

[TAP:WNUME] Big-O

```
public static boolean contains(int[] nums, int x) {  
    return containsHelper(nums, x, nums.length);  
}  
private static boolean containsHelper(int[] nums, int x,  
    int curIdx){  
    if (curIdx == 0)  
        return false;  
    return nums[curIdx]==x || containsHelper(nums, x, curIdx-1);  
}
```

- What is the time complexity of the code above?

- > A. $O(n)$
- > B. $O(\log n)$
- > C. $O(n \log n)$
- D. $O(n^2)$
- E. Whatever

$$\begin{aligned} T(n) &= k + T(n-1) \\ &= k + k + T(n-2) \\ &\vdots \\ &= k + k + \dots + T(0) \\ &= O(n) \end{aligned}$$

Handwritten notes:
- An arrow points from "nums.length" in the code to n in the equation.
- n is written as $\frac{n}{2}$ above $n-1$.
- n is written as $\frac{n}{2^m}$ above $n-2$.
- A bracket under the final $O(n)$ is labeled $O(\log n)$ with $m = \log n$ written below it.

$$\begin{aligned} \frac{n}{2^m} &= 1 \\ n &= 2^m \\ \log n &= m \end{aligned}$$

Administrative Details

- Lab 1
 - I apologize for not having it returned yet
 - Feedback will show up on github as a Pull Request (PR)
 - PRs give you the option to view comments line-by-line, and respond to comments
- (New workflow this semester, so it is taking time to get the kinks worked out. It should be faster turnaround than printouts once it is working.)

Agenda

- ⦿ Induction
 - List

(Proof by) Induction

- The mathematical cousin of recursion is induction:

Base case:

- prove for 0, 1, ...

Inductive case:

- assume true for $n-1$
prove true for n

Recursion

- In recursion, we always use the same basic approach/structure
 - base case
 - recursive case

Mathematical Induction

- Prove that for every $n \geq 0$

$$2^0 + 2^1 + 2^2 + \dots + 2^n = 2^{n+1} - 1$$

Base case:

$$n=0: 2^0 = 2^{0+1} - 1$$

$$1 = 2 - 1$$

$$1 = 1 \quad \checkmark$$

Inductive case:

$$\text{Assume } 2^0 + 2^1 + \dots + 2^{n-1} = 2^{n-1+1} - 1 \quad (\text{IH})$$

$$\text{Show } 2^0 + \dots + 2^n = 2^{n+1} - 1$$

$$\begin{aligned} 2^0 + \dots + 2^n &= 2^n - 1 + 2^n && \text{by IH} \\ &= 2(2^n) - 1 \\ &= 2^{n+1} - 1 \quad \checkmark \quad \square \end{aligned}$$

Mathematical Induction

- Prove that for every $n \geq 0$

$$0 + 1 + \dots + n = \frac{n(n+1)}{2}$$

$n=0$ $0 = \frac{0(0+1)}{2}$ ✓

assume $0 + \dots + (n-1) = \frac{(n-1)(n-1+1)}{2}$ (IH)

show $0 + \dots + n = \frac{n(n+1)}{2}$

$0 + \dots + n = \frac{n(n-1)}{2} + n$ by IH

$$= \frac{n(n-1)}{2} + \frac{2n}{2}$$

$$= \frac{n^2 + n}{2}$$

$$= \frac{n(n+1)}{2}$$

□

Agenda

- Induction
- ⦿ List

The List Interface

```
interface List {  
    size()  
    isEmpty()  
    contains(e)  
    get(i)  
    set(i, e)  
    add(i, e)  
    remove(i)  
    addFirst(e)  
    getLast()  
    .  
    .  
    .  
}
```

Vector implements List

*·
· " "
·
·*

Singly Linked List

Pros and Cons of Vectors

Pros *constant time*
Random access

- Fast access to elements

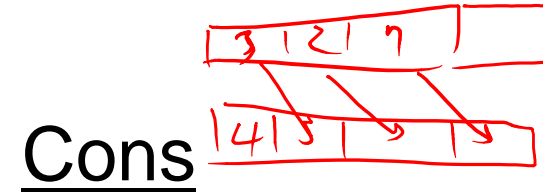
Array

- An array is stored in consecutive memory locations:

```
int[] nums;  
nums = new int[5];
```

*mem location = mem location of Array + (index * size of type)*

9



Cons

- Slow updates to front of list
- Hard to predict time for add (depends on internal array size)
- Potentially wasted space



- Dynamically Resizable? *yes, but inefficient*

Singly Linked List

- There are two key components of Lists

- The list itself

- Instance variables

- (Pointer to) the head node of the list

of elements

- Methods

- Those declared in the List interface

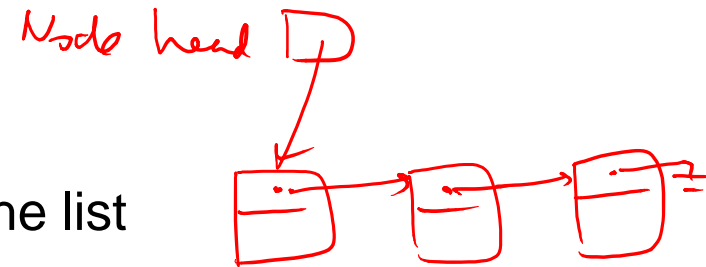
- Nodes

- Instance variables

- data
- (Pointer to) the “next” element

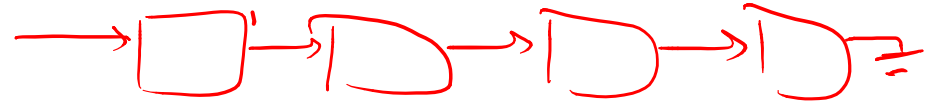
- Define methods

- Getters and setters



Singly Linked List Methods

```
head != null  
// pre: 0 ≤ index < numElements  
public E get(int index) {  
    Node finger = head;  
    for (int i = 0; i < index; i++)  
        finger = finger.next();  
    return finger.value();  
}
```



Singly Linked List Methods

```
public E set(E d, int index) {
```

]

```
E ori = finger.value();  
finger.setValue(d);
```

```
return ori;
```

```
}
```



Singly Linked List Methods

```
public void add(E d, int index) {
```

