

# CSCI 136

## Data Structures & Advanced Programming

Lecture 26

Fall 2017

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# Administrative Details

- Lab 10: Two Towers is online
  - Individual lab again this week
- Final Exam location: TBL 112
  - It's on Dec. 14<sup>th</sup>, 9:30--noon

# Last Time

- Binary Search Tree Implementation details
- *Balanced* Binary search trees
  - AVL Trees
    - Height is  $O(\log n)$ , so all operations are  $O(\log n)$
  - Red-Black Trees
    - Different height-balancing idea: height is  $O(\log n)$
    - All operations are  $O(\log n)$
  - Splay Trees
    - No guaranteed balance; good *amortized* performance
    - Any sequence of  $m$  operations take  $O(m \log n)$  time

# Today's Outline

Much less esoteric...

- Bit operations
  - Useful in general, but required for Lab 10
- Introduction To Graphs
  - Basic Definitions and Properties
  - Applications and Problems

# Representing Numbers

- Humans usually think of numbers in base 10
- But even though we write `int x = 23;` the computer stores `x` as a sequence of 1s and 0s

- Recall Lab 3:

```
public static String printInBinary(int n) {  
    if (n <= 1)  
        return "" + n%2;  
  
    return printInBinary(n/2)+n%2;  
}
```

- 00000000 00000000 00000000 00010111

# Bitwise Operations

- We can use *bitwise* operations to manipulate the 1s and 0s in the binary representation
  - Bitwise 'and': &
  - Bitwise 'or': |
- Also useful: bit shifts
  - Bit shift left: <<
  - Bit shift right: >>

# & and |

- Given two integers  $a$  and  $b$ , the bitwise *or* expression  $a | b$  returns an integer s.t.
  - At each bit position, the result has a 1 if that bit position had a 1 in EITHER  $a$  OR  $b$
  - $3 | 6 = ?$
- Given two integers  $a$  and  $b$ , the bitwise *and* expression  $a \& b$  returns an integer s.t.
  - At each bit position, the result has a 1 if that bit position had a 1 in BOTH  $a$  AND  $b$
  - $3 \& 6 = ?$

## >> and <<

- Given two integers  $a$  and  $i$ , the expression  $(a \ll i)$  returns  $(a * 2^i)$ 
  - Why? It shifts all bits **left** by  $i$  positions
  - $1 \ll 4 = ?$
- Given two integers  $a$  and  $i$ , the expression  $(a \gg i)$  returns  $(a / 2^i)$ 
  - Why? It shifts all bits **right** by  $i$  positions
  - $1 \gg 4 = ?$
  - $97 \gg 3 = ?$  ( $97 = 1100001$ )
- Be careful about shifting left and “overflow”!!!



# Revisiting `printlnBinary(int n)`

- How would we rewrite a recursive `printlnBinary` using bit shifts and bitwise operations?

```
public static String printlnBinary(int n) {  
    if (n <= 1) {  
        return "" + n;  
    }  
    return printlnBinary(n >> 1) + (n & 1);  
}
```

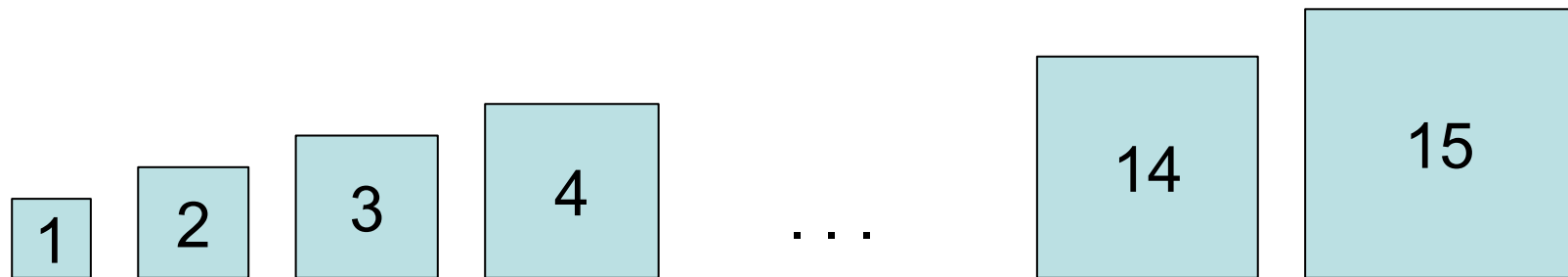
# Revisiting `printlnBinary(int n)`

- How would we write an iterative `printlnBinary` using bit shifts and bitwise operations?

```
public static String printlnBinary(int n,
                                   int width) {
    String result = "";
    for(int i = 0; i < width; i++)
        if ((n & (1<<i)) == 0)
            result = 0 + result;
        else
            result = 1 + result;
    return result;
}
```

# Lab 8: Two Towers

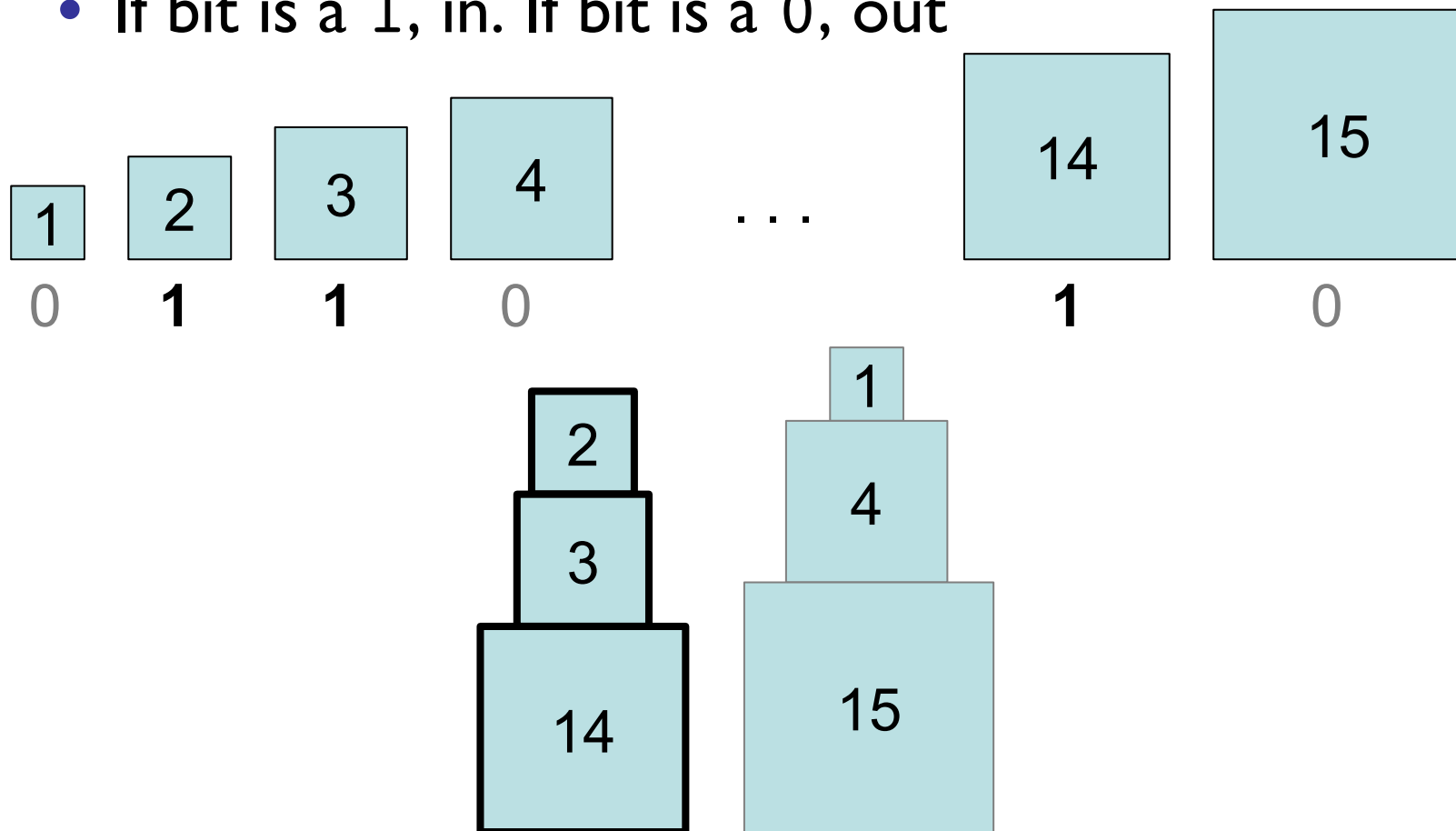
- **Goal:** given a set of blocks, iterate through all possible subsets to find the *best* set



- “Best” set produces the most balanced towers
- **Strategy:** create an iterator that uses the bits in a binary number to represent subsets

# Lab 8: Two Towers

- A block can either be in the set or out
  - If bit is a 1, in. If bit is a 0, out



# Questions?

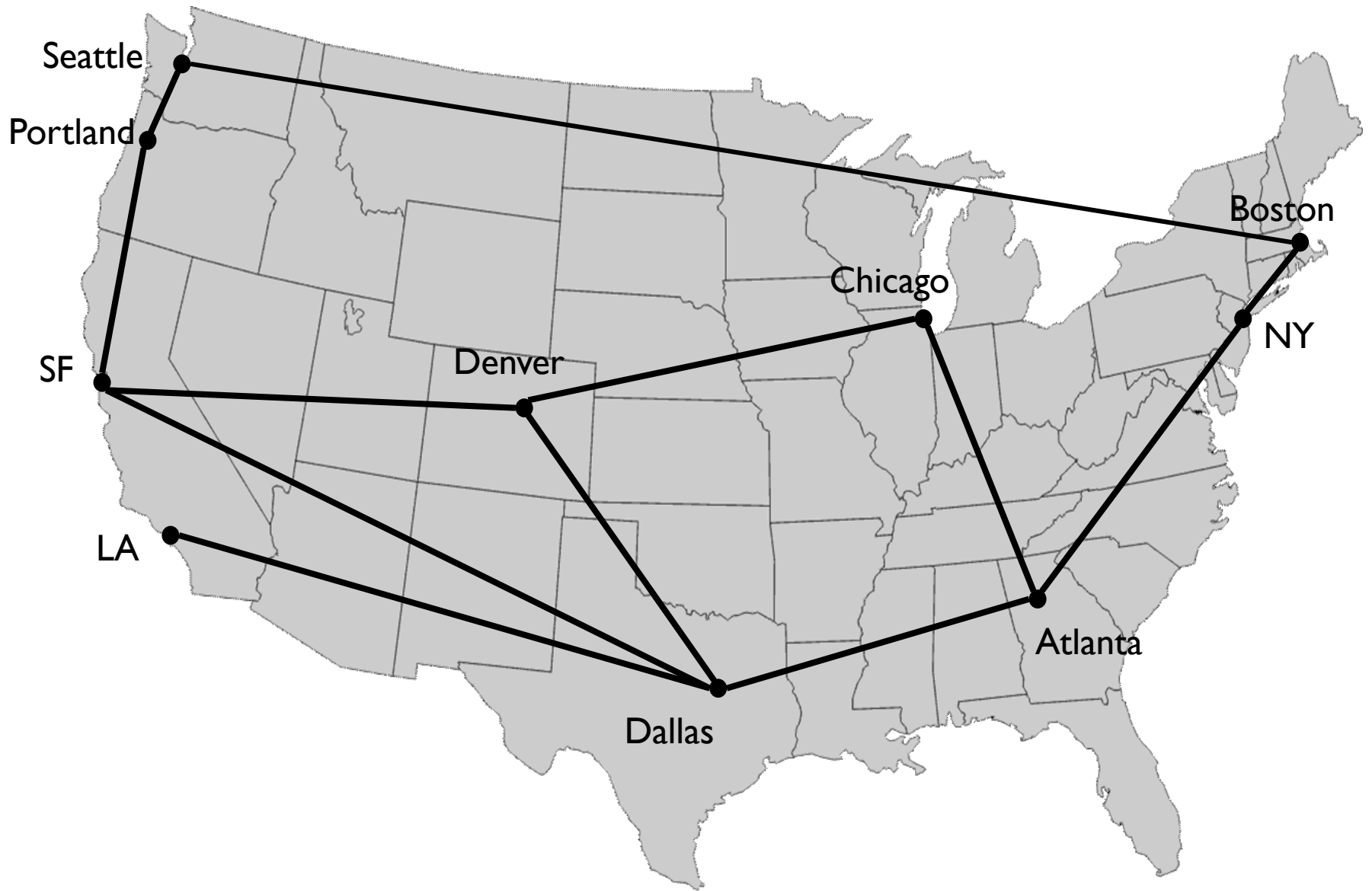
- We will write a “SubsetIterator” to enumerate all possible subsets of a `Vector<E>`
- We will use `SubsetIterator` to solve two problems
  - Two Towers
  - Identify all Subsequences of a String that are words
    - Use your `LexiconTrie!` (or an `OrderedStructure`)

# Graphs Describe the World<sup>1</sup>

- Transportation Networks
- Communication Networks
- Social Networks
- Molecular structures
- Dependency structures
- Scheduling
- Matching
- Graphics Modeling
- ....

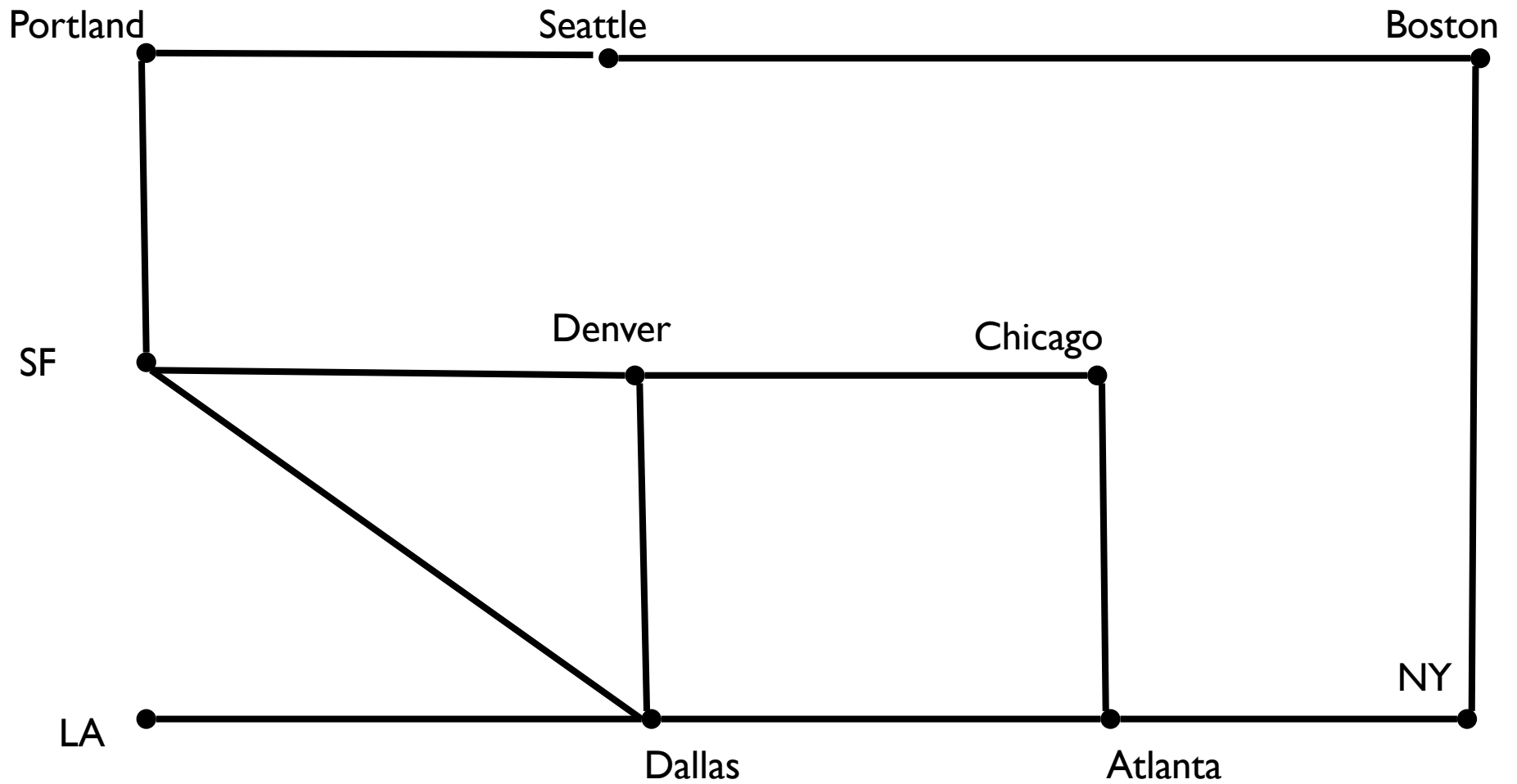


Nodes = subway stops; Edges = subway lines



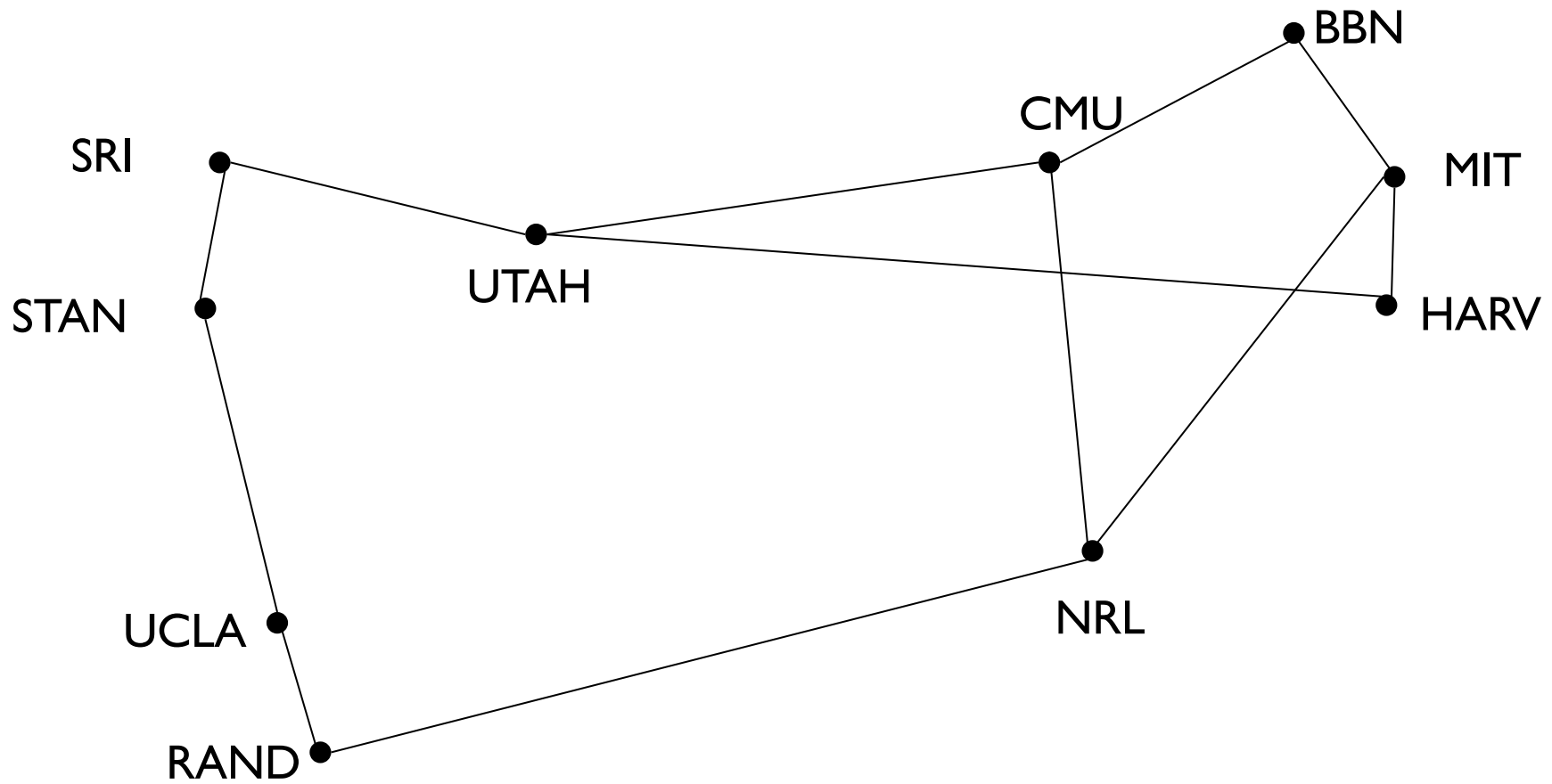
Nodes = cities; Edges = rail lines connecting cities



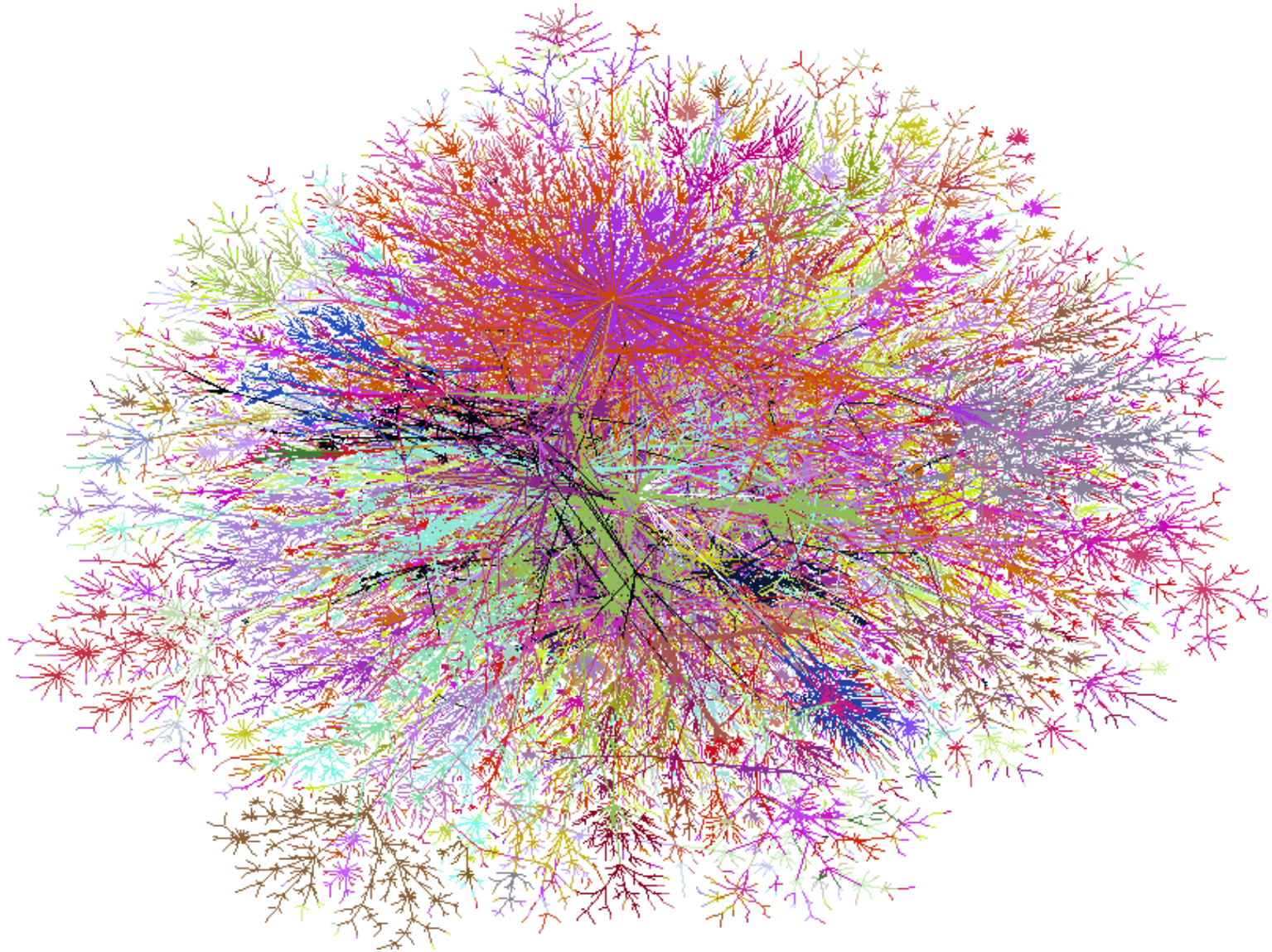


Note: Connections in graph matter, not precise locations of nodes

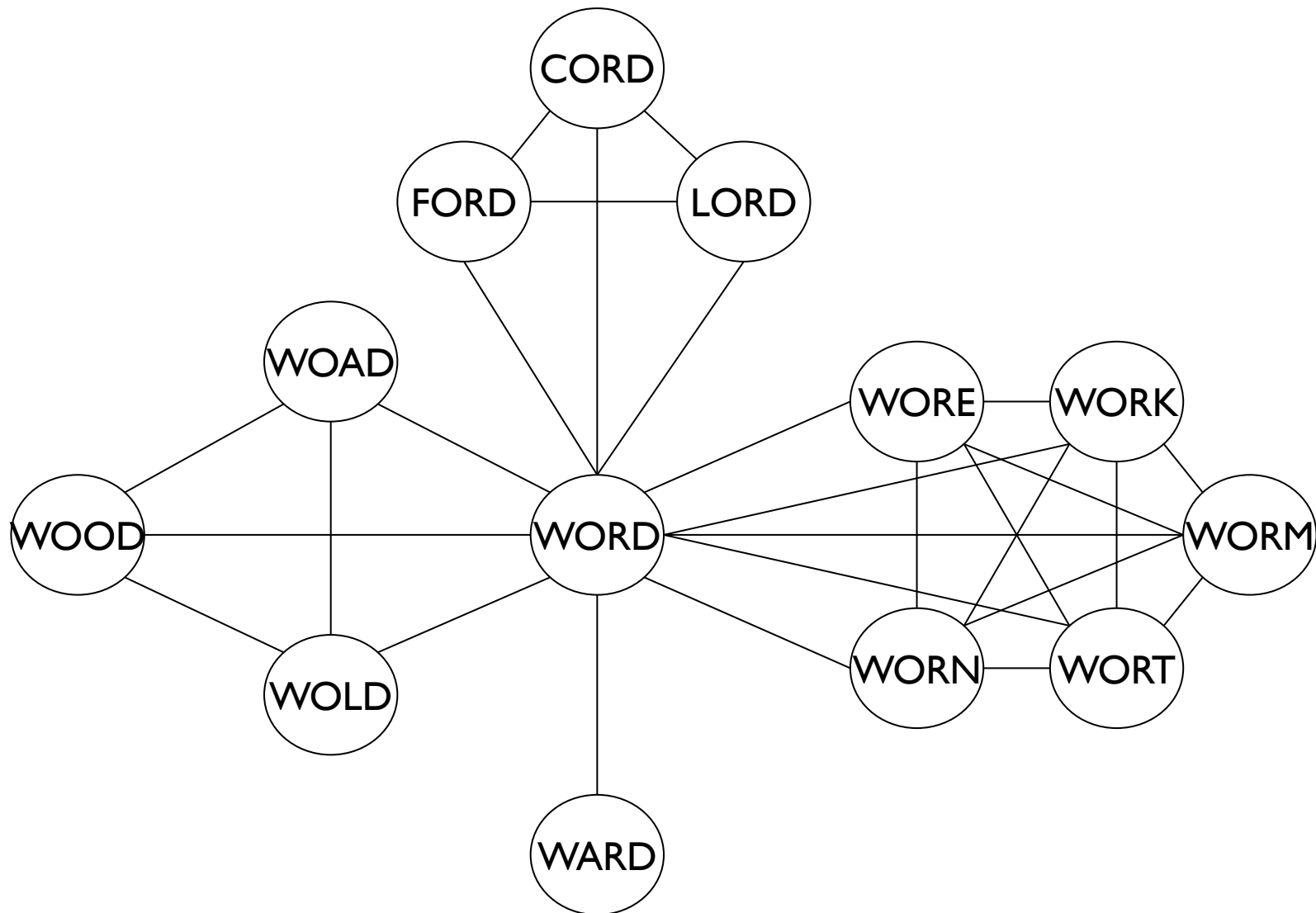
# Internet (~1972)



# Internet (~1998)

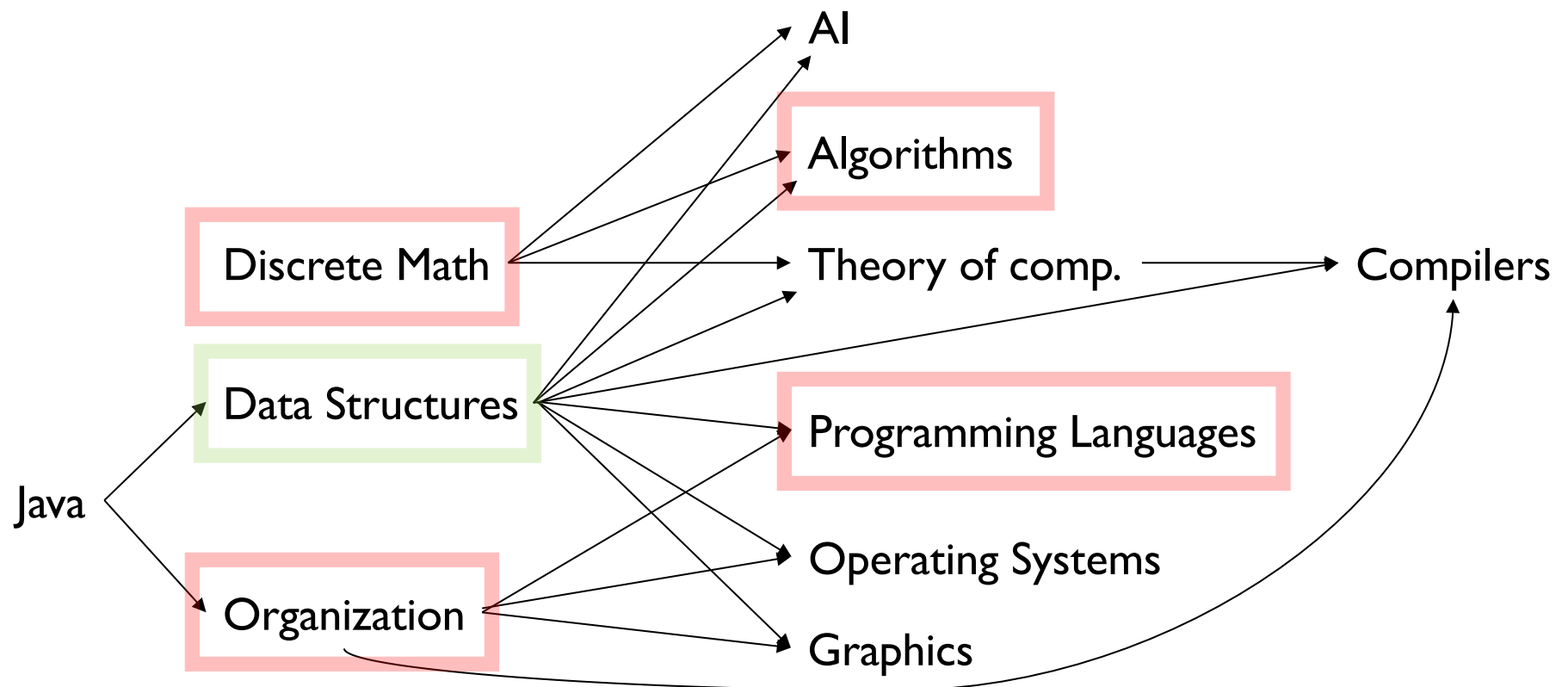


# Word Game



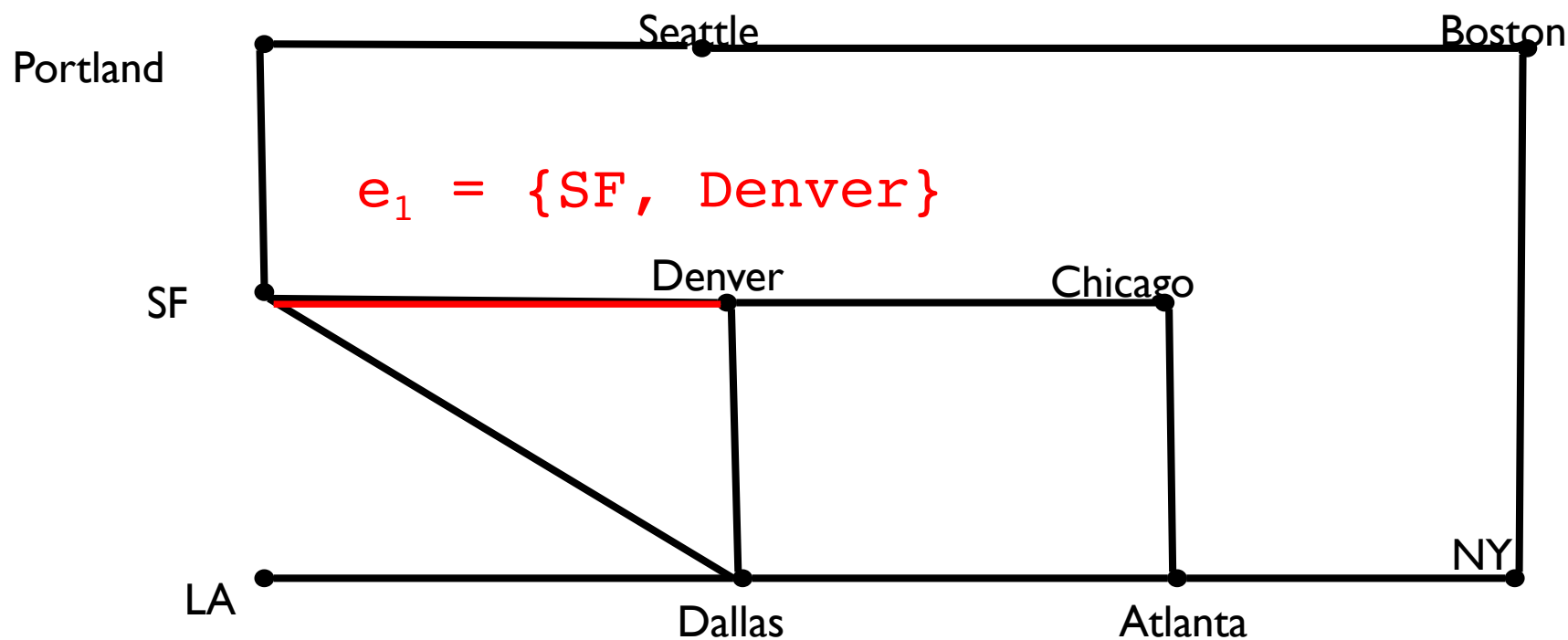
Nodes = words; Edges = words that differ by exactly one letter

# CS Pre-requisite Structure (subset)



Nodes = courses; Edges = prerequisites \*\*\*

# Basic Definitions & Concepts



**Definition:** An *undirected graph*  $G = (V, E)$  consists of two sets

- $V$  : the *vertices* of  $G$ , and  $E$  : the *edges* of  $G$
- Each edge  $e$  in  $E$  is defined by a set of two vertices: its *incident vertices*.
- We write  $e = \{u, v\}$  and say that  $u$  and  $v$  are *adjacent*.

# Walking Along a Graph

- A *walk* from  $u$  to  $v$  in a graph  $G = (V, E)$  is an *alternating* sequence of vertices and edges

$$u = v_0, e_1, v_1, e_2, v_2, \dots, v_{k-1}, e_k, v_k = v$$

such that each  $e_i = \{v_i, v_{i+1}\}$  for  $i = 1, \dots, k$

- Note a walk starts and ends on a vertex
- If no *edge* appears more than once then the walk is called a *path*
- If no *vertex* appears more than once then the walk is a *simple path*

# Walking In Circles

- A *closed walk* in a graph  $G = (V, E)$  is a walk  
 $v_0, e_1, v_1, e_2, v_2, \dots, v_{k-1}, e_k, v_k$   
such that  $v_0 = v_k$  (it ends at the starting  $v$ )
- A *circuit* is a *path* where  $v_0 = v_k$ 
  - Circuit vs. closed walk? Circuit has no repeat edges
- A *cycle* is a *simple path* where  $v_0 = v_k$ 
  - Circuit vs. cycle? Cycle has no repeated vertices.
- The *length* of any of these is the number of *edges* in the sequence