

CSCI 136

Data Structures & Advanced Programming

Lecture 13
Spring 2018
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Administrative Details

- Lab 5 Posted
 - Sorting with Comparators
- Midterm Wednesday March 14
 - Held in your scheduled Lab (same time and place)
 - Study guide and sample exam
 - Review session

Last Time

- The Comparable Interface
 - Including: how to write a generic static method
 - Generic Linear and Binary Search methods
- “Basic” Sorting
 - Bubble sort

Today's Outline

- “Basic” Sorting Wrapup
 - Bubble, Insertion, Selection Sorts
- Comparator: interface for flexible sorting
- More Efficient Sorting Algorithms
 - MergeSort
 - QuickSort

Basic Sorting Algorithms

- BubbleSort
 - Swaps consecutive elements of $a[0..k]$ until largest element is at $a[k]$; Decrements k and repeats
- InsertionSort
 - Assumes $a[0..k]$ is sorted and moves $a[k+1]$ across $a[0..k]$ until $a[0..k+1]$ is sorted
 - Increments k and repeats
- SelectionSort
 - Finds largest item in $a[0..k]$ and swaps it with $a[k]$
 - Decrements k and repeats

Sorting Preview: Bubble Sort

- Simple sorting algorithm that works by ascending through the list to be sorted, comparing two items at a time, and swapping them if they are in the wrong order
- Repeated until no swaps are needed
- Gets its name from the way larger elements "bubble" to the end of the list

Bubble Sort

5 1 3 2 9

- First Pass:

- (**5** 1 3 2 9) → (1 **5** 3 2 9)
- (1 **5** 3 2 9) → (1 3 **5** 2 9)
- (1 3 **5** 2 9) → (1 3 2 **5** 9)
- (1 3 2 **5** 9) → (1 3 2 5 **9**)

- Second Pass:

- (**1** 3 2 5 9) → (**1** 3 2 5 9)
- (1 **3** 2 5 9) → (1 2 **3** 5 9)
- (1 2 **3** 5 9) → (1 2 3 **5** 9)

- Third Pass:

- (**1** 2 3 5 9) → (**1** 2 3 5 9)
- (1 **2** 3 5 9) → (1 **2** 3 5 9)

- Fourth Pass:

- (**1** 2 3 5 9) → (**1** 2 3 5 9)

<http://www.youtube.com/watch?v=lyZQPjUT5B4>

Bubble Sort

- Simple sorting algorithm that works by ascending through the list to be sorted, comparing two items at a time, and swapping them if they are in the wrong order
- Repeated until no swaps are needed
- Gets its name from the way larger elements "bubble" to the end of the list
- Time complexity?
 - $O(n^2)$
- Space complexity?
 - $O(n)$ total (no additional space is required)

Sorting Preview: Insertion Sort

- Simple sorting algorithm that works by building a sorted list one entry at a time
- Sorted list in low region of the array
- To-be-sorted part in upper region
- Each time you “grow” your sorted region, you swap it backwards into its sorted location

Sorting Preview: Insertion Sort

- 5 7 0 3 4 2 6 1
- 5 7 0 3 4 2 6 1
- 0 5 7 3 4 2 6 1
- 0 3 5 7 4 2 6 1
- 0 3 4 5 7 2 6 1
- 0 2 3 4 5 7 6 1
- 0 2 3 4 5 6 7 1
- 0 1 2 3 4 5 6 7

Red: sorted region.

Each round, swap the first unsorted item back into sorted region

Sorting Preview: Insertion Sort

- Less efficient on large lists than more advanced algorithms
- Advantages:
 - Simple to implement and efficient on small lists
 - Efficient on data sets which are already substantially sorted
- Time complexity
 - $O(n^2)$
- Space complexity
 - $O(n)$

Sorting Preview: Selection Sort

The algorithm works as follows:

- Find the maximum value in the list
- Swap it with the value in the last position
- Repeat the steps above for remainder of the list (ending at the second to last position)

Sorting Preview: Selection Sort

- 11 3 27 5 16 Swap 27 with 16
- 11 3 16 5 27 Swap 16 with 5
- 11 3 5 16 27 Swap 11 with 5
- 5 3 11 16 27 Swap 5 with 3
- 3 5 11 16 27 Done!

Sorting Preview: Selection Sort

- Similar to insertion sort
- Performs worse than insertion sort in general
- Noted for its simplicity and performance advantages when compared to complicated algorithms
- Time Complexity:
 - $O(n^2)$
- Space Complexity:
 - $O(n)$

Basic Sorting Algorithms

(All Run in $O(n^2)$ Time)

- BubbleSort
 - Always performs cn^2 comparisons and might need to perform cn^2 swaps
- InsertionSort
 - Might need to perform cn^2 comparisons and cn^2 swaps
- SelectionSort
 - Always performs cn^2 comparisons but only $O(n)$ swaps

Swap!

- The “Basic” sorts all use a utility method: swap.
How would you implement swap?

```
private static void swap(int[] a, int i, int j) {  
    int temp = a[i];  
    a[i] = a[j];  
    a[j] = temp;  
}
```


Aside: Lower Bound Notation

Definition: A function $f(n)$ is $\Omega(g(n))$ if for some constant $c > 0$ and all $n \geq n_0$

$$f(n) \geq c g(n)$$

So, $f(n)$ is $\Omega(g(n))$ exactly when $g(n)$ is $O(f(n))$

The previous slide says that all three sorting algorithms have time complexity

- $O(n^2)$: Never use more than cn^2 operations
- $\Omega(n^2)$: Sometimes use at least cn^2 operations

When $f(n)$ is $O(g(n))$ and $f(n)$ is $\Omega(g(n))$ we write:

$$f(n) \text{ is } \Theta(g(n))$$

Comparators

- Limitations with Comparable interface?
 - Comparable permits 1 order between objects
 - What if compareTo() isn't the desired ordering?
 - What if Comparable isn't implemented?
- Solution: Comparators

Comparators (Ch 6.8)


- A comparator is an object that contains a method that is capable of comparing two objects
- Sorting methods can be written to apply a Comparator to two objects when a comparison is to be performed
- Different comparators can be applied to the same data to sort in different orders or on different keys

```
public interface Comparator <E> {  
    // pre: a and b are valid objects  
    // post: returns a value <, =, or > than 0 determined by  
    // whether a is less than, equal to, or greater than b  
    public int compare(E a, E b);  
}
```

Example

Note that Patient does
not implement
Comparable or
Comparator!

```
class Patient {  
    protected int age;  
    protected String name;  
    public Patient (String n, int a) { name = n; age = a; }  
    public String getName() { return name; }  
    public int getAge() { return age; }  
}
```



```
class NameComparator implements Comparator <Patient>{  
    public int compare(Patient a, Patient b) {  
        return a.getName().compareTo(b.getName());  
    }  
    // Note: No constructor; a "do-nothing" constructor is added by Java  
}
```

```
public void sort(T a[], Comparator<T> c) {  
    ...  
    if (c.compare(a[i], a[max]) > 0) {...}  
}
```

```
sort(patients, new NameComparator());
```

Comparable vs Comparator

- Comparable Interface for class X
 - Permits just one order between objects of class X
 - Class X must implement a compareTo method
 - Changing order requires rewriting compareTo
 - And then recompiling class X
- Comparator Interface
 - Allows creation of “compator classes” for class X
 - Class X isn’t changed or recompiled
 - Multiple Comparators for X can be developed
 - Ex: Sort Strings by length (alphabetically for same-length)
 - Ex: Sort names by last name instead of first name

Selection Sort with Comparator

```
public static <E> int findPosOfMax(E[] a, int last,
                                   Comparator<E> c) {
    int maxPos = 0      // A wild guess
    for(int i = 1; i <= last; i++)
        if (c.compare(a[maxPos], a[i]) < 0)
            maxPos = i;
    return maxPos;
}

public static <E> void selectionSort(E[] a, Comparator<E> c) {
    for(int i = a.length - 1; i > 0; i--) {
        int big= findPosOfMin(a,i,c);
        swap(a, i, big);
    }
}
```

- The same array can be sorted in multiple ways by passing different Comparator<E> values to the sort method;

Merge Sort

- A **divide and conquer** algorithm
- Merge sort works as follows:
 - **Base case:**
 - If the list is of length 0 or 1, then it is already sorted. Return the sorted list.
 - Divide the unsorted list into two sublists of about half the size of original list.
 - **Recursive call:**
 - Sort each sublist by re-applying merge sort.
 - Merge the two sublists back into one sorted list.

Merge Sort

- [8 14 29 1 17 39 16 9]
- [8 14 29 1] [17 39 16 9] split
- [8 14] [29 1] [17 39] [16 9] split
- [8] [14] [29] [1] [17] [39] [16] [9] split
- [8 14] [1 29] [17 39] [9 16] merge
- [1 8 14 29] [9 16 17 39] merge
- [1 8 9 14 16 17 29 39] merge

Transylvanian Merge Sort Folk Dance

Merge Sort

- How would we implement it?
- Pseudocode:

```
//recursively mergesorts A[from..To] “in place”  
void recMergeSortHelper(A[], int from, int to)  
    if ( from < to )  
        // find midpoint  
        mid = (from + to)/2  
        //sort each half  
        recMergeSortHelper(A, from, mid)  
        recMergeSortHelper(A, mid+1, to)  
        // merge sorted lists  
        merge(A, from, to)
```

But `merge` hides a number of important details.... 25