# CSCI 136 Data Structures \& Advanced Programming 

Lecture 30
Fall 2017
Instructors: Bills are Back

## Last Time

- Graph Data Structures: Implementation
- Adjacency Array Implementation Details
- Featuring many Iterators!


## Today's Outline

- Greedy Algorithms for Optimization
- Lab II : Exam Scheduling
- Defining the problem
- Sketching a design
- Adjacency List Implementation Details
- More Fundamental Graph Properties
- An Important Algorithm: Minimum-cost spanning subgraph


## Lab II Overview: <br> Graph Algorithms using structure5

## Greedy Algorithms

- A greedy algorithm attempts to find a globally optimum solution to a problem by making locally optimum (greedy) choices
- Example: Graph Coloring
- A (proper) coloring of a graph $G=(V, E)$ is an assignment of a value (color) to each vertex so that adjacent vertices get different values (colors)
- Typically one strives to minimize the number of colors used


## Greedy Coloring : Math

Here's a greedy coloring algorithm
Build a collection $C=\left\{C_{l}, \ldots, C_{k}\right\}$ of sets of vertices
$i=0 ; C_{i}=\{ \} / /$ empty set
while $G$ is has more vertices
for each vertex u in $G$
if $u$ is not adjacent to any vertex of $C_{i}$ remove u from $G$ and add u to $C_{i}$
add $C_{i}$ to $C$
$i^{++}$;
Return C as the coloring

## Greedy Coloring : CS

Here's a greedy coloring algorithm
Create a structure C to hold a collection of lists while $G$ is not empty
pick a vertex $v$ in $G$; create an empty list $L$; add v to $L$ for each vertex $u \neq v$ in $G$
if $u$ is not adjacent to any vertex of $L$ add u to $L$
remove all vertices of $L$ from $G$ add L to $C$
Return C as the coloring

## Greedy Coloring



## Greedy Coloring

## Some observations

- Each list (color class) $L$ is a set of vertices no two of which are adjacent (an independent set)
- Each color class is maximal: cannot be made any larger
- The hope is that this results in fewer colors being needed
- But the solution is not always optimum!
- This is a very hard problem
- The coloring problem is the same as finding a partition of the vertex set into independent sets
- Partition means union of disjoint sets


## Lab II : Exam Scheduling

Find a schedule (set of time slots) for exams so that

- No student has two exams in the same slot
- Every course is in a slot
- The number of slots is as small as possible

This is just the graph coloring problem in disguise!

- Each course is a vertex
- Two vertices are adjacent if the courses share students
- A slot must be an independent set of vertices (that is, a color class)


## Lab II Notes: Using Graphs

- Create a new graph in structure5
- GraphListDirected, GraphListUndirected,
- GraphMatrixDirected, GraphMatrixUndirected
- Graph<V,E> conflictGraph = new GraphListUndirected<V,E>();


## Lab II : Useful Graph Methods

- void add(V label)
- add vertex to graph
- void addEdge(V vtx1, V vtx2, E label)
- add edge between vtxI and vtx2
- Iterator<V> neighbors(V vtx1)
- Get iterator for all neighbors to vtxl
- boolean isEmpty()
- Returns true iff graph is empty
- Iterator<V> iterator()
- Get vertex iterator
- V remove(V label)
- Remove a vertex from the graph
- E removeEdge(V vLabel1, V vLabel2)
- Remove an edge from graph


## Adjacency List : Directed Graph

| A | $\longrightarrow B$ | $\rightarrow \mathrm{C}$ | $\rightarrow$ G | $\longrightarrow \mathrm{H}$ |
| :---: | :---: | :---: | :---: | :---: |
| B | $\longrightarrow \mathrm{D}$ | $\rightarrow$ G | $\rightarrow \mathrm{H}$ |  |
| C | $\rightarrow$ B | $\rightarrow D$ |  |  |
| D |  |  |  |  |
| E | $\rightarrow D$ | $\rightarrow \mathrm{H}$ |  |  |
| F | $\rightarrow \mathrm{C}$ | $\rightarrow D$ |  |  |
| G | $\rightarrow F$ |  |  |  |
| H | $\rightarrow E$ |  |  |  |



The vertices are stored in an array V[]
V[] contains a linked list of edges having a given source

## Adjacency List : Undirected Graph




The vertices are stored in an array V[] V[] contains a linked list of edges incident to a given vertex

## GraphList

- Maintain an adjacency list of edges at each vertex (no adjacency matrix)
- Keep only outgoing edges for directed graphs
- Support both directed and undirected graphs (GraphListDirected, GraphListUndirected)


## Vertex and GraphListVertex

- We use the same Edge class for all graph types
- We extend Vertex to include an Edge list
- GraphListVertex class adds to Vertex class
- A Structure to store edges adjacent to the vertex protected Structure<Edge<V,E>> adjacencies; // adjacent edges
- adjacencies is created as a SinglyLinkedList of edges
- Several methods

```
public void addEdge(Edge<V,E> e)
public boolean containsEdge(Edge<V,E> e)
public Edge<V,E> removeEdge(Edge<V,E> e)
public Edge<V,E> getEdge(Edge<V,E> e)
public int degree()
// and methods to produce Iterators...
```


## GraphListVertex

```
public GraphListVertex(V key){
    super(key); // init Vertex fields
    adjacencies = new SinglyLinkedList<Edge<V,E>>();
}
public void addEdge(Edge<V,E> e){
    if (!containsEdge(e)) adjacencies.add(e);
}
public boolean containsEdge(Edge<V,E> e){
    return adjacencies.contains(e);
}
public Edge<V,E> removeEdge(Edge<V,E> e) {
    return adjacencies.remove(e);
}
```


## GraphListVertex Iterators

```
// Iterator for incident edges
public Iterator<Edge<V,E>> adjacentEdges() {
    return adjacencies.iterator();
}
// Iterator for adjacent vertices
public Iterator<V> adjacentVertices() {
        return new GraphListAIterator<V,E>
        (adjacentEdges(), label());
}
```

GraphListAlterator creates an Iterator over vertices based on the Iterator over edges produced by adjacentEdges ( )

## GraphListAlterator

## GraphListAlterator uses two instance variables

protected AbstractIterator<Edge<V,E>> edges; protected V vertex;

```
public GraphListAIterator(Iterator<Edge<V,E>> i, V v) {
    edges = (AbstractIterator<Edge<V,E>>)i;
    vertex = v;
}
public V next() {
    Edge<V,E> e = edges.next();
    if (vertex.equals(e.here()))
            return e.there();
        else { // could be an undirected edge!
        return e.here();
}
```


## GraphListElterator

GraphListEIterator uses one instance variable
protected AbstractIterator<Edge<V,E>> edges;
GraphListEIterator
-Takes the Map storing the vertices

- Uses it to build a linked list of all edges
-Gets an iterator for this linked list and stores it, using it in its own methods


## GraphList

- To implement GraphList, we use the GraphListVertex (GLV) class
- GraphListVertex class
- Maintain linked list of edges at each vertex
- Instance vars: label, visited flag, linked list of edges
- GraphList abstract class
- Instance vars:
- Map<V,GraphListVertex<V,E>> dict; // label -> vertex
- boolean directed; // is graph directed?
- How do we implement key GL methods?
- GraphList(), add(), getEdge(), ...

```
protected GraphList(boolean dir){
    dict = new Hashtable<V,GraphListVertex<V,E>>();
    directed = dir;
}
public void add(V label) {
    if (dict.containsKey(label)) return;
    GraphListVertex<V,E> v = new
        GraphListVertex<V,E>(label);
    dict.put(label,v);
}
public Edge<V,E> getEdge(V label1, V label2) {
    Edge<V,E> e = new Edge<V,E> (get(labell),
    get(label2), null, directed);
    return dict.get(label1).getEdge(e);
}
```


## GraphListDirected

- GraphListDirected (GraphListUndirected) implements the methods requiring different treatment due to (un)directedness of edges
- addEdge, remove, removeEdge, ...

```
// addEdge in GraphListDirected.java
// first vertex is source, second is destination
    public void addEdge(V vLabel1, V vLabel2, E label) {
        // first get the vertices
        GraphListVertex<V,E> v1 = dict.get(vLabel1);
        GraphListVertex<V,E> v2 = dict.get(vLabel2);
        // create the new edge
        Edge<V,E> e = new Edge<V,E>(v1.label(), v2.label(), label, true);
        // add edge only to source vertex linked list (aka adjacency list)
        v1.addEdge(e);
}
```

```
public V remove(V label) {
    //Get vertex out of map/dictionary
    GraphListVertex<V,E> v = dict.get(label);
    //Iterate over all vertex labels (called the map "keyset")
    Iterator<V> vi = iterator();
    while (vi.hasNext()) {
        //Get next vertex label in iterator
        V v2 = vi.next();
        //Skip over the vertex label we're removing
        //(Nodes don't have edges to themselves...)
        if (!label.equals(v2)) {
        //Remove all edges to "label"
        //If edge does not exist, removeEdge returns null
        removeEdge(v2,label);
    }
    }
    //Remove vertex from map
    dict.remove(label);
    return v.label();
}
```

```
public E removeEdge(V vLabel1, V vLabel2) {
    //Get vertices out of map
    GraphListVertex<V,E> v1 = dict.get(vLabel1);
    GraphListVertex<V,E> v2 = dict.get(vLabel2);
    //Create a "temporary" edge connecting two vertices
    Edge<V,E> e = new Edge<V,E>(v1.label(), v2.label(), null, true);
    //Remove edge from source vertex linked list
    e = v1.removeEdge(e);
    if (e == null) return null;
    else return e.label();
}
```

