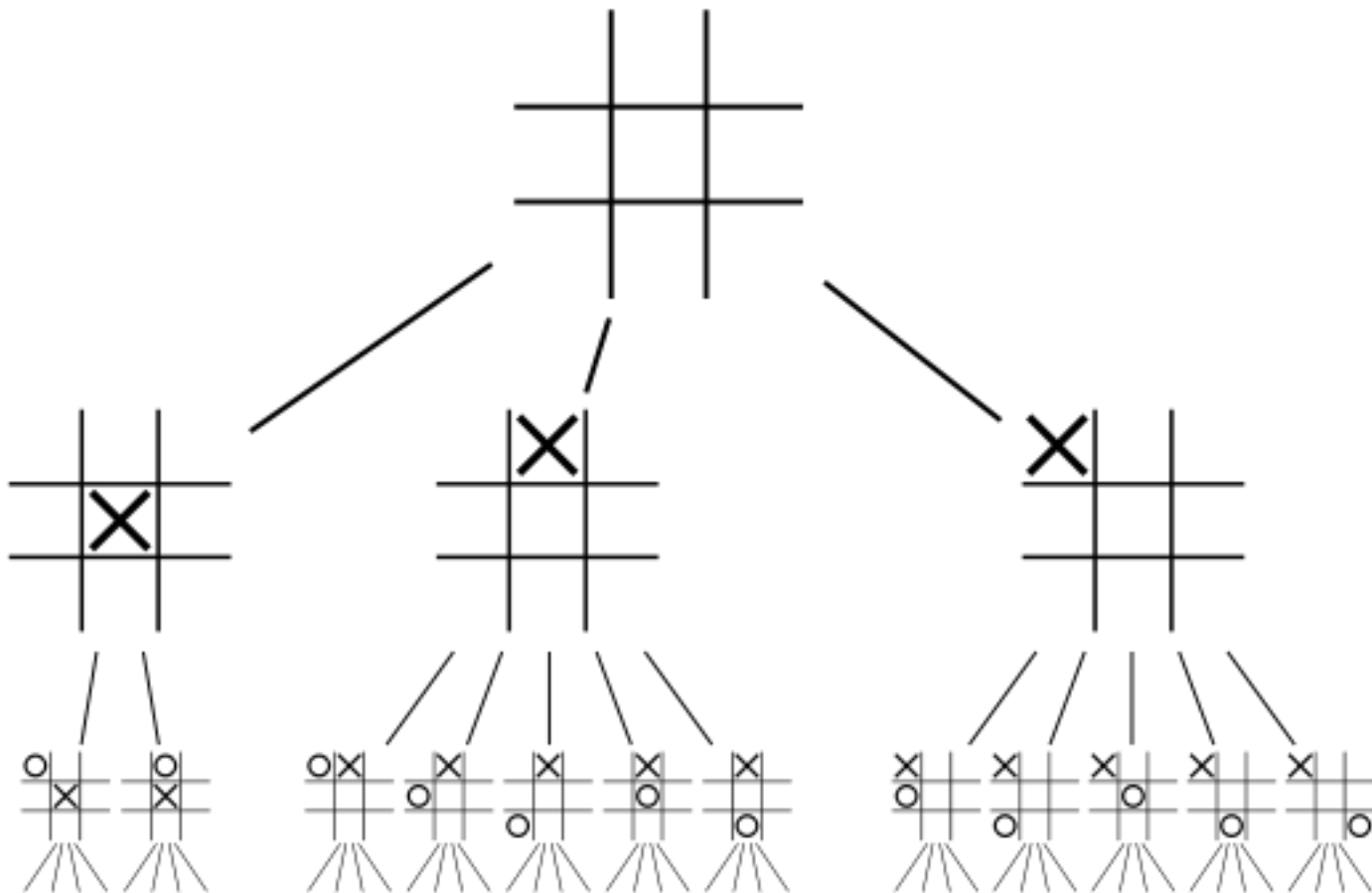
purple if setLeft() and setRight()

updates the parent

Today's Outline

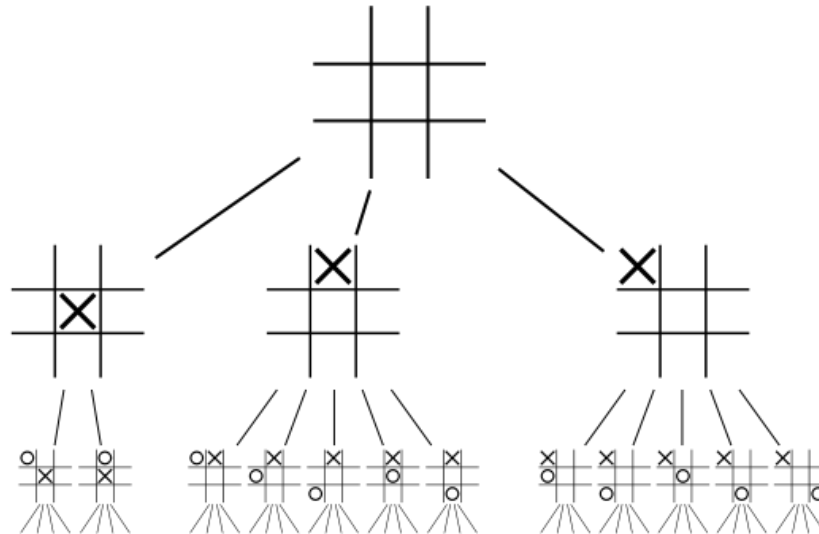
- Balanced Binary Search Trees
 - AVL Tree
 - Red-Black Tree
 - (Splay Tree)
- • Game Trees

Game Trees



Game Trees

- Nodes are positions in a game (game state)
- Edges are moves (transition from one game state to another)
 - All nodes at a given level represent moves by the same player
- Leaf nodes represent ending board states (winner or tie)
 - # of leaf nodes = # of ways a game can be played

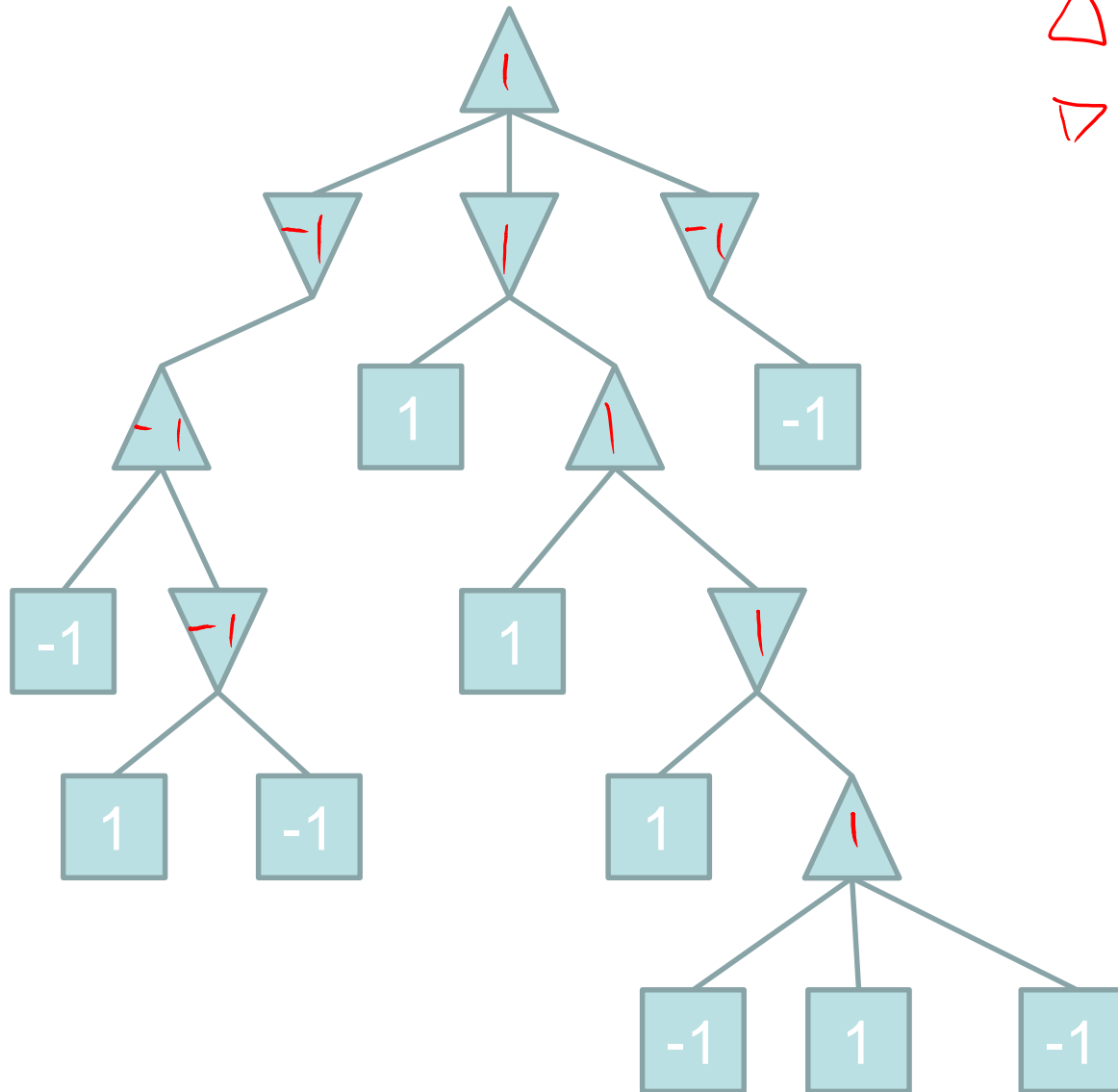


Game Trees

- In AI, often search the game tree and use an algorithm like **minimax** to choose the next “best move”
 - Chess, checkers, Go, etc.



Minimax Example



\triangle max \leftarrow you
 ∇ min \leftarrow opponent

Modern Game AI



Hexapawn

