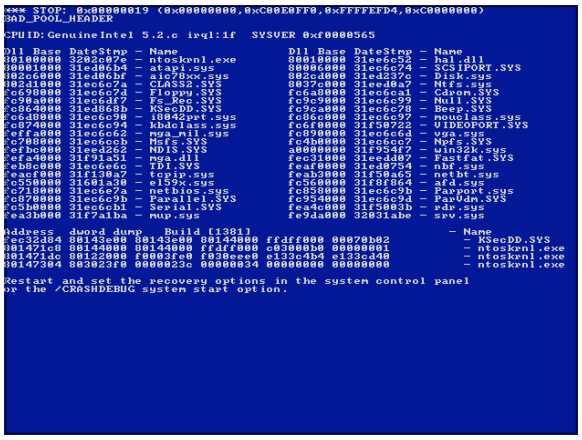
Squashing the Bugs: Tools for Building Better Software

Stephen Freund
Williams College

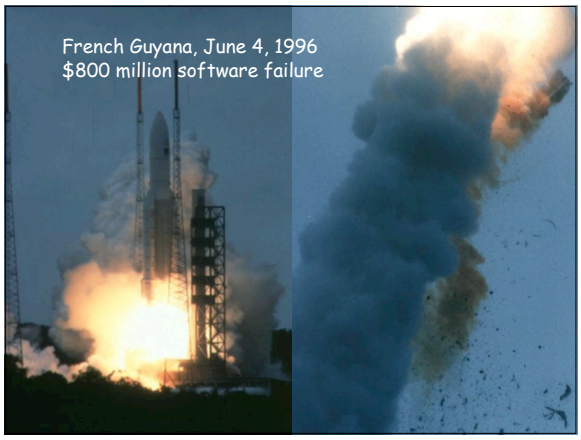


The Blue Screen of Death (BSOD)

USS Yorktown

- Smart Ship
  - 27 Pentium-based PCs
  - Windows NT 4.0
- September 21, 1997:
  - data entry error caused a "Divide-By-0" error
  - entire system failed
  - ship dead in the water for over 2 hours

[Wired 1997]



Mars Climate Orbiter

Purpose: Collect data. Relay signals from Mars Polar Lander ($165M)

Failure: Smashed into Mars (1999)

Bug: Failed to convert English to metric units

Mars Polar Lander

Purpose: Lander to study the Mars climate ($120M)

Failure: Smashed into Mars (2000)

Bug: Spurious signals f sensors caused premat engine shutoff

Therac25 Radiation Therapy

Purpose: Computer-controlled radiation therapy machine

Failure: gave fatal radiation doses to 2 cancer patients (1986)

Bug: timing bug

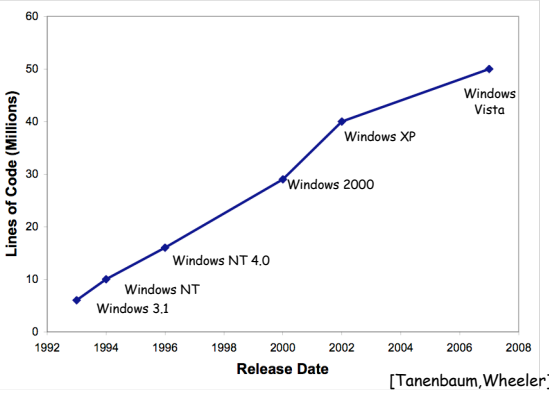
USS Vincennes

Failure: Shot down an Airbus jet that was mistaken for a F-14. 290 people died. (1988)

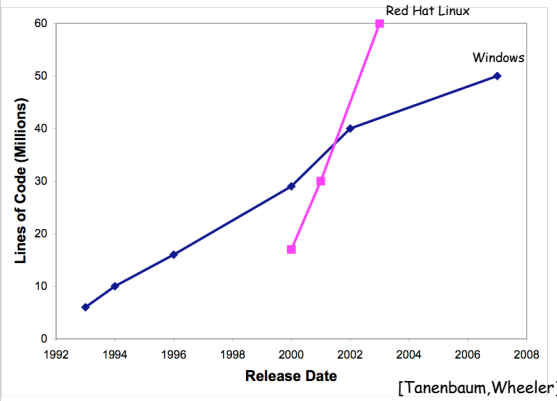
Bug: tracking software displayed cryptic and misleading output

Software Complexity

- Huge costs
  - $60 billion / year [NIST 2003]
  - convenience, security, safety
- Why is it so hard to get right?
  - systems are too large to understand completely
  - small mistakes can compromise whole system
  - many failure modes
    - unreliable network, media, hardware, users,...
    - half of code is to handle failures
  - programs evolve over time



[Tanenbaum,Wheeler]



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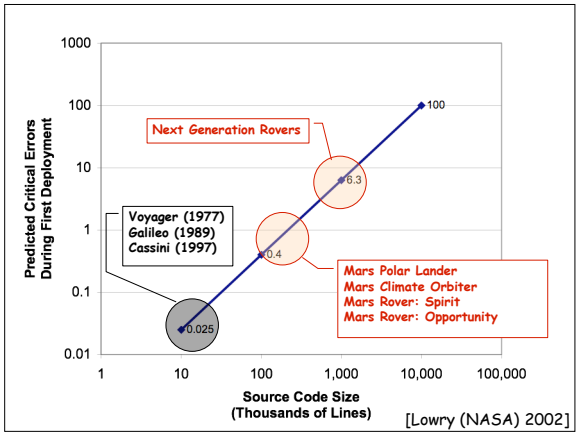
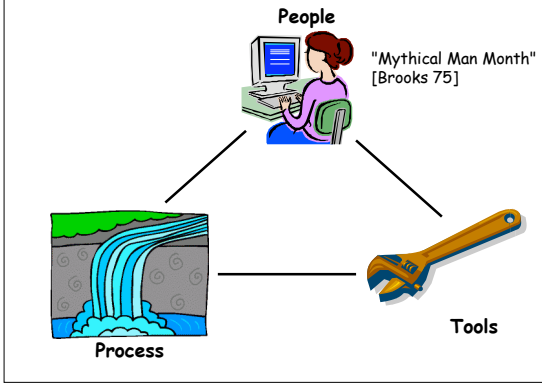


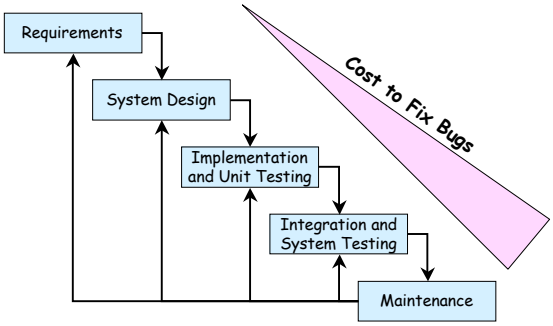
Managing Software Complexity



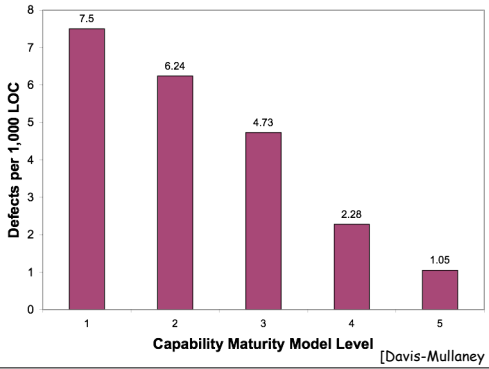
Software Processes

- Organization of large software project
  - other examples: army, company, college admin
  - planning, responsibilities, interactions
- Benefits
  - avoid miscommunication, misunderstandings
  - predict time and cost
  - recognize problems early
- Process failures do happen (frequently...)
  - Air Traffic Control system, $2.6Billion, 1983-1996

Waterfall Model

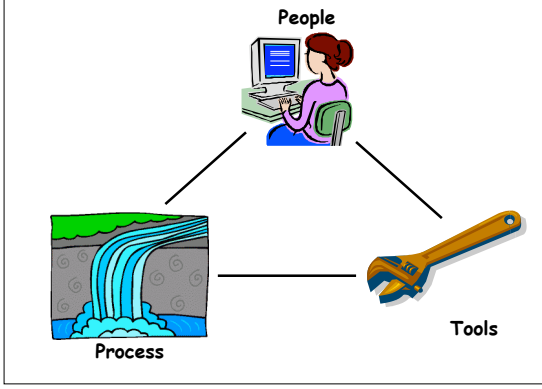


Defect Density in Delivered Software

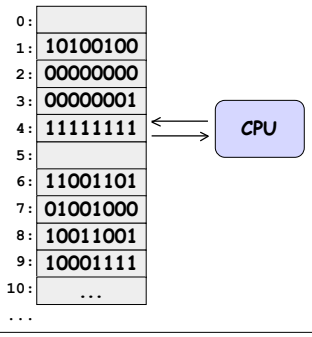


[Davis-Mullaney 2003]

Managing Software Complexity



Computer Architecture



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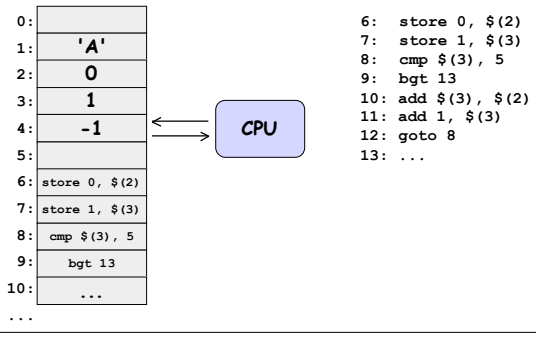

Abstraction In Programming Languages

```
sum = 0;
i = 1;
while (i <= 5) {
  sum = sum + i;
  i = i + 1;
}
```

Compiler

```
store 0, $(2)
store 1, $(3)
cmp $(3), 5
bgt 13
add $(3), $(2)
add 1, $(3)
goto 8
...
```

- High-level programming languages:
  - hide details of hardware
  - simplify programming process
- Can add language features to avoid mistakes

Memory Management

- Must allocate and free memory for arrays in C:

```
name = new char[8];
get_string(name);
authenticate(name);
free(name);
```

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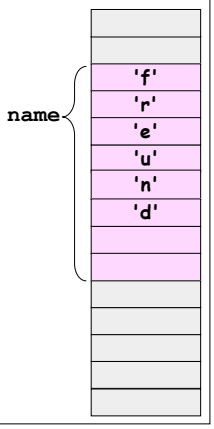
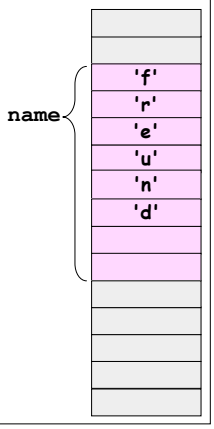


Memory Management

- Potential errors:
  - forget to free memory
  - free memory twice
  - free too early?
  - overwrite memory block



Garbage Collection

- Never explicitly free memory

```
name = new char[8];
get_string(name);
authenticate(name);
```

- Program periodically pauses to find and reclaim unused memory
- Used in Java, C#, Python, ...

Tradeoffs in Design

- No free lunch
  - features impact performance / expressiveness
  - can be prohibitively expensive (or impossible!)
- Other features
  - object-oriented languages / modules
  - exception handling
  - threads

Program Checking

- Design algorithm to automatically identify errors
- Example 1: type errors

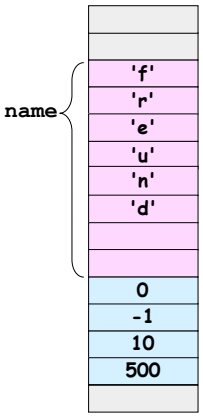
```
WebPage w = new WebPage("http://...");
int x = w - 3;   ← BAD

String s = "hello";
if (s < 100) ...  ← BAD
```

- Example 2: buffer overruns in C programs

Buffer Overruns in C

```
name = new char[8];
get_string(name);
authenticate(name);
free(name);
```



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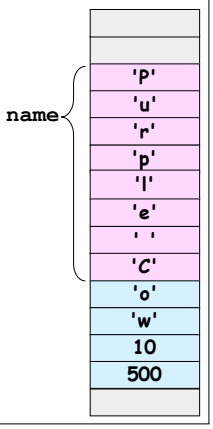
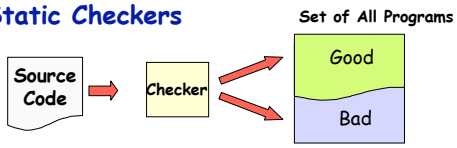


Dynamic Monitors

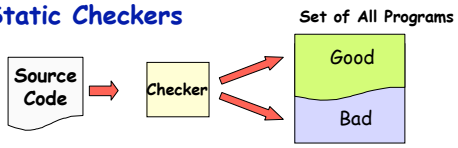
- Check program as it executes
  - example: buffer overruns
- Pros:
  - identify cause of bug faster than testing
  - easy to add to development process
- Cons:
  - coverage
  - performance
  - must run whole program

Static Checkers



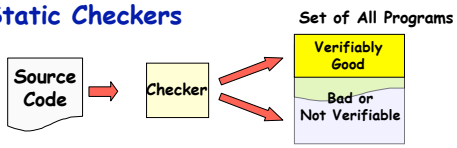
- Pros:
  - catch errors sooner (even before running)
  - check program for all inputs / all possible paths
- Cons:

Static Checkers



- Pros:
  - catch errors sooner (even before running)
  - check program for all inputs / all possible paths
- Cons:
  - cannot distinguish "good" from "bad" exactly
  - computers *cannot* compute everything
    - undecidability of the Halting Problem [Turing 1936]

Static Checkers



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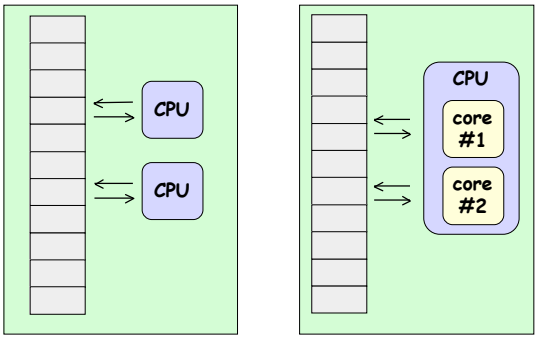
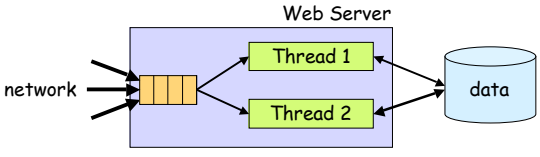
Multi-Processors and Multi-Core Chips



Concurrent Programming With Threads

- Decompose into pieces that run in parallel
- Improve throughput



Demo

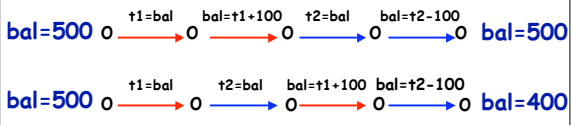
Multithreaded Program Execution

Thread 1

```
...
t1 = bal;
bal = t1 + 100;
...
```

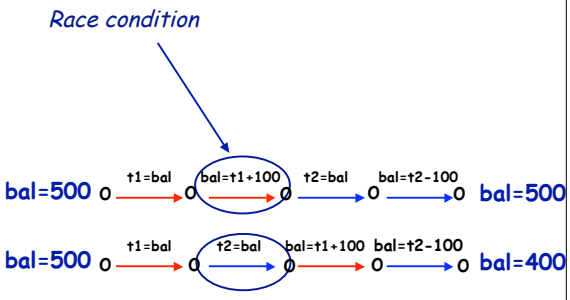
Thread 2

```
...
t2 = bal;
bal = t2 - 100;
...
```



Multithreaded Program Execution

Race condition

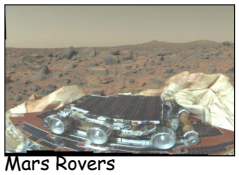


Race Conditions

- Common
- Hard to find via testing
  - scheduler dependent
- Memory, files, printers, ...

Therac-25

2003 Blackout ($6 Billion)



Mars Rovers

Preventing Race Conditions Using Locks

Thread 1

```
acquire(m);
t1 = bal;
bal = t1 + 100;
release(m);
```

Thread 2

```
acquire(m);
t2 = bal;
bal = t2 - 100;
release(m);
```

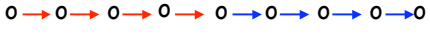
Preventing Race Conditions Using Locks

Thread 1

```
synchronized(m) {
  t1 = bal;
  bal = t1 + 100;
}
```

Thread 2

```
synchronized(m) {
  t2 = bal;
  bal = t2 - 100;
}
```



Demo

Checkers For Race Conditions

```
int bal; //# guarded_by m
```

Thread 1

```
synchronized(m) {
  t1 = bal;
  bal = t1 + 100;
}
```

Thread 2

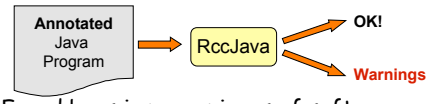
```
synchronized(m) {
  t2 = bal;
  bal = t2 - 100;
}
```

- Good: Threads always hold `m` when accessing `bal`
- Bad: Thread accesses `bal` without holding `m`

Tools: Eraser [Savage et al. 97], RccJava, ...

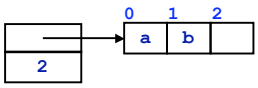
Statically Checking Real Systems

- RccJava [with Flanagan (UCSC), Peter Applegate '03]



- Found bugs in many pieces of software
  - commercial products, Java libraries, web servers
  - false positive rate of 75-80%
    - benign races, other forms of synchronization
    - better than not finding errors

java.util.Vector



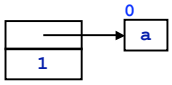

```
class Vector {
  Object elementData[] // guarded_by this
  int elementCount     // guarded_by this

  synchronized int lastIndexOf(Object elem, int n) {
    for (int i = n ; i >= 0 ; i--)
      if (elem.equals(elementData[i])) return i;
    return -1;
  }

  int lastIndexOf(Object elem) {
    return lastIndexOf(elem, elementCount - 1); // race!
  }

  synchronized void trimToSize() { ... }
  synchronized boolean remove(int index) { ... }
}
```

java.util.Vector

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  }

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}
```

Demo

Atomicity Violations

```
int bal; //# guarded_by m
```

Thread 1

```
synchronized(m) {
  t1 = bal;
}

synchronized(m) {
  bal = t1 + 100;
}
```

Thread 2

```
synchronized(m) {
  t2 = bal;
  bal = t2 - 100;
}
```

Atomicity Violations

```
int bal; //# guarded_by m
```

Thread 1

```
atomic {
  synchronized(m) {
    t1 = bal;
  }

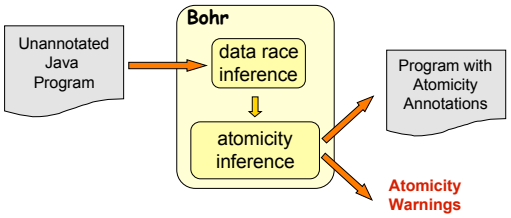
  synchronized(m) {
    bal = t1 + 100;
  }
}
```

Thread 2

```
atomic {
  synchronized(m) {
    t2 = bal;
    bal = t2 - 100;
  }
}
```

Bohr

- Compute `atomic` annotations automatically
- Identify methods that may suffer interference
- With Masha Lifshin '05



The (Long) Road to Reliable Software

- Bugs are a real problem
- Checking tools will improve life for everyone
- Industry starting to adopt checkers

- Lots of problems (and fun) left
  - tools often hard to use, imprecise
  - simple tools pave way for more sophisticated
  - need teachable design methodologies

Thanks

- Pete Applegate '03
- Masha Lifshin '05
- Cormac Flanagan (UCSC)
- Martín Abadi (UCSC and Microsoft)
- Shaz Qadeer (Microsoft)

- NSF/NASA HDCCSR Program